

# The Topology of Bendless Orthogonal Three-Dimensional Graph Drawing



**David Eppstein**  
Computer Science Dept.  
Univ. of California, Irvine

# What's in this talk?

Unexpected equivalence between a style of graph drawing  
and a type of topological embedding

3d grid drawings in which each vertex has three perpendicular edges

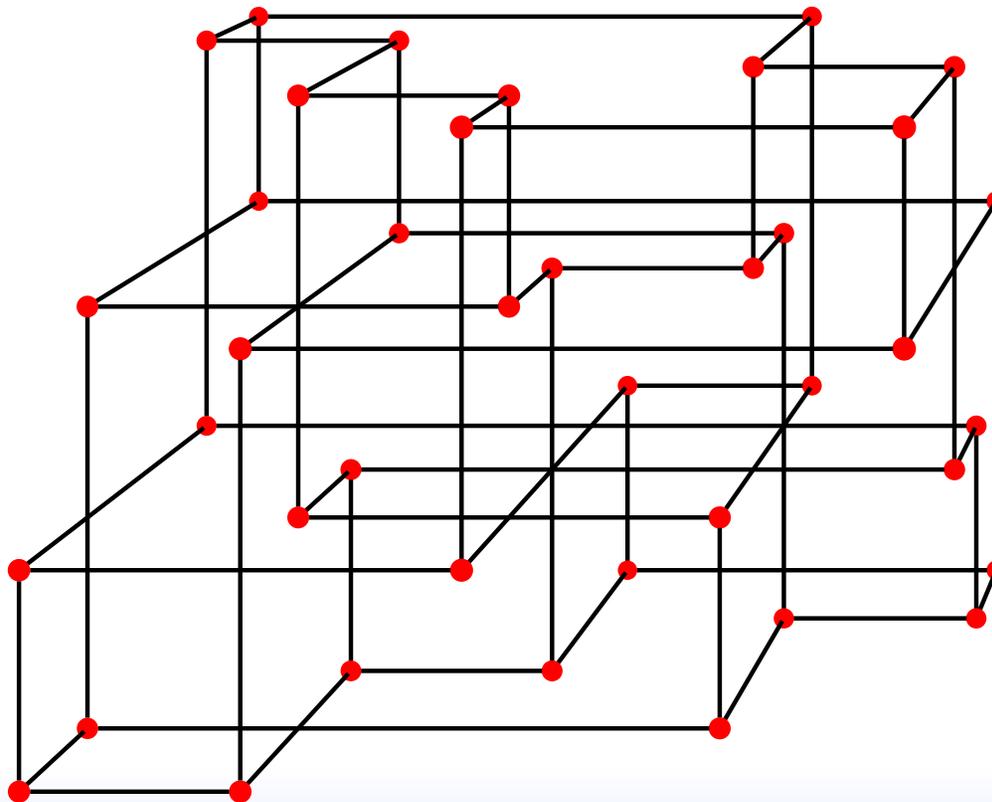
2d surface embeddings in which the faces meet nicely and may be 3-colored

...and its algorithmic consequences

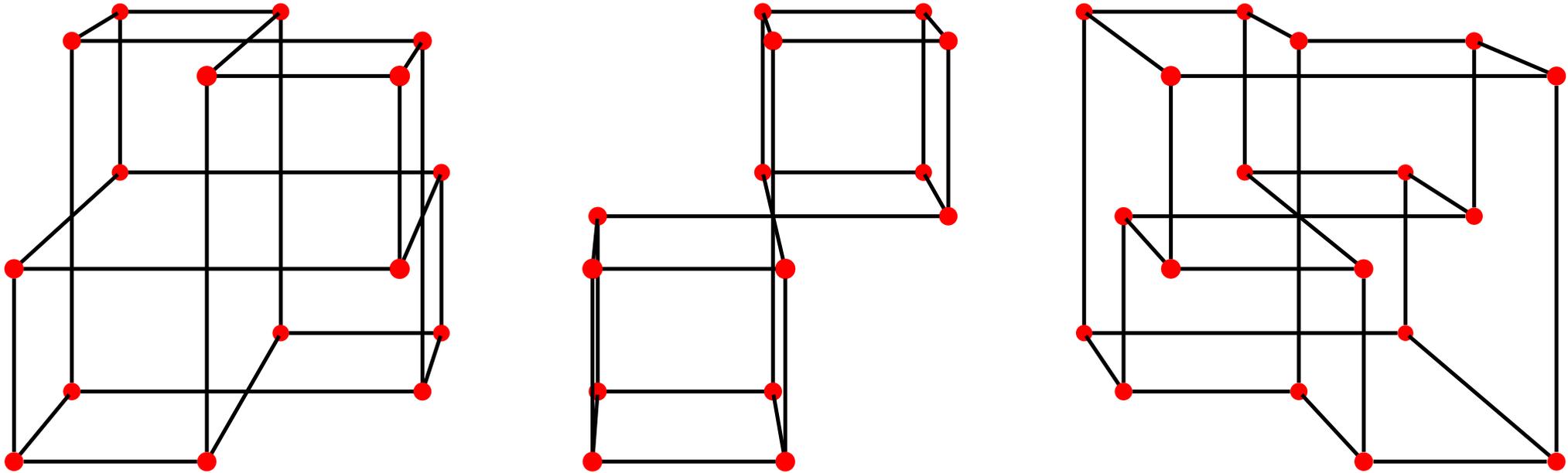
# xyz graphs

Let  $S$  be a set of points in three dimensions such that each axis-aligned line contains zero or two points of  $S$

Draw an edge between any two points on an axis-aligned line



## Three xyz graphs within a 3 x 3 x 3 grid

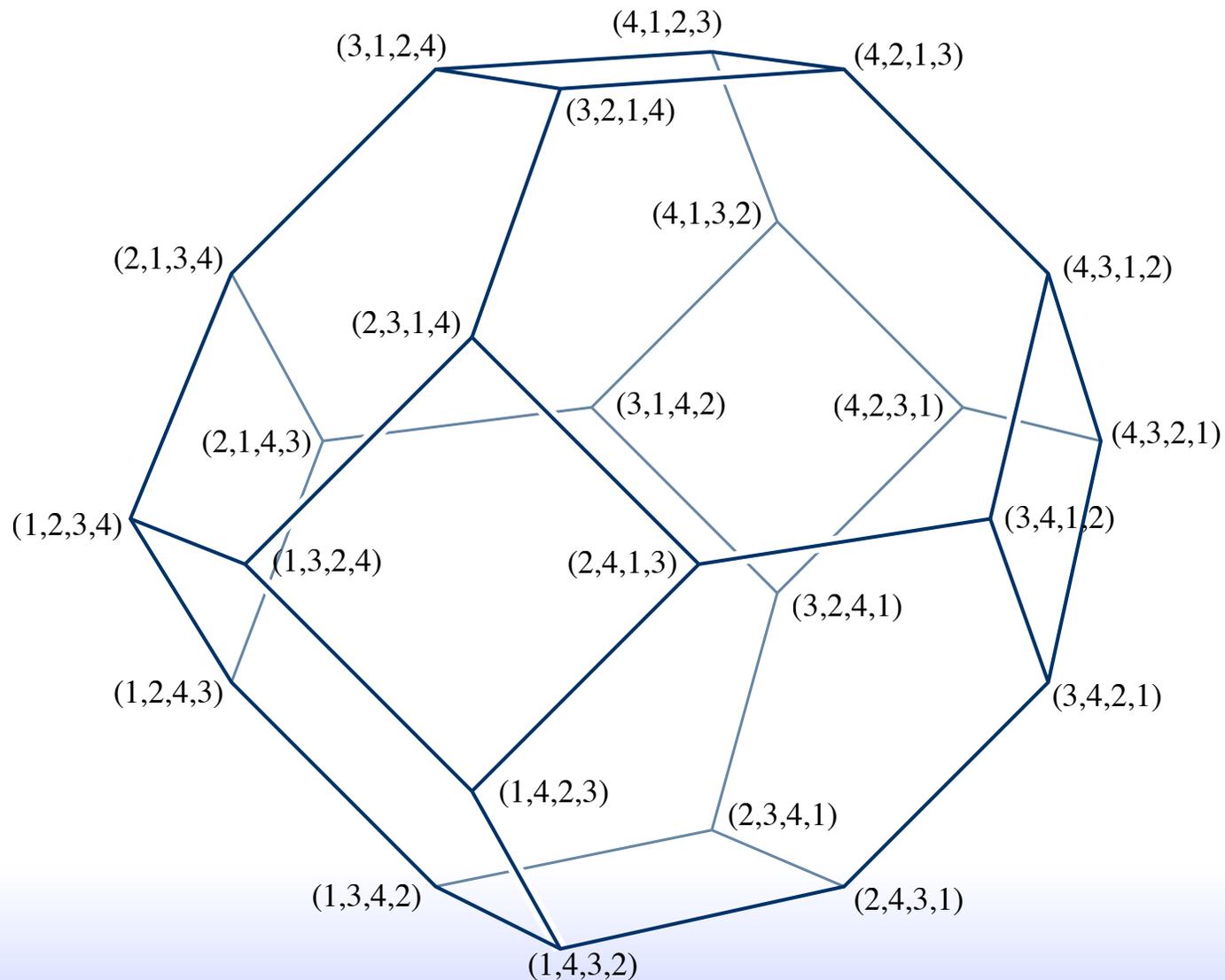


Note that edges are allowed to cross

Crossings differ visually from vertices as vertices never have two parallel edges

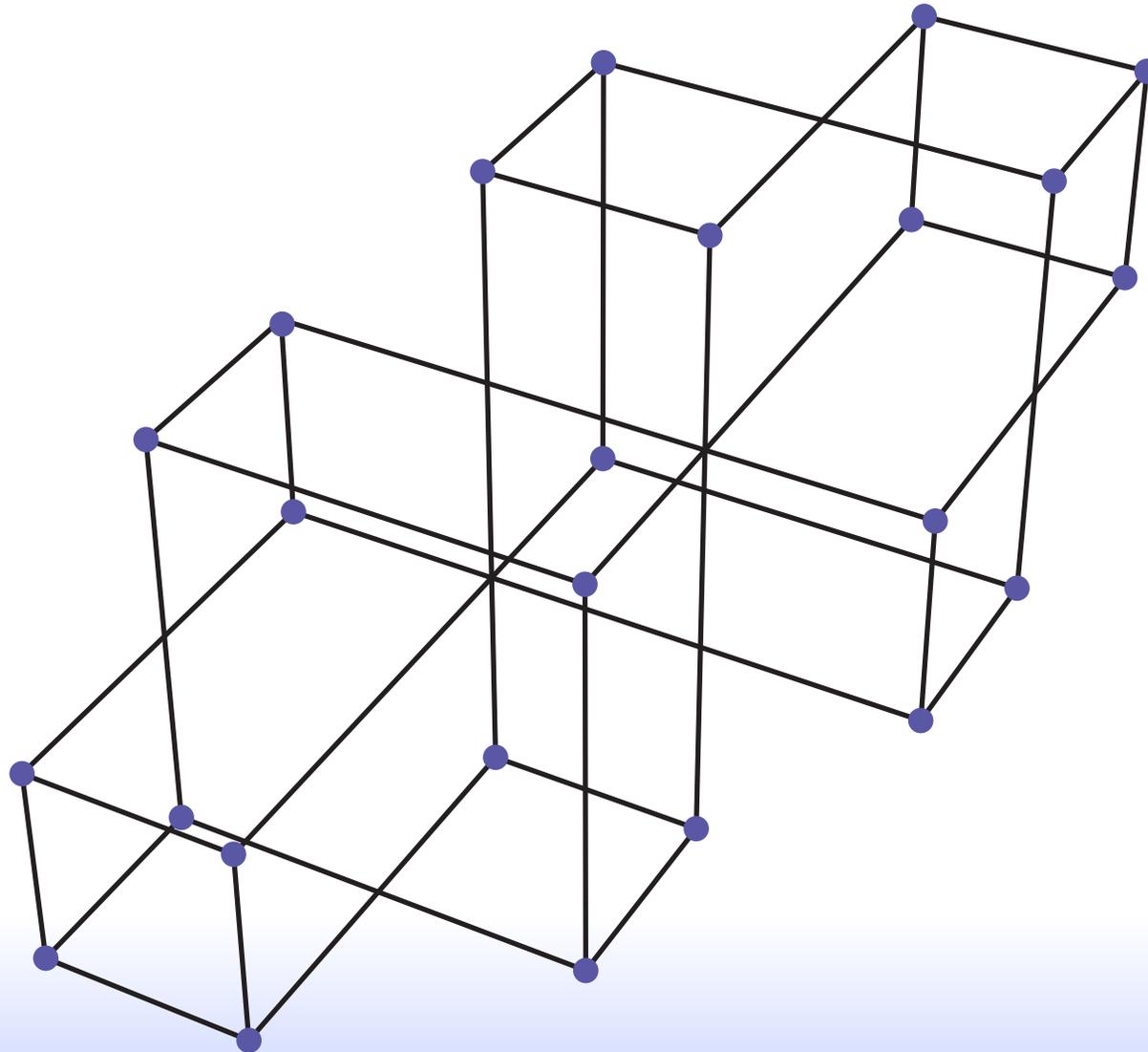
# The permutohedron

Convex hull of all permutations of  $(1,2,3,4)$  in 3-space  $x+y+z+w=10$   
Forms a truncated octahedron



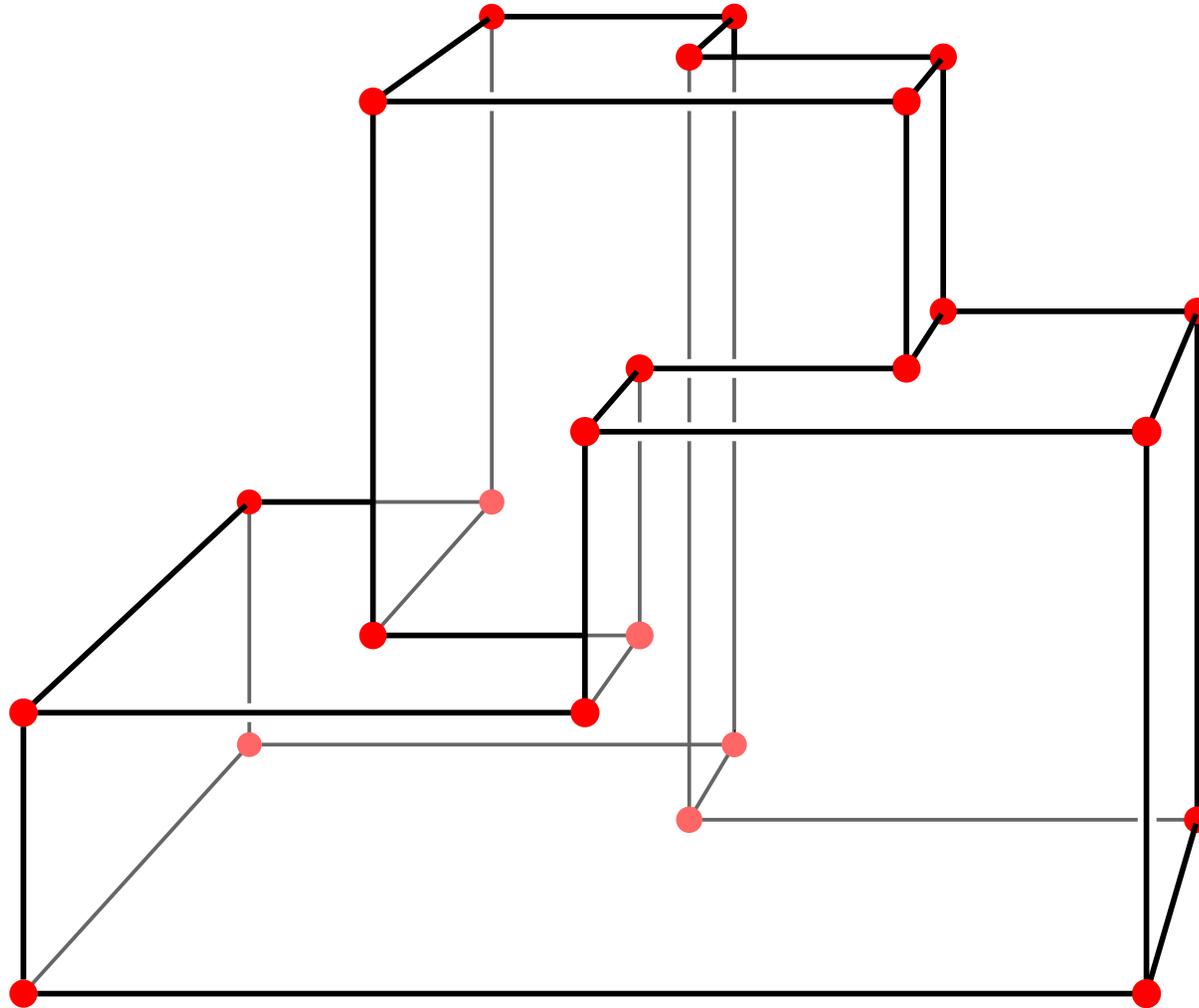
# Inverting the permutohedron

Move each permutation vertex to its inverse permutation affine transform so that the edges are axis-aligned



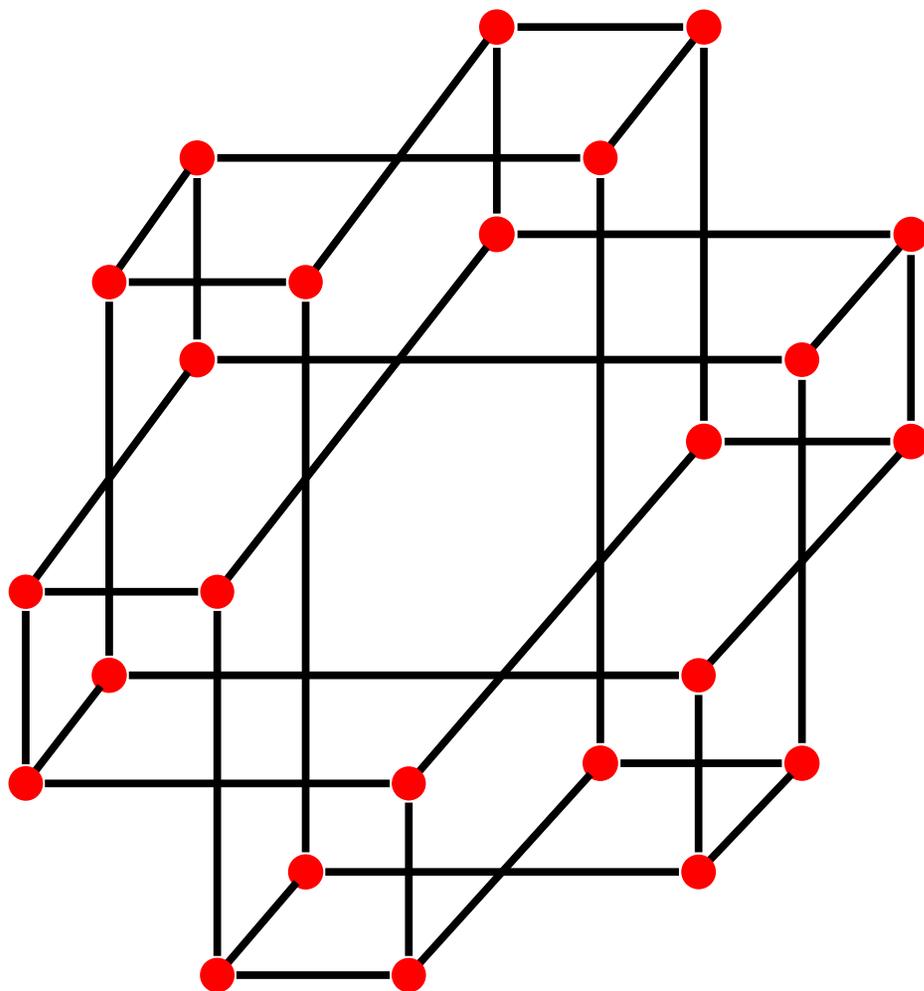
# A polyhedron for the inverse permutohedron

Rearrange face planes to form nonconvex topological sphere



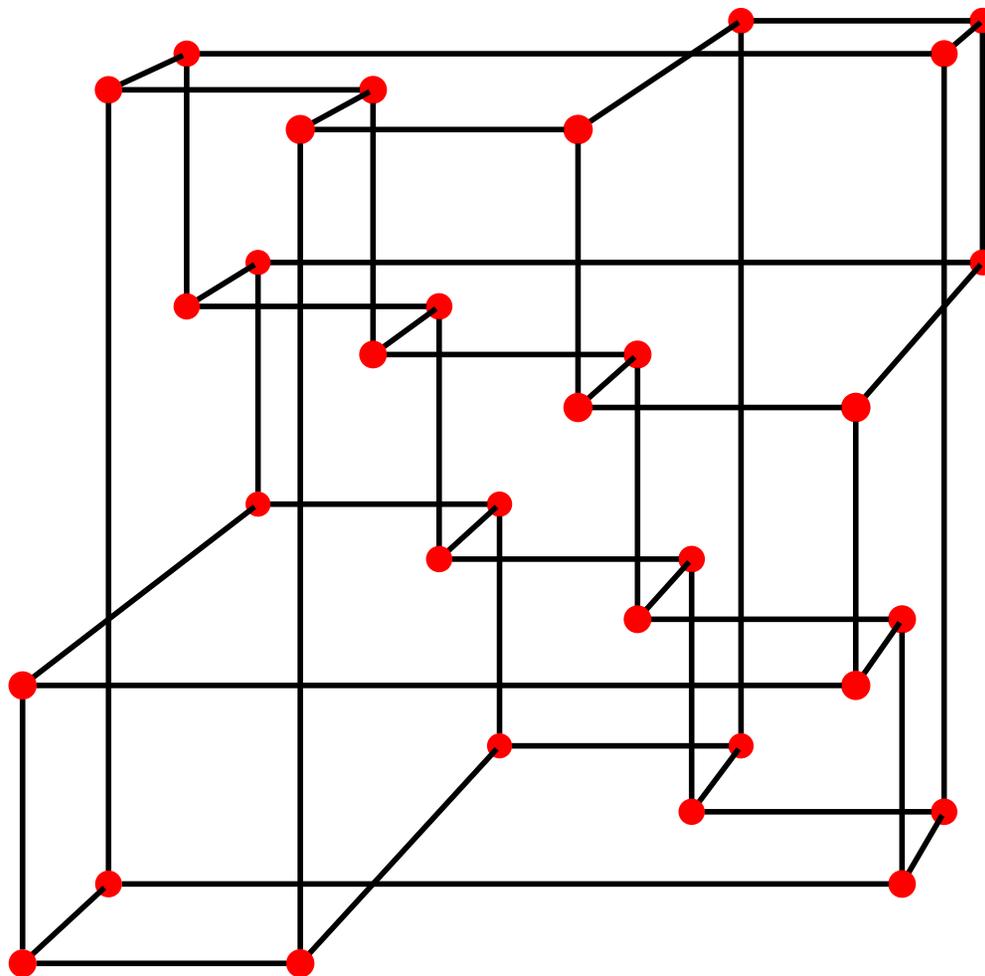
# A different xyz graph on 4-element permutations

Project  $(x,y,z,w)$  to  $(x,y,z)$



# xyz graphs with many vertices in a small bounding box

In  $n \times n \times n$  box, place points such that  $x+y+z = 0$  or  $1 \pmod n$



$n = 4$ , the Dyck graph

# Basic properties of xyz graphs

3-regular (each vertex has exactly three edges)

Triangle-free  
and 5-cycle-free  
(but may have longer odd cycles)

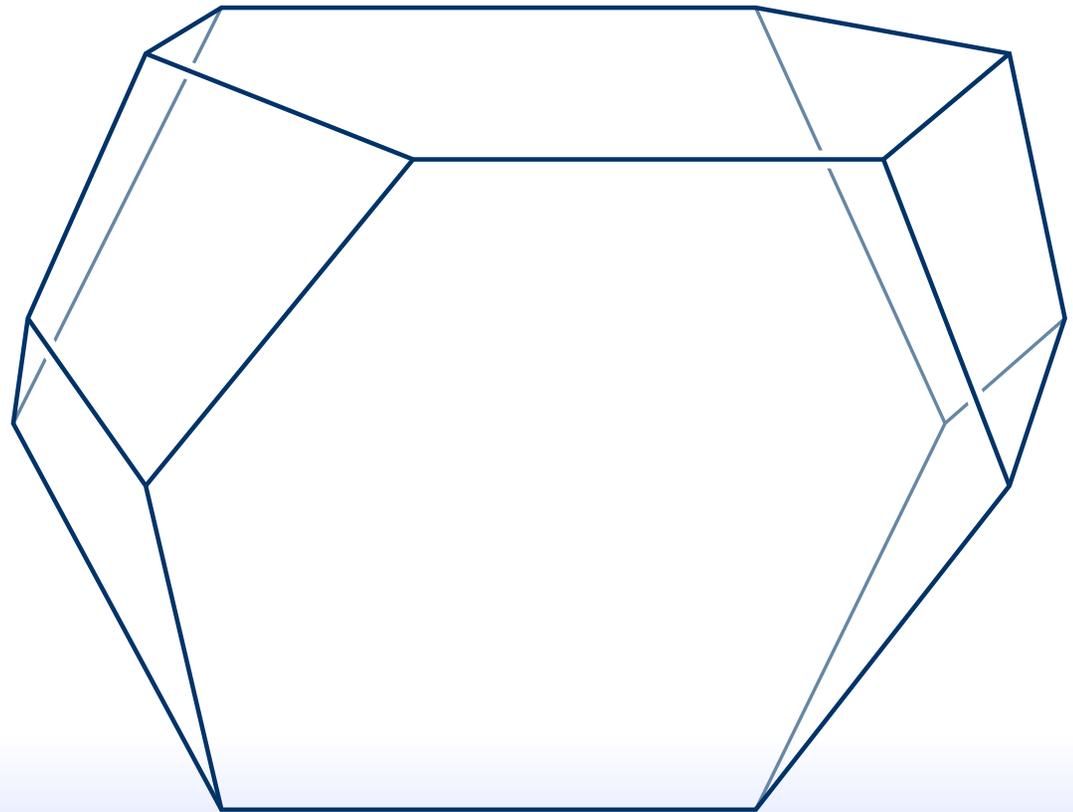
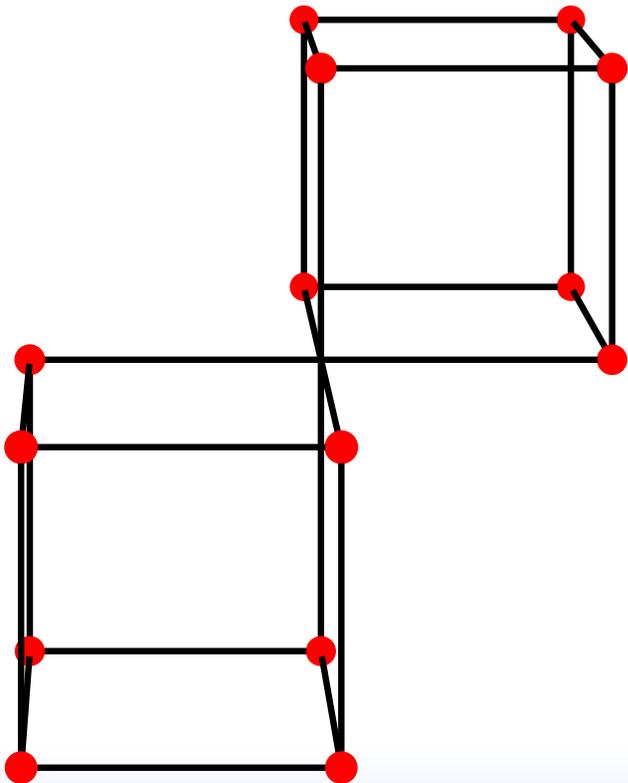
3-connected  
(can replace any edge by paths of alternating parallel and perpendicular edges,  
with two different choices of perpendicular direction)

Are these (or similar simple properties) sufficient to characterize them?

# From xyz graphs to surface embeddings

Edges parallel to any coordinate plane  
form degree-two subgraph (collection of cycles)

Form a face of a surface for each cycle



# Basic properties of xyz surfaces

All faces are topological disks (by construction)

If two faces meet, they lie on perpendicular planes  
the planes meet in a line  
and the faces meet in an edge lying on that line

The faces may be given three colors  
(by the direction of the planes they lie in)  
and are thus properly 3-colored

# From xyz surfaces to xyz graphs

Let  $G$  be a 3-regular graph embedded on a surface, so that

- faces are topological disks
- any two intersecting faces meet in a single edge
- the faces are properly 3-colored  
(say, red, blue, and green)

Number the faces of each color

Assign coordinates of a vertex:

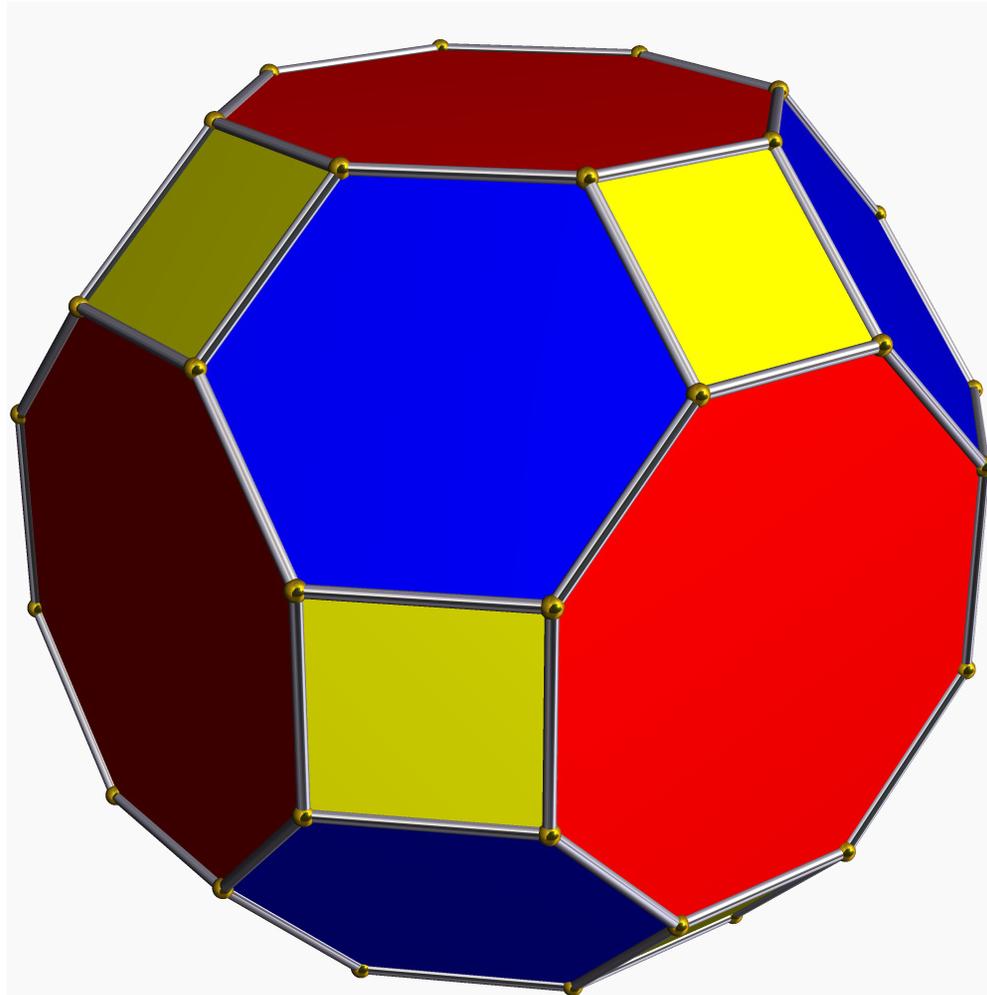
$x$  = the number of its red face

$y$  = the number of its blue face

$z$  = the number of its green face

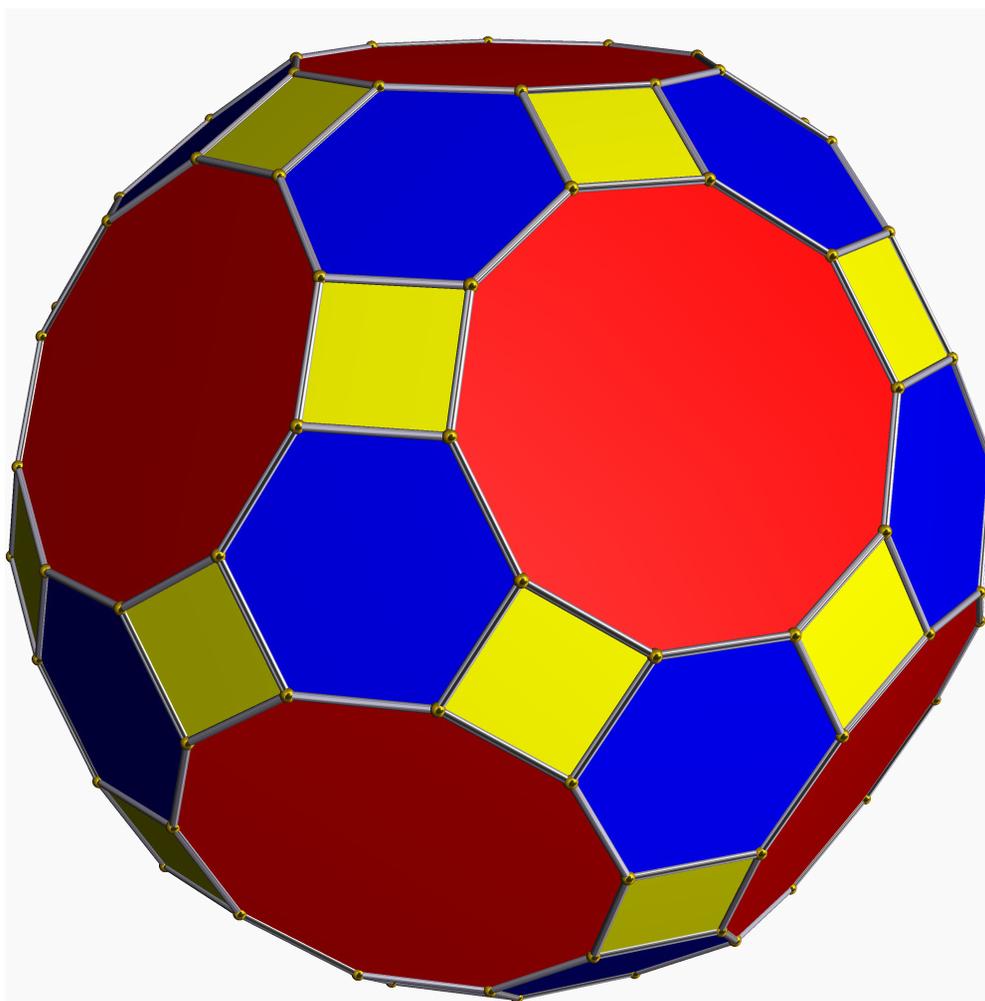
The result is an xyz graph embedding!

# Great rhombicuboctahedron



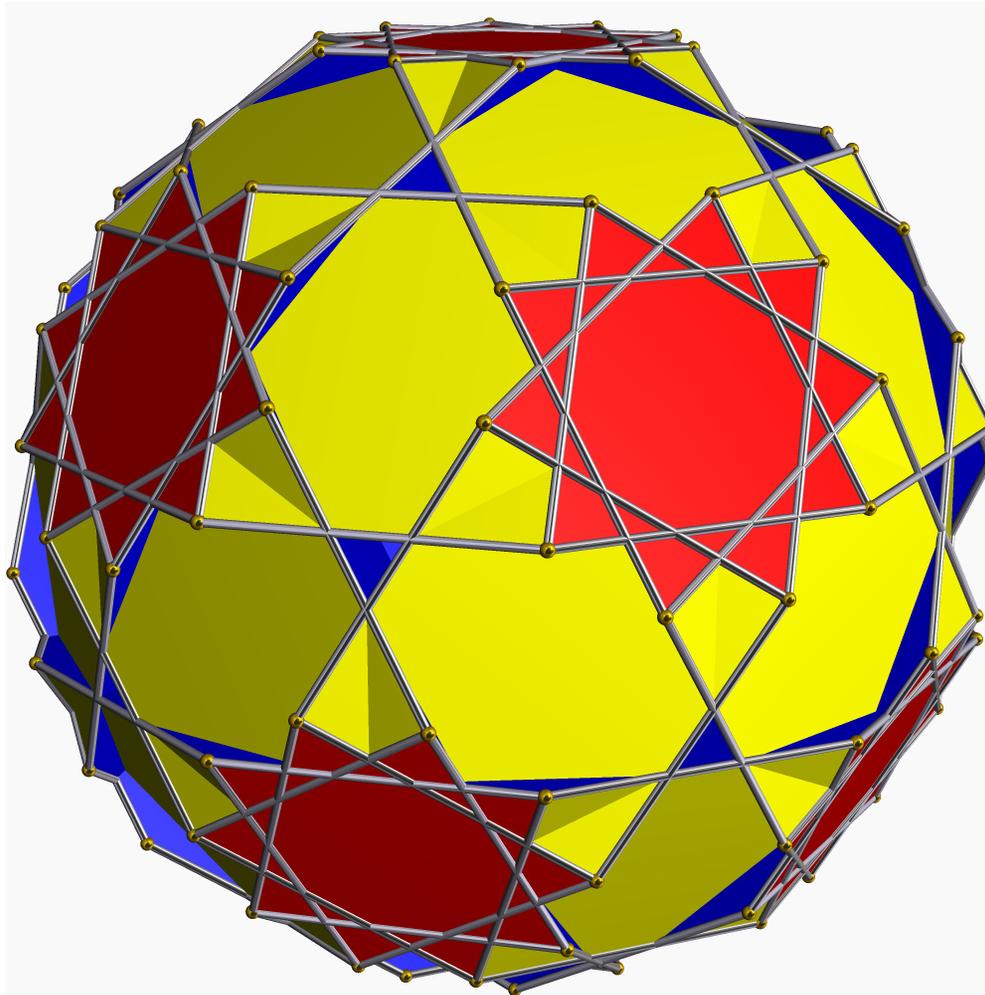
By Robert Webb using Great Stella, <http://www.software3d.com/Stella.html>  
image from [http://commons.wikimedia.org/wiki/Image:Great\\_rhombicuboctahedron.png](http://commons.wikimedia.org/wiki/Image:Great_rhombicuboctahedron.png)

# Great rhombicosidodecahedron



By Robert Webb using Great Stella, <http://www.software3d.com/Stella.html>  
image from [http://commons.wikimedia.org/wiki/Image:Great\\_rhombicosidodecahedron.png](http://commons.wikimedia.org/wiki/Image:Great_rhombicosidodecahedron.png)

# Truncated dodecadodecahedron

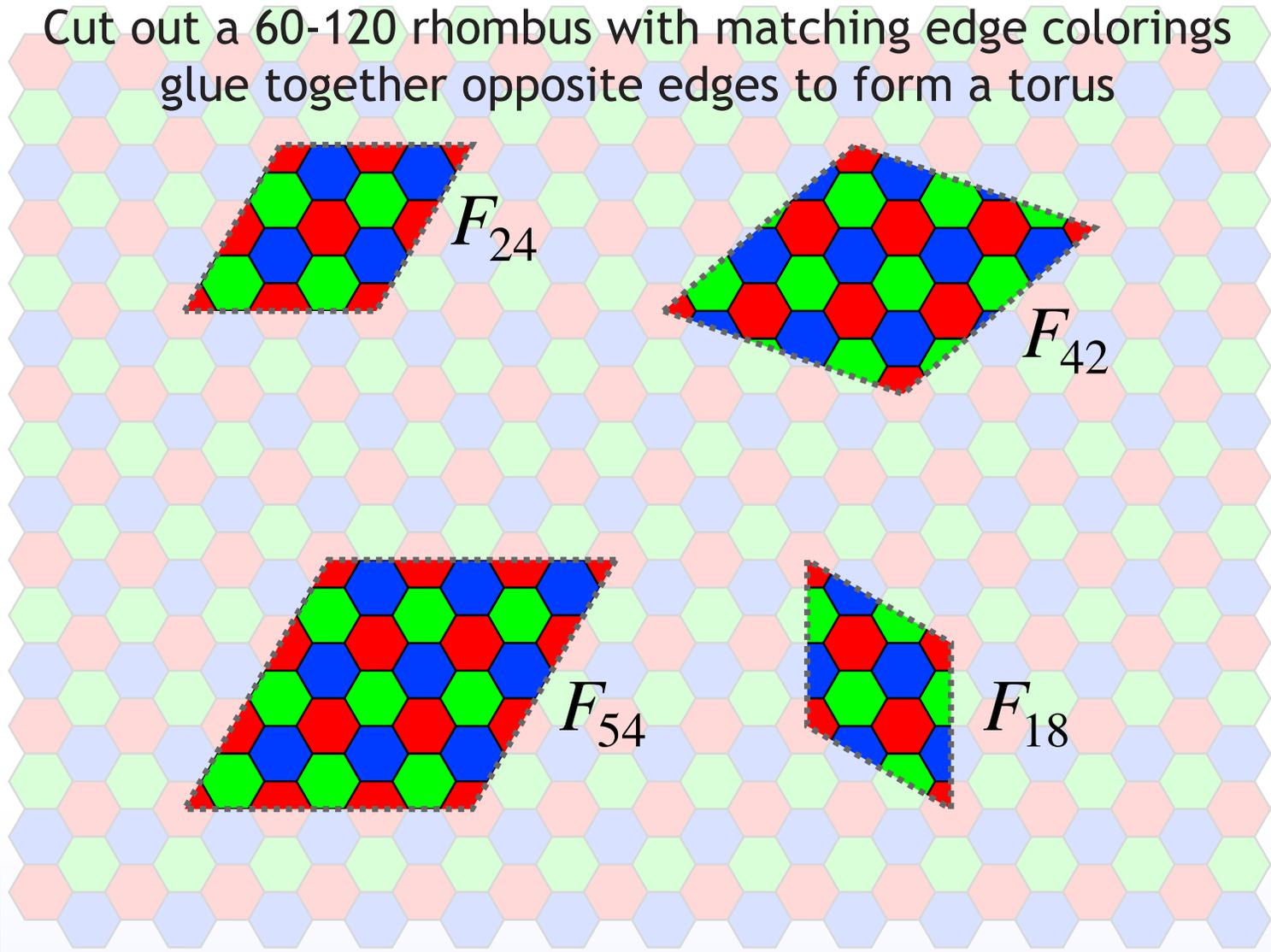


By Robert Webb using Great Stella, <http://www.software3d.com/Stella.html>  
image from [http://commons.wikimedia.org/wiki/Image:Truncated\\_dodecadodecahedron.png](http://commons.wikimedia.org/wiki/Image:Truncated_dodecadodecahedron.png)

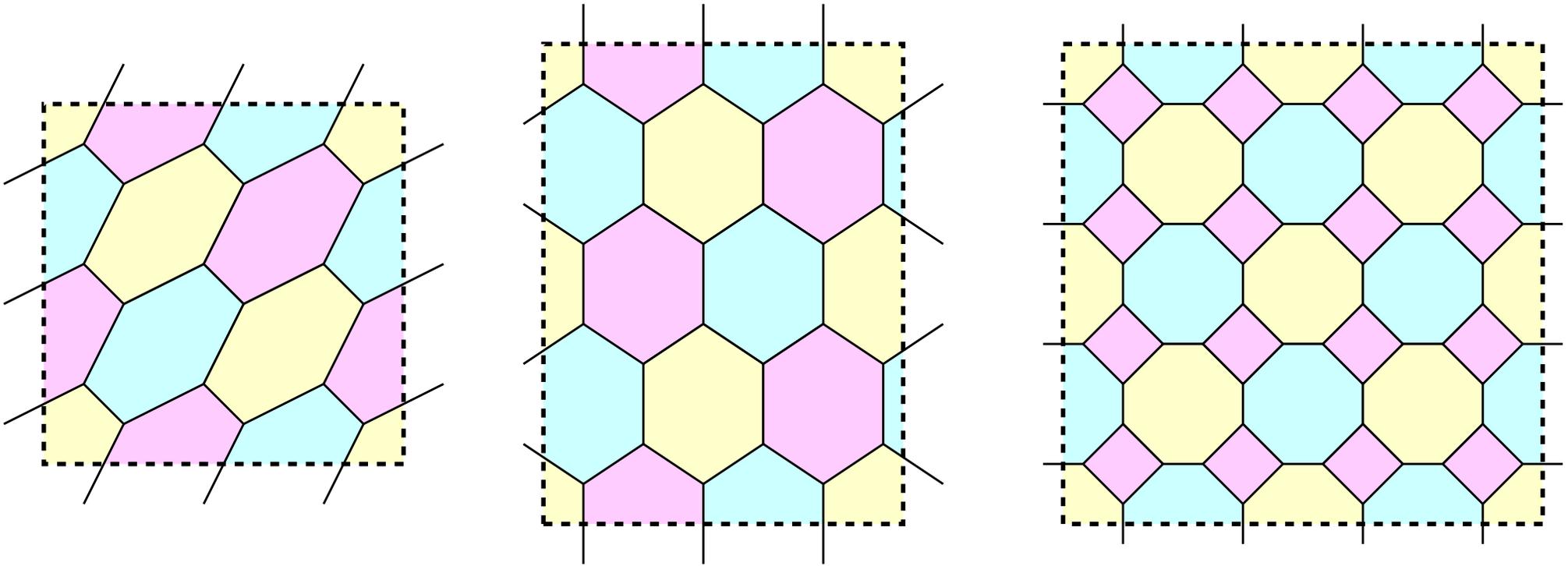
# Symmetric graphs on the torus

Start with regular tiling of plane by 3-colored hexagons

Cut out a 60-120 rhombus with matching edge colorings  
glue together opposite edges to form a torus



## More xyz tori



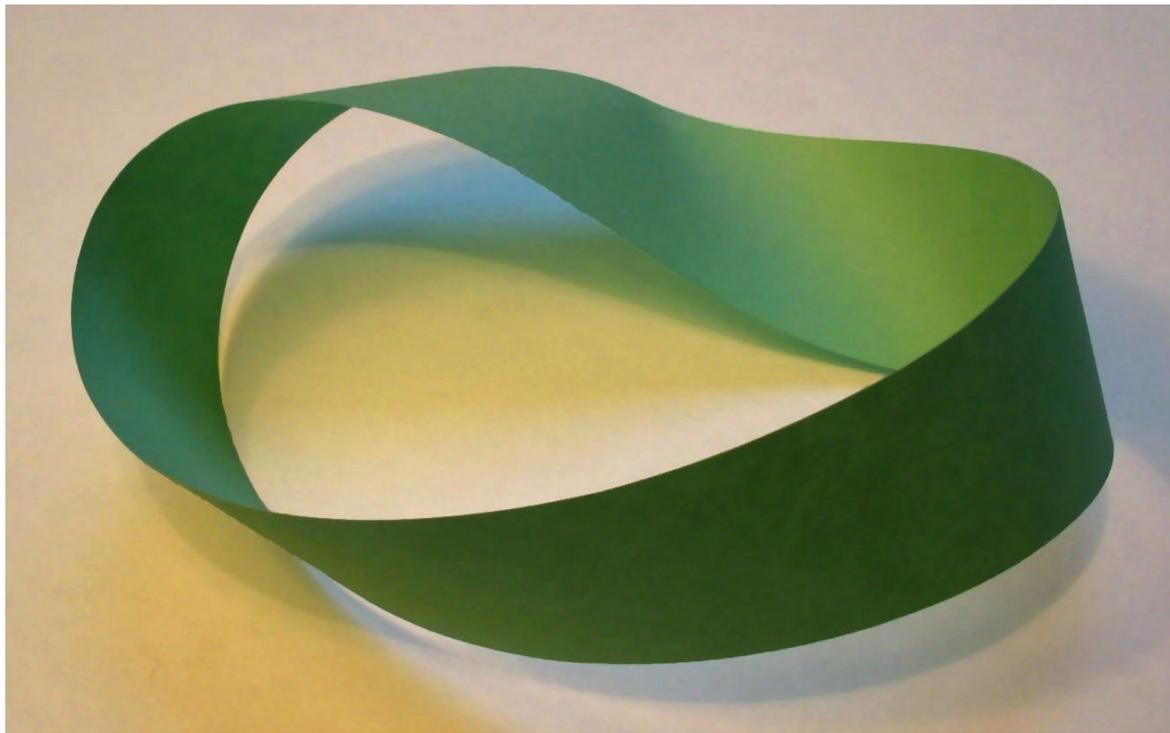
Leftmost example is order-4 cube-connected cycles network  
Embedding generalizes to any even order CCC

# Bipartiteness and orientability

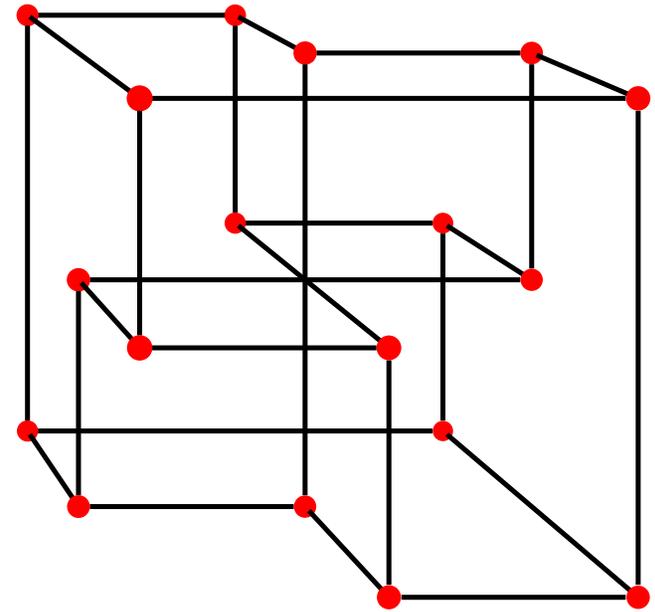
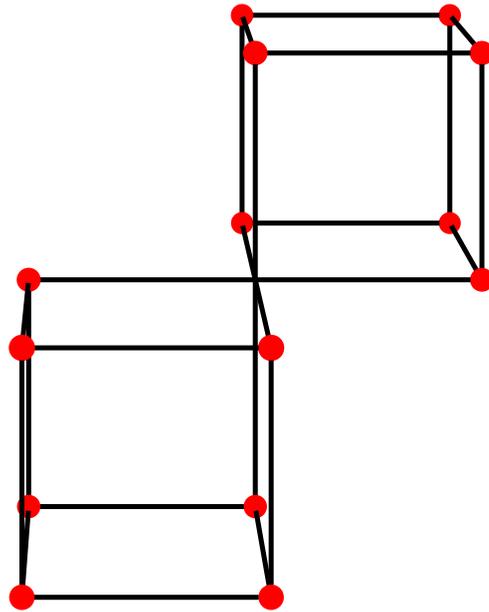
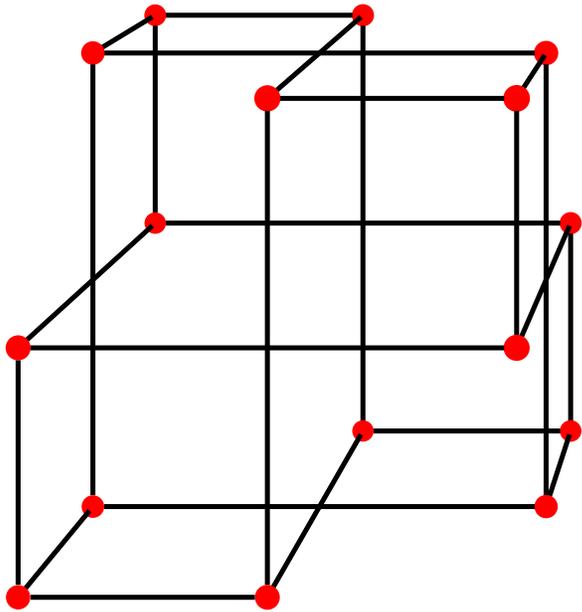
Theorem: Let  $G$  be an xyz graph.  
Then  $G$  is bipartite if and only if the corresponding  
xyz surface is orientable

Orientable surfaces: sphere, torus, ...

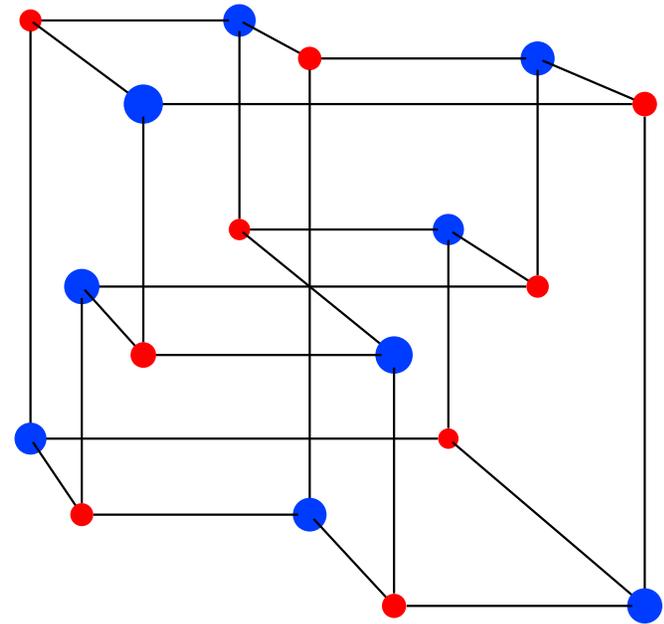
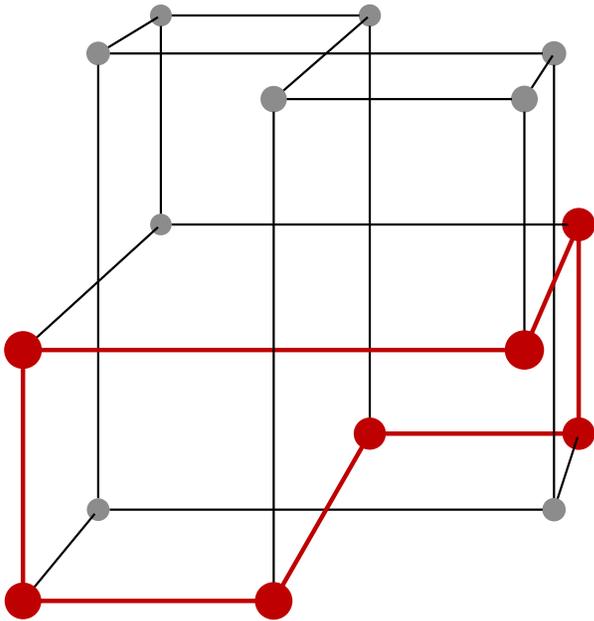
Non-orientable surfaces: Möbius strip, projective plane, Klein bottle, ...



**Puzzle: which of these three graphs is not bipartite?**



# Puzzle solution



# Planar xyz graphs

Lemma: If  $G$  is a planar xyz graph, its xyz surface must be a topological sphere

Therefore, every planar xyz graph is 3-connected and bipartite

Known:

every 3-connected planar graph is the skeleton of a polyhedron  
(so faces meet at most in single edges)

every bipartite polyhedron has 3-colorable faces

Therefore: a planar graph has an xyz embedding  
if and only if it's 3-connected and bipartite

## Testing if a surface embedding is xyz

Choose arbitrarily two colors for two adjacent faces

Propagate colors:

If some face has neighbors of two colors, assign it the third color

Must successfully color all faces of any xyz surface  
(colors are forced by triples of faces along a path connecting any two faces)

**So embedding is xyz iff faces intersect properly and coloring succeeds**

## Testing if a partition of edges into parallel classes is xyz

Find the xyz surface embedding that would correspond to the partition

Check that faces intersect properly and color it

# Testing if a graph has an xyz embedding

Try all partitions of its edges into three matchings

Backtracking algorithm:

order vertices so all but two have both incoming and outgoing edges

assign edges of first vertex to matchings, arbitrarily

for each remaining vertex, in order:  
try all assignments of its incident edges to matchings  
that are consistent with previous choices

Vertex with two incoming edges has only one choice

Vertex with two outgoing edges has two choices

So number of search paths  $\leq 2^{n/2-1}$  and total time =  $O(2^{n/2})$

# Implementation

123 lines of Python

<http://www.ics.uci.edu/~eppstein/PADS/xyzGraph.py>

Successfully run on graphs on **up to 54 vertices**

Could probably benefit from additional optimization:

- Faster test for each edge partition
- Early backtrack for bad partial partitions



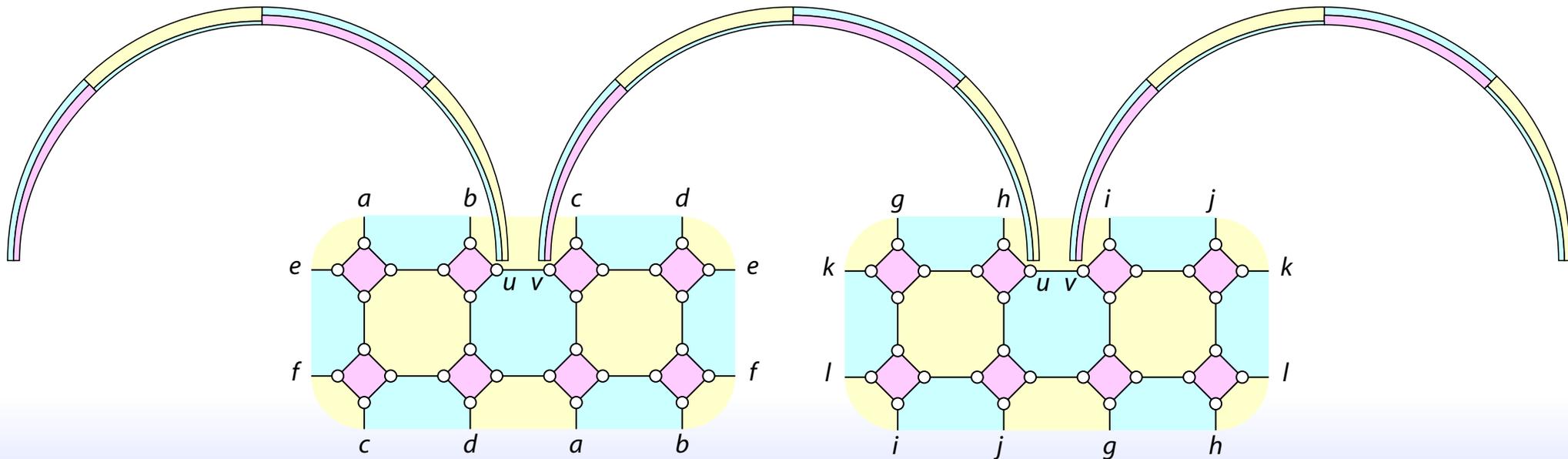
# Recognizing xyz graphs is NP-complete

Proof idea: reduction from graph 3-coloring

Represent color as orientation of an edge of a connector gadget

Vertex of graph to be colored becomes planar graph in possible xyz graph

Edge in graph to be colored becomes edge gadget formed from three connectors and two ambiguous tori



# Conclusions

Interesting type of 3d graph drawing

Equivalence with 2d surface embedding leads to some deep theory

It's NP-complete

but...

that doesn't prevent us from implementing algorithms and finding drawings