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# Databases Systems

## ICS 184

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# CS 184 Course Web Server

- ◆ All course information will be posted on line
- ◆ URL:
  - <http://www.ics.uci.edu/~dbclass/ics184/index.html>
- ◆ Class Notes available before class on the Web.

# Course Info

## ◆ TAs

- Koushik Niyogi

  - Office and office Hours:

    - ◆ Monday and Wednesday : 2:00-3:00 PM

- Rajat Mathur

  - Office and office Hours: TBA

## ◆ Instructor

- Office Hours:

- Tue 11-12 pm

- (send email)

- Email: [sharad@ics.uci.edu](mailto:sharad@ics.uci.edu)

- to contact me urgently, send email and mark subject line as CS 184 URGENT

# Desiderata

- ◆ Course Text: (either of following two books will suffice)
  - A First Course in Database Systems, Ullman and Widom
    - ◆ we will cover the entire book
  - Database Systems Concepts, Silberschatz, Korth, and Sudarshan
    - ◆ we will cover chapters 1-9
- ◆ Software:
  - Course will involve significant programming.
  - You will get exposure to database programming in DB2

# Desiderata (cont.)

## ◆ Course Requirements:

- Problem sets ~ approx. every week to 10 days.

- ◆ Total not to exceed 8

- Midterm

- Final (comprehensive)

- Grades:

- ◆ Problem sets 15%

- ◆ Personal Database Assignment (project) 15%

- ◆ Midterm 30%

- ◆ Final 40%

# Policies

## ◆ Late Assignments

- No grace period after due date... except under exceptional circumstances
- job interviews, out of town trip, breaks etc do not qualify as exceptional circumstances!

## ◆ Working in Groups

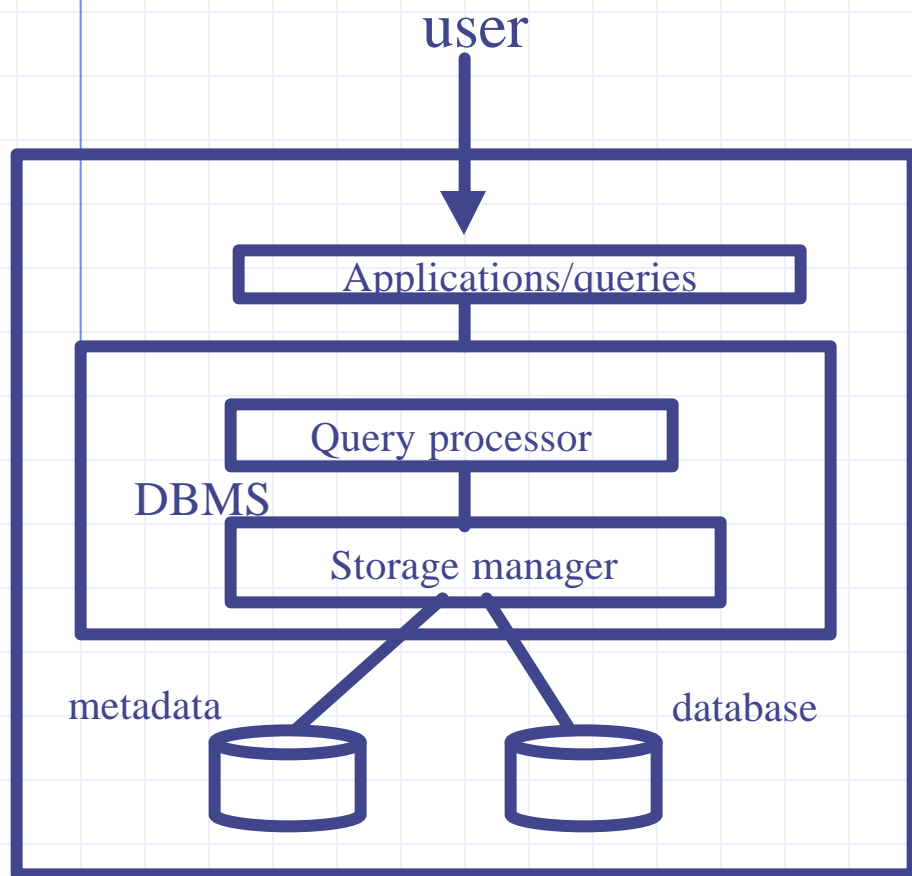
- do your homework problem sets in group size not to exceed 3
- learn more
- get better grades
- get used to working in groups (important to employers)

## ◆ Do exams individually!!

# Material Covered in CS 184

- ◆ Four aspects of studying DBMSs
  - Modeling and design of databases
    - ◆ allows exploration of issues before committing to an implementation
  - Programming: queries and DB operations like update.
    - ◆ SQL == “intergalactic dataspeak”
  - DBMS implementation
  - Effect of technology and application advances to database technology.
- ◆ CS 184 == (1) + (2)
- ◆ CS 214 == (3)
- ◆ CS215 == (3) + (4)

# Database Management Environment



**Database:** collection of interrelated information about world being modeled

**DBMS:** general purpose software to define, create, modify, retrieve, delete and manipulate a database

**Vendors:** Informix, Oracle, O2, Sybase, IBM, DEC



# Traditional DBMS Goals

- ◆ *Efficient management* of (faster than files) *large amounts of* (gigabytes) of *persistent* (outlasts creator), *reliable* (outlasts crashes) *shared information* (multiple users).
- ◆ DBMS Users:
  - small and large corporations
    - ◆ E-commerce companies, banks, airlines, transportation companies, corporate databases, government agencies, defense.
    - ◆ Anyone you can think of!

# Databases and File Systems

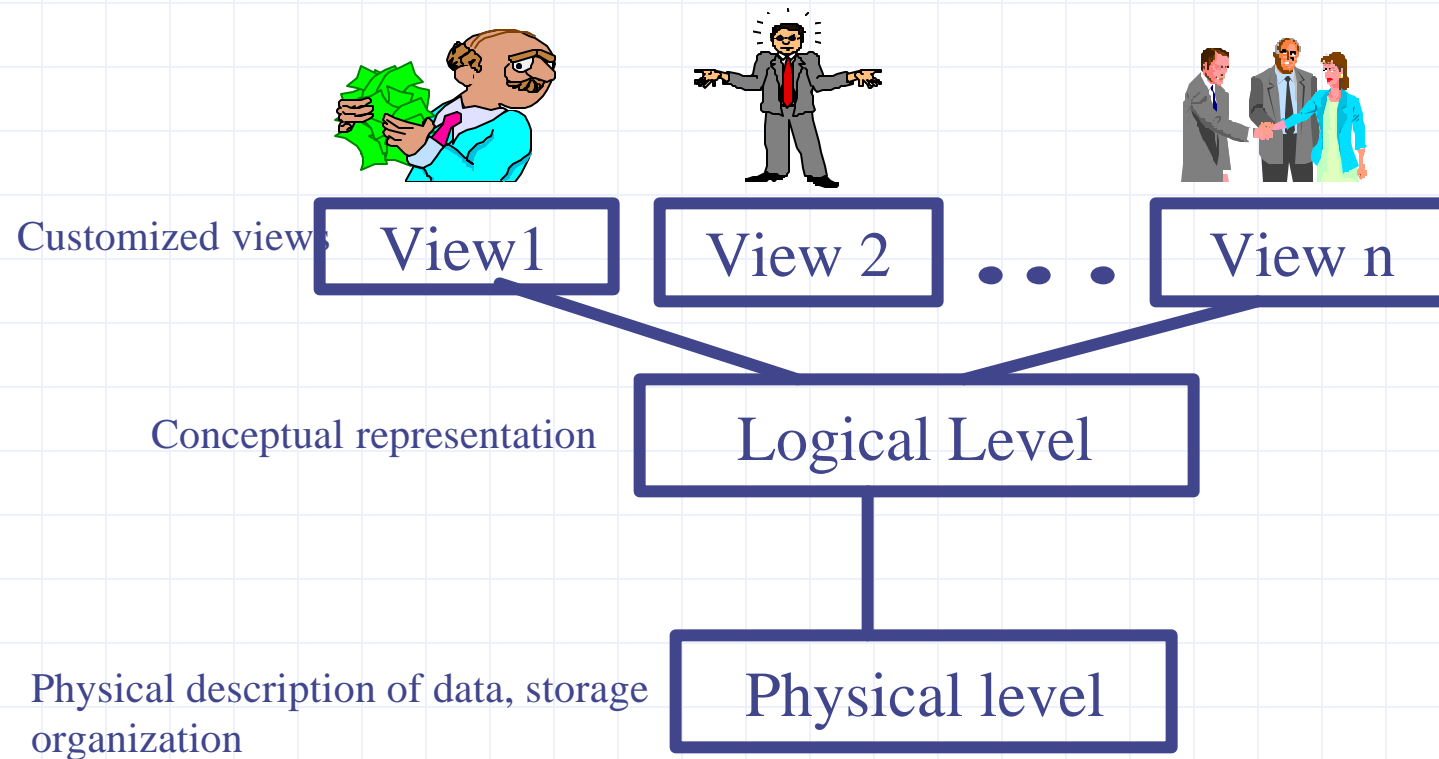
- ◆ DBMSs evolved from file systems.
- ◆ DBMSs provide many features that traditional file systems do not.
  - Support for concurrent access and data sharing. Data consistency in presence of concurrency
  - Reliability in presence of failures and system crashes.
  - Efficient associative access to very large amounts of data
  - A high level Query language (SQL) to define, create, access, and manipulate data. Support for unanticipated queries
  - support for multiple data views
  - security and authorization
  - data abstraction
  - prevention of data redundancy and inconsistencies

# Data Abstraction

- ◆ *program data independence*:
  - ability to hide details of how data is stored and maintained from application programs
- ◆ *program-operation independence*:
  - ability to hide details of operation implementation from application programs (Object-Orientation)

# Data Abstraction

- ◆ Hiding system complexity, physical storage details from users and application programs



# Schemas and Instances

- ◆ Instance:
  - set of data currently instantiated in database
  - changes frequently
- ◆ Schema:
  - overall design, structure, and constraints over the database
  - referred to as metadata
  - changes infrequently
- ◆ Example:

## Schema

### Tables

Emp (ename, dep#)

Dept(dep#, dname, mgr)

### Constraints

each department has  
a single manager

## Instance

Emp

(John, 10), (Cindy, 15), (Martha, 10)

dept

(10, Toy, John), (15, Sales, Cindy)

# Data Model

- ◆ Set of concepts and tools used to describe the database schema
- ◆ Different schemas at different abstraction levels:
  - *physical schema*: describes physical organization of data
  - *logical schema*: describes data at conceptual level
  - *sub schema*: defines data at view level
- ◆ Different models used describe schemas at different abstraction levels

# Types of Data Models

## ◆ Object based Logical Models

- Used to describe schema at view and logical levels.
- Support abstract view of data as objects, relationships, constraints
- Example: Entity Relationship Model, Functional data Model, Semantic Model, Object Oriented Model like ODL

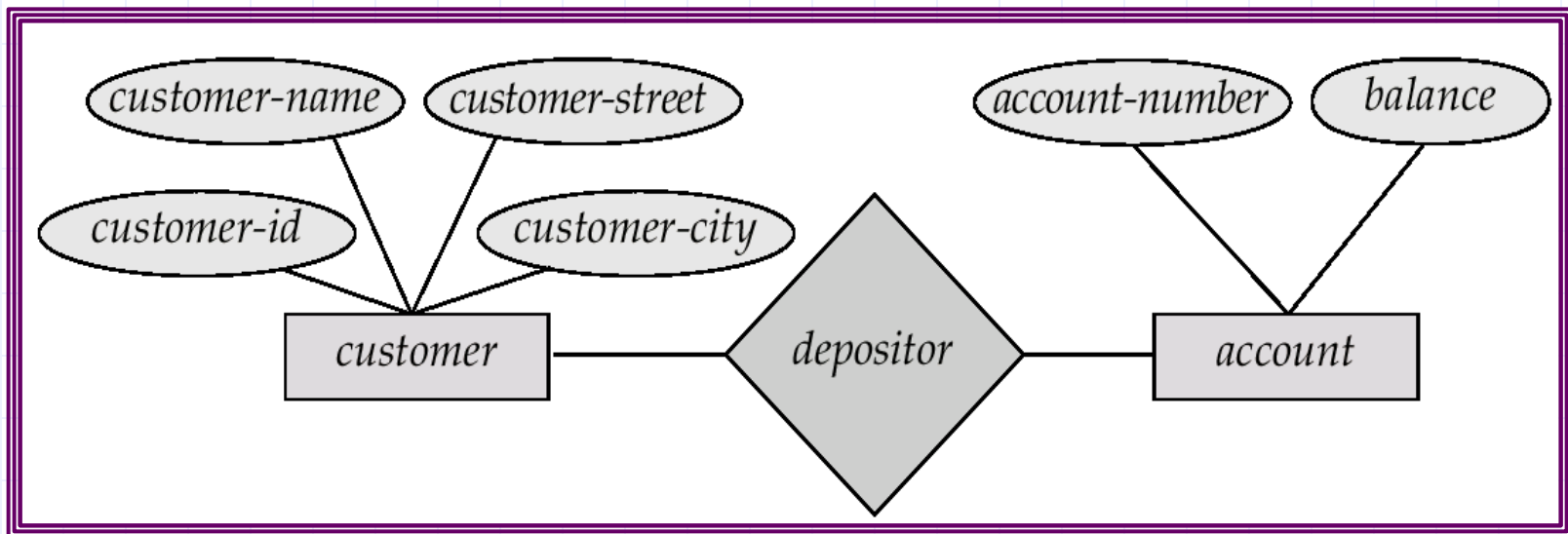
# Types of Data Model (cont.)

- ◆ Record-Based Logical Models
  - Used to define data at view and logical levels.
  - Provide a high level description of implementation
  - Examples: Relational Model, Hierarchical Model, Network Model
- ◆ Physical Models
  - Used to describe data at implementation level.
  - Examples: Frame Memory Model, Unifying Model



# Entity-Relationship Model

Example of schema in the entity-relationship model




# Entity Relationship Model (Cont.)

- ◆ E-R model of real world
  - Entities (objects)
    - ◆ E.g. customers, accounts, bank branch
  - Relationships between entities
    - ◆ E.g. Account A-101 is held by customer Johnson
    - ◆ Relationship set *depositor* associates customers with accounts
- ◆ Widely used for database design
  - Database design in E-R model usually converted to design in the relational model (coming up next) which is used for storage and processing

# Relational Model

- ◆ Example of tabular data in the relational model

Attributes



<i>Customer-id</i>	<i>customer-name</i>	<i>customer-street</i>	<i>customer-city</i>	<i>account-number</i>
192-83-7465	Johnson	Alma	Palo Alto	A-101
019-28-3746	Smith	North	Rye	A-215
192-83-7465	Johnson	Alma	Palo Alto	A-201
321-12-3123	Jones	Main	Harrison	A-217
019-28-3746	Smith	North	Rye	A-201

# A Sample Relational Database

<i>customer-id</i>	<i>customer-name</i>	<i>customer-street</i>	<i>customer-city</i>
192-83-7465	Johnson	12 Alma St.	Palo Alto
019-28-3746	Smith	4 North St.	Rye
677-89-9011	Hayes	3 Main St.	Harrison
182-73-6091	Turner	123 Putnam Ave.	Stamford
321-12-3123	Jones	100 Main St.	Harrison
336-66-9999	Lindsay	175 Park Ave.	Pittsfield
019-28-3746	Smith	72 North St.	Rye

(a) The *customer* table

<i>account-number</i>	<i>balance</i>
A-101	500
A-215	700
A-102	400
A-305	350
A-201	900
A-217	750
A-222	700

(b) The *account* table

<i>customer-id</i>	<i>account-number</i>
192-83-7465	A-101
192-83-7465	A-201
019-28-3746	A-215
677-89-9011	A-102
182-73-6091	A-305
321-12-3123	A-217
336-66-9999	A-222
019-28-3746	A-201

(c) The *depositor* table

# Classification of DBMSs based on Data Model

## ◆ Relational DBMSs:

- modeling concept: tables and constraints on tables
- Query Language: SQL
- Applications: suited for traditional business processing applications

## ◆ Object Oriented DBMSs

- modeling concepts: objects, classes, inheritance
- Query Language: object oriented OQL
- Applications: suited for CAD databases, CASE databases, office automation

## ◆ Object Relational DBMSs:

- incorporate OO concepts into relational model
- similar functionality as OODBMSs though different in implementations
- Language extended to process objects.

# DBMS Languages

- ◆ Data Definition Language (DDL)
  - DDL = the language used to describe a schema
  - Data dictionary/directory = a compiled description of a schema
- ◆ Data Manipulation Language (DML)
  - DML = Language users use to ask questions about (query) the database, and to change the data in the database.
- ◆ Storage Definition Language (SDL)
  - SDL = language to define the internal schema
- ◆ View Definition Language (VDL)
  - VDL = view definition language

# Data Definition Language (DDL)

- ◆ Specification notation for defining the database schema

- E.g.

```
create table account (  
    account-number char(10),  
    balance integer)
```

- ◆ DDL compiler generates a set of tables stored in a *data dictionary*
- ◆ Data dictionary contains metadata (i.e., data about data)
  - database schema
  - Data *storage and definition* language
    - ◆ language in which the storage structure and access methods used by the database system are specified
    - ◆ Usually an extension of the data definition language

# Data Manipulation Language (DML)

- ◆ Language for accessing and manipulating the data organized by the appropriate data model
  - DML also known as query language
- ◆ Two classes of languages
  - Procedural – user specifies what data is required and how to get those data
  - Nonprocedural – user specifies what data is required without specifying how to get those data
- ◆ SQL is the most widely used query language

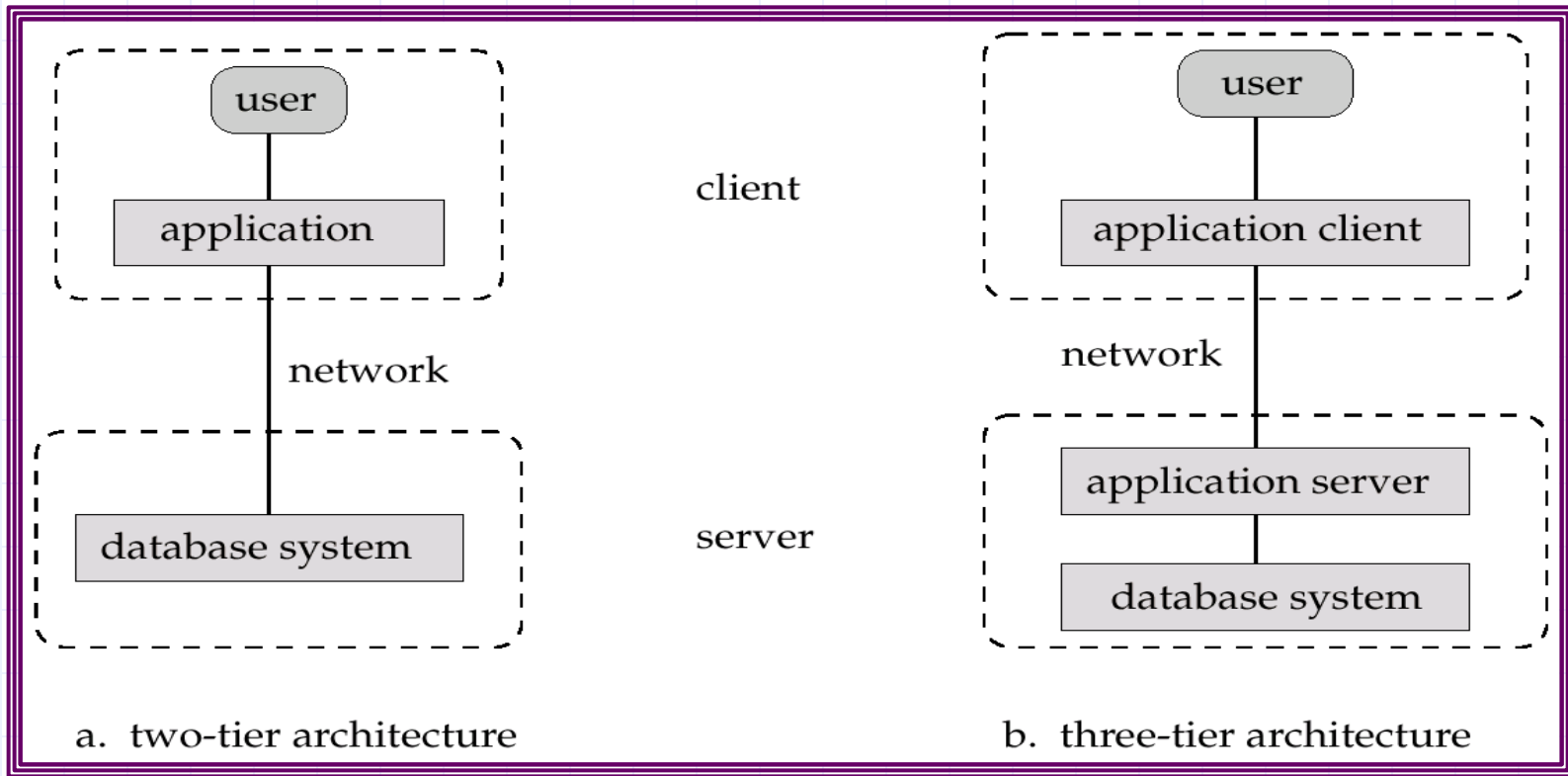


# SQL

- ◆ SQL: widely used non-procedural language
  - E.g. find the name of the customer with customer-id 192-83-7465

```
select  customer.customer-name
from    customer
where   customer.customer-id =
        '192-83-7465'
```
- ◆ Basic SQL has limited expressability
  - cannot implement any arbitrary function in SQL
- ◆ Application programs generally access databases through one of
  - Language extensions to allow embedded SQL
  - Application program interface (e.g. ODBC/JDBC) which allow SQL queries to be sent to a database

# Application Architectures



- **Two-tier architecture:** E.g. client programs using ODBC/JDBC to communicate with a database
- **Three-tier architecture:** E.g. web-based applications, and applications built using “middleware”

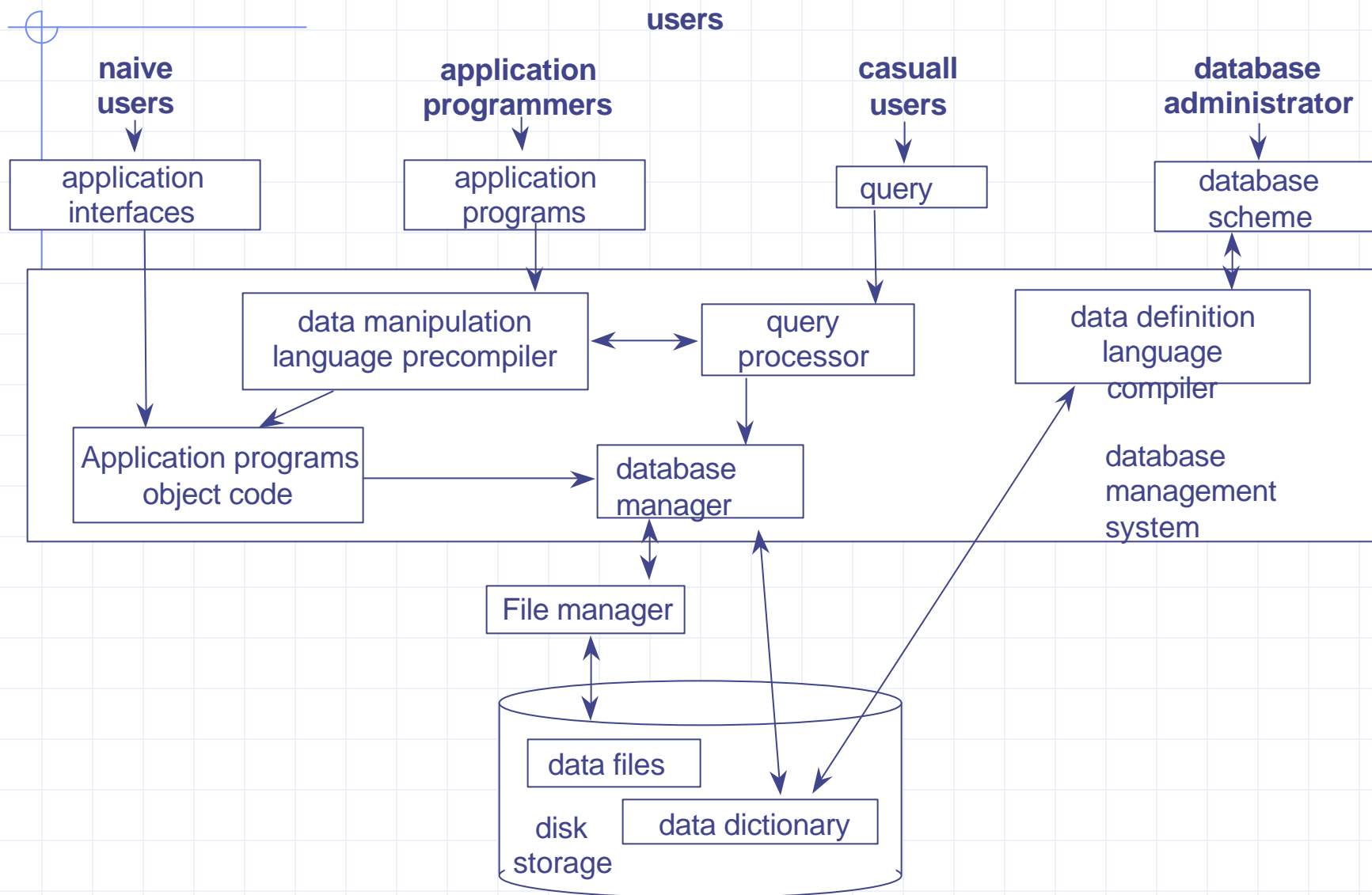
# DBMS Interface

- ◆ Provides users means to interact with database.
- ◆ Where we are today:
  - Menu driven interface, Graphical interface, Forms based interface, Natural Language Interface, WWW connectivity.
- ◆ Interface design is a tremendous challenge not only for DBMS researchers but to HCI and human cognition researchers.
- ◆ Future interfaces to databases:
  - virtual reality, immersive environments, speech, natural languages, gestures, handwriting, eye tracking brain waves, tactile interfaces, multimodal input and outputs -- combination of more than one modality.

# People Involved with DBMSs

- ◆ DBMS designers and implementers
- ◆ tool designers
- ◆ database administrator (DBA)
  - DBA = 'super-user' for a database, similar to a system administrator.
  - DBA can define schemas, views, authorization, indexes, tuning parameters, etc.
- ◆ application programmers
- ◆ database designers
  - interact with users to define database at all levels
- ◆ database and system operators.
- ◆ end users
- ◆ large number of jobs available for each of the above tasks!!

# DBMS Architecture



# Key Database Technologies

## ◆ Data Models

- allow specification of database structure at all the levels of abstraction

## ◆ Design tools

- that help in the database design process. These tools automates or facilitate some aspects of the design

## ◆ Access Methods

- data structures to support efficient access of data on disk

## ◆ Query Optimization and Processing

- efficient query processing techniques for good query performance. These techniques usually minimize the amount of disk I/O

## ◆ Transaction processing techniques

- to support concurrent access and reliability in the presence of failures

# Need for Query Optimization

- ◆ Consider two tables
  - Employee(ename, salary, department)
    - ◆ say 1000 entries
  - Manager(mname, department)
    - ◆ say 10 entries
- ◆ Query:
  - List the names of employees for the department of which Sharad is the manager

# Strategies 1

- ◆ For each entry M in Manager

- read record M

- For each entry E in employees

- ◆ read Entry E

- If (E.department == M.department) and (M.mname = "sharad")  
        print E.ename

- ◆ Cost Analysis:

- Outer loop 10 iterations. 1 read operation each time.

- Inner loop 1000 iterations. 1 entry read each time.

- total number of reads =  $10 + 1000 * 10 = 10,010$ .



## Strategy 2

- ◆ For each entry M in Manager
  - If M.mname = "sharad"  
temp = M.department
- ◆ For each entry E in Employees
  - ◆ If E.department == temp  
print E.ename
- ◆ Cost Analysis:
  - first loop 10 iterations. 1 read operation each time.
  - Second loop 1000 iterations. 1 read operation each times.
  - Total number of reads = 1010.

# Transaction Concept

- ◆ ***Atomicity:***

- all or nothing execution.

- ◆ ***Consistency:***

- execution of a transaction leaves system state as well as the state of the real world consistent.

- ◆ ***Isolation:***

- partial effects of a transaction are hidden from each other.

- ◆ ***Durability:***

- a successful transactions effects survives future system malfunctions.

# Example of Transaction

- ◆ *Withdraw \$100 checking account using an ATM.*
- ◆ Atomicity:
  - account debited if and only if t user gets money from the ATM
- ◆ Consistency:
  - balance of account is always positive.
- ◆ Isolation:
  - concurrent execution of withdraw, deposit, transfers does not result in an incorrect balance of account.
- ◆ Durability:
  - After withdraw terminates, and the ATM dispenses money account reflects that \$100 withdrawn despite failures.

# Motivation of Isolation

- ◆ Consider two transactions--
  - read account A, debit the value by \$100 and write the new value to A.
  - read account A, credit the value by \$200 and write the new value to A.
- ◆ Let initial value of A be \$1000.
- ◆ Final value should be \$1100.
- ◆ Consider the following execution if concurrency is permitted:
  - `read1(A,1000) read2(A,1000) write2(A,1200) write1(A,900)`
  - for the above execution the value of A is 900!

# Importance of the Transactions

- ◆ Transaction concept supports:
  - simple failure semantics -- either all the effects of the transaction appear or none do-- all or nothing
  - an isolated view of the world -- protection from partial effects of concurrently executing transactions
- ◆ Makes application development easy
  - complex, possibly distributed applications that share data and resources can be developed without explicitly dealing with synchronization and fault-tolerance.

# Transactions versus Other Concurrent Programming Environments

- ◆ concurrent programs prevalent in a variety of other areas in CS
  - operating systems, parallel programming, distributed systems.
- ◆ support Powerful Language Constructs:
  - to specify concurrent behavior of applications (programmers responsibility to deal with failures and concurrency issues).
- ◆ In contrast transactions relieve the application programmers of these tasks.