

Proof of correctness for Power(a, n).

Power(a, n) $\xrightarrow{\text{red arrow}}$ a^n
 a real #
 n non-neg int.

• \Rightarrow if $(n=0)$ return 1.

\Rightarrow $\left\{ \begin{array}{l} \text{if } (n \text{ is even}) \\ \text{return } (\text{Power}(a, n/2))^2 \\ \text{if } (n \text{ is odd}) \\ \text{return } (\text{Power}(a, \lfloor n/2 \rfloor))^2 \cdot a \end{array} \right.$

End.

Theorem For any real # a , and any non-neg int n , Power(a, n) returns a^n

$\forall n, \forall a \text{ Power}(a, n) = a^n$
 $\mathcal{Q}(n).$

Proof By induction on n .

Base: $n=0$ $a^0 = 1$
 Power(a, 0) returns 1

Induction step: Assume for $k = 0, 1, \dots, n-1$, and $n \geq 1$.
Power(a, k) returns a^k .
 Prove Power(a, n) returns a^n .

Case 1 n is even and $n \geq 1$.
 $n = 2k$ for some int k . $n \geq 2$
 $k \geq 1$.

Recursive call is $\text{Power}(a, n/2)$
 $= \text{Power}(a, k)$. \rightarrow

Show $k \in \{0, 1, \dots, n-1\}$.
 $0 \leq k \leq n-1$.

Show $k \leq n-1$. $n \geq 2$.
 $\frac{n}{2} \leq n-1$.

By the I.H.

$$\text{Power}(a, n/2) = a^{n/2}$$

$$\begin{aligned} \text{Power}(a, n) &\text{ returns } [\text{Power}(a, n/2)]^2 \\ &= (a^{n/2})^2 = a^{n/2 \cdot 2} = a^n \end{aligned}$$

$$\begin{aligned} 2 &\leq n && +n \\ 2+n &\leq n+n \\ n &\leq 2n-2 \\ n &\leq 2(n-1) \\ n/2 &\leq n-1 \end{aligned}$$

Case 2: n is odd and $n \geq 1$

$$2k+1 = n \quad k \geq 0.$$

$$\begin{aligned} \text{Recursive call: } \lfloor n/2 \rfloor &= \lfloor \frac{2k+1}{2} \rfloor \\ &= \lfloor k + 1/2 \rfloor = k. \end{aligned}$$

$$\textcircled{n = 2k + 1} \quad \underline{k \geq 0} \quad n \geq 1.$$

Prove all $\text{Power}(a, \lfloor n/2 \rfloor) = \text{Power}(a, k)$

$$\underline{0 \leq k \leq n-1}$$

$$\begin{aligned} n &= 2k + 1 \\ \underline{n-1} &= \underline{2k} \geq k \end{aligned}$$

By the I.H. $\text{Power}(a, \lfloor n/2 \rfloor) = \underline{\underline{\text{Power}(a, k) = a^k}}$

$$\text{Power}(a, n) \text{ returns } [\text{Power}(a, k)]^2 \cdot a$$

$$\geq (a^k)^2 \cdot a =$$

$$a^{2k} \cdot a = a^{2k+1} = a^n \quad \square$$

Recursive def of nested parens.

Base: $()$ properly nested

Recursive rule: If x is properly nested then
 (x) is P.N.

If x & y are P.N.
then xy is P.N.

Thm If s is properly nested then # of
left parens in s is equal to the # of right parens in s .

If s is a sequence of parens,
 $N[(, s)] = \#$ of left parens in s .

$$N[), s] = \# \text{ right parens.}$$

Thm $\forall n \geq 2$ If s is properly nested seq of parens $\text{w/ } n \text{ symbols}$
then $N[(, s] = N[), s]$.

Proof By ind on the length of the seq.

Base case: $n=2$ $s = ()$
 $N[(, s] = N[), s] = 1$.

Assume for $k=2, \dots, n-1$.

any properly nested seq of length k has the same # of left parens + right parens.
Prove claim holds true for length n sequences.

If s is properly nested + length n .

- ① $s = (x)$ where x is properly nested.
- ② $s = xy$ where $x + y$ are properly nested.

Case 1 $s = (x)$ By I.H. $N[(, x] = \underline{N[), x]}$

$$N[(, s] = 1 + N[(, x] = \underline{1 + N[), x]} = N[), s].$$

Case 2 $s = xy$. By I.H. $N[(, x] = N[), x]$
 $N[(, y] = N[), y]$

$$N[(, s] = N[(, x] + N[(, y] = \frac{N[), x] + N[), y]}{N[), s]} \quad \square$$

Number theory: $a \mid b$.

$a \mid b$ a "divides" b .

b is a multiple of a .
 a is a factor/divisor of b .
 $b = k \cdot a$ $k \text{ int.}$

$$6 \mid 48 \quad 48 = 6 \cdot 8.$$

$$\begin{array}{l} -3 \mid 90 \\ 5 \mid -45 \end{array}$$

$$90 = -3(-30)$$

$$6 \mid 49$$

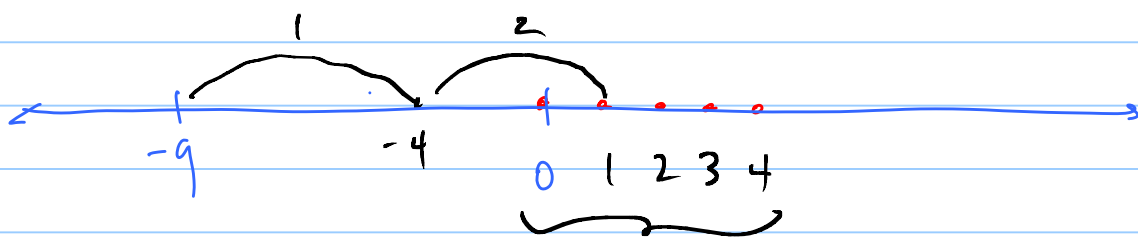
Modular arithmetic

$$17 \bmod 3 = 2.$$

$$-17 \bmod 3 = 1$$

For any int n and any $d \geq 1$.
There are unique integers q, r

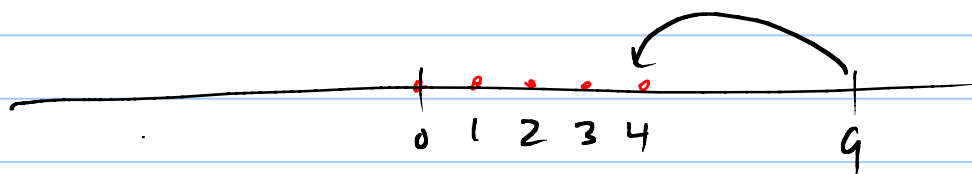
$$n = \underline{q} \cdot d + r \quad r \in \{0, \dots, d-1\}.$$



$$\begin{array}{l} -9 \bmod 5 = 1 \\ -9 \text{ div } 5 \end{array}$$

$$-9 = \underbrace{-2}_q \cdot 5 + 1.$$

$$\begin{aligned} 9 \bmod 5 &= 4 \\ \uparrow \text{div } 5 &= 1 \end{aligned}$$



$$n \text{ div } d = \lfloor n/d \rfloor \quad \lfloor -9/5 \rfloor = -2$$

$$n \bmod d = n - (n \text{ div } d) \cdot d.$$

$$-9 - (-2) \cdot 5 = 1.$$