Discussion Section Week 1

• Intro

• Course Project Information

- Constraint Satisfaction Problems
 - Sudoku
 - Backtracking Search Example
 - Heuristics for guiding Search Example

Intro

- Teaching Assistant
 - Junkyu Lee (June Queue Lee)
 - Office Hour
 - Friday 11:00 AM ~ 12:00 PM
 - DBH 4099
- Reader
 - Minhaeng Lee (Min Heng Lee)
 - Office Hour
 - Thursday 2:00 PM ~ 3:00 PM
 - DBH 4219

Course Project Information

- Fri., 15 Jan., 11:59pm: Project Team Formation
- Sun., 24 Jan., 11:59pm: Project **Problem** Generator
- Sun., 31 Jan., 11:59pm: Project Backtracking Search
- Sun., 14 Feb., 11:59pm: Project Forward Checking
- Sun., 21 Feb., 11:59pm: Project Arc Consistency
- Sun., 28 Feb., 11:59pm: Project MRV & DH Heuristic
- Sun., 6 Mar.., 11:59pm: Project LCV Heuristic
- Sun., 13 Mar., 11:59pm: Final Project

You will lose 10% of your Project grade for every day or fraction thereof it is late

Course Project Information

- Fri., 15 Jan., 11:59pm: Project Team Formation
 - How Many Members ?
 - We will post a google doc next week on EEE message board

You Will Be Expected to Know

- Basic definitions (section 6.1)
 What is a CSP?
- Backtracking search for CSPs (6.3)
- Variable ordering or selection (6.3.1)
 - Minimum Remaining Values (MRV) heuristic
 - Degree Heuristic (DH) (to <u>unassigned</u> variables)
- Value ordering or selection (6.3.1)
 Least constraining value (LCV) heuristic

What is CSP?

• Task

• Model

What is CSP?

• Task/goal for solving CSP

- Given a set of constraints,

- Find a solution that satisfy all constraints
- Find all solutions that satisfy all constraints
- Count number of solutions

• .

What is CSP?

• How to model/express CSP problems?

- variable and its domain

- constraints, relations, functions

• allowed (partial) combinations of variable values

Constraint Satisfaction Problems

- What is a CSP?
 - Finite set of variables X_1, X_2, \dots, X_n
 - Nonempty domain of possible values for each variable $D_1, D_2, ..., D_n$
 - Finite set of constraints C_1, C_2, \ldots, C_m
 - Each constraint C_i limits the values that variables can take,
 - e.g., $X_1 \neq X_2$
 - Each constraint C_i is a pair <scope, relation>
 - Scope = Tuple of variables that participate in the constraint.
 - Relation = List of allowed combinations of variable values. May be an explicit list of allowed combinations. May be an abstract relation allowing membership testing and listing.

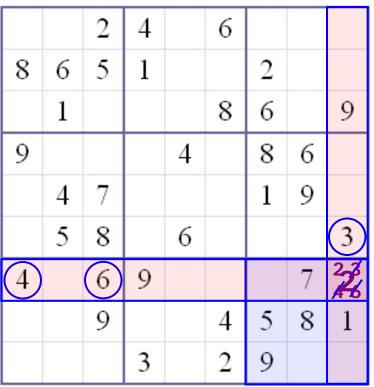
• CSP benefits

- Standard representation pattern
- Generic goal and successor functions
- Generic heuristics (no domain specific expertise).

Sudoku

Example: Sudoku (constraint propagation)

Constraint propagation



•Variables: 81 slots

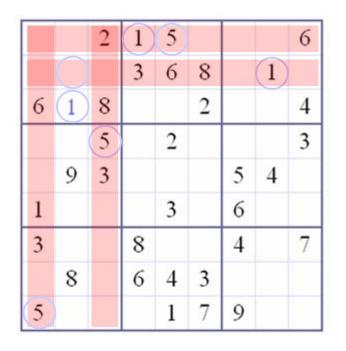
•Domains = {1,2,3,4,5,6,7,8,9}

•Constraints: •27 not-equal

Each row, column and major block must be all different

"Well posed" if it has unique solution: 27 constraints CS 275 Winter 2016, Constraint Networks, Rina Dechter

Sudoku (inference)



Each row, column and major block must be all different

"Well posed" if it has unique solution

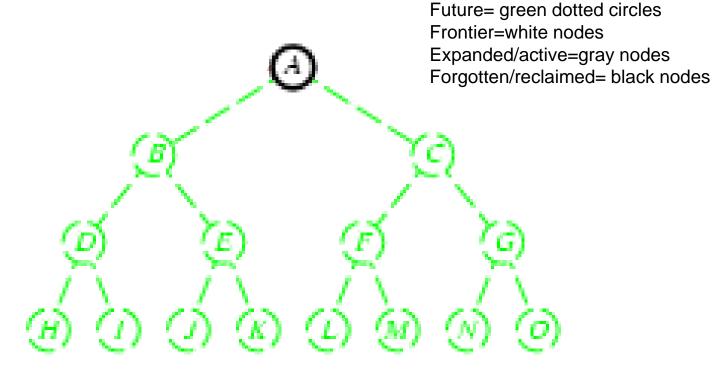
CS 275 Winter 2016, Constraint Networks, Rina Dechter

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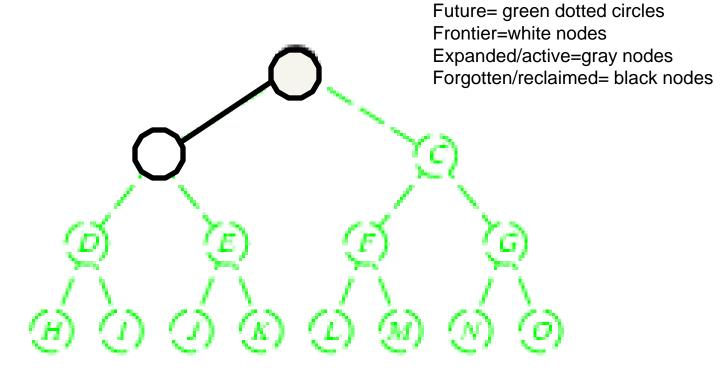
Backtracking search

- Similar to Depth-first search
 - At each level, picks a single variable to explore
 - Iterates over the domain values of that variable
- Generates kids one at a time, one per value
- Backtracks when a variable has no legal values left
- Uninformed algorithm
 - No good general performance

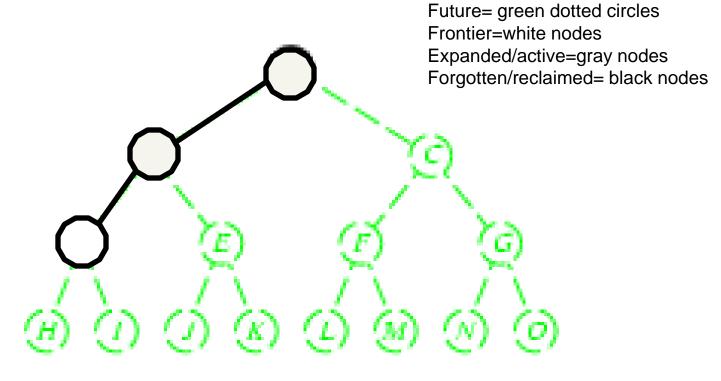
- Generate only one child at a time.
- Goal-Test when inserted.
 - For CSP, Goal-test at bottom



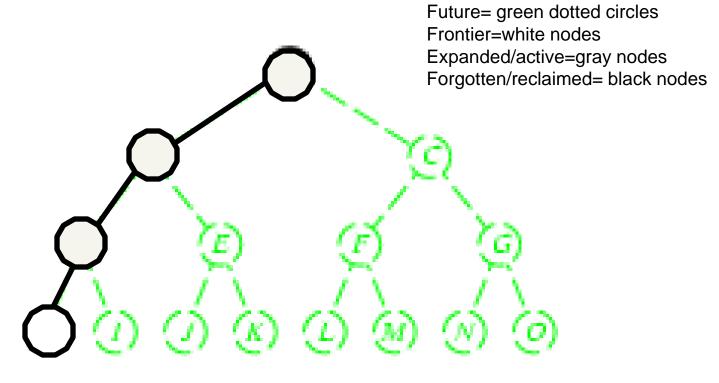
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 - For CSP, Goal-test at bottom



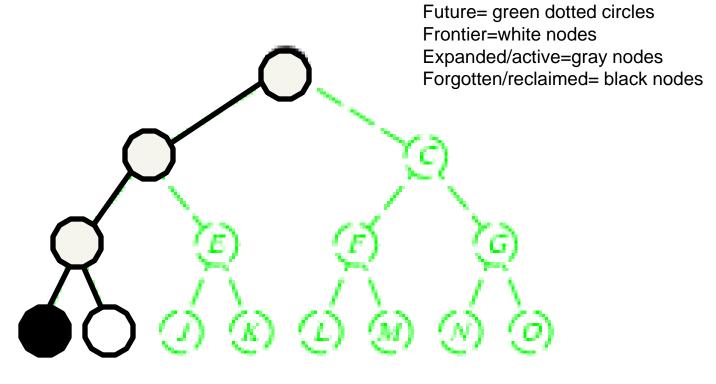
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Future= green dotted circles Frontier=white nodes Expanded/active=gray nodes Forgotten/reclaimed= black nodes (x) \mathcal{C}

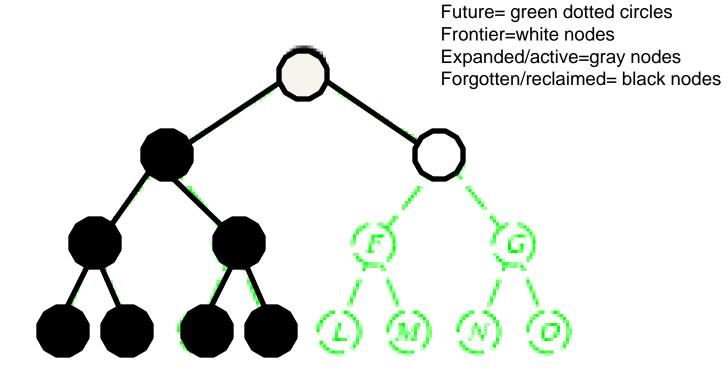
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Backtracking search (Figure 6.5)

function BACKTRACKING-SEARCH(csp) return a solution or failure
 return RECURSIVE-BACKTRACKING({}, csp)

function RECURSIVE-BACKTRACKING(assignment, csp) return a solution or failure if assignment is complete then return assignment $var \leftarrow SELECT-UNASSIGNED-VARIABLE(VARIABLES[csp], assignment, csp)$ for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do if value is consistent with assignment according to CONSTRAINTS[csp] add {var=value} to assignment result \leftarrow RECURSIVE-BACTRACKING(assignment, csp) if result \neq failure then return result remove {var=value} from assignment

then

return failure

Depth First Search

Depth-first search

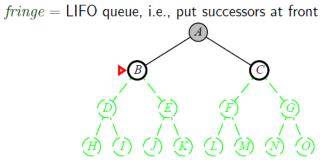
Expand deepest unexpanded node

Implementation:

fringe = LIFO queue, i.e., put successors at front

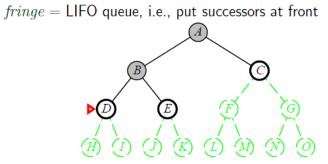
Expand deepest unexpanded node

Implementation:

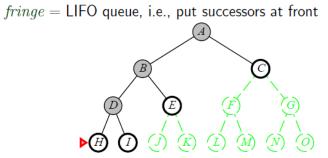


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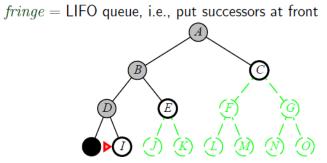
Expand deepest unexpanded node



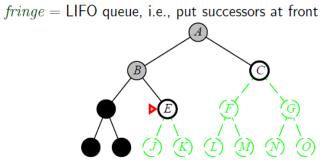
Expand deepest unexpanded node



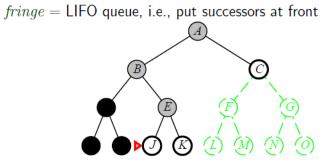
Expand deepest unexpanded node



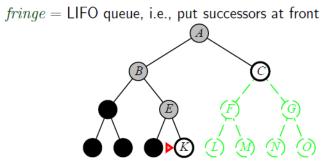
Expand deepest unexpanded node



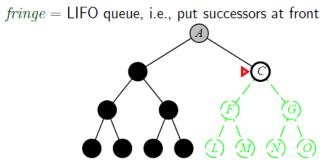
Expand deepest unexpanded node



Expand deepest unexpanded node



Expand deepest unexpanded node



Improving CSP efficiency

- Previous improvements on uninformed search
 → introduce heuristics
- For CSPS, general-purpose methods can give large gains in speed, e.g.,
 - Which variable should be assigned next?
 - In what order should its values be tried?
 - Can we detect inevitable failure early?
 - Can we take advantage of problem structure?

Note: CSPs are somewhat generic in their formulation, and so the heuristics are more general compared to methods in Chapter 4

Heuristic

- Selecting Variable
 - Minimum remaining values (MRV)
 - choose variable with the fewest legal moves
 - Degree heuristic for next variable
 - select variable that is involved in the largest number of constraints on other unassigned variables
 - useful as a tie breaker after MRV.
- Selecting Value
 - Least constraining value (LCV)
 - given a variable choose the least constraining value

Backtracking search (Figure 6.5)

function BACKTRACKING-SEARCH(csp) return a solution or failure
 return RECURSIVE-BACKTRACKING({}, csp)

function RECURSIVE-BACKTRACKING(assignment, csp) return a solution or failure

if assignment is complete then return assignment

var ← SELECT-UNASSIGNED-VARIABLE(VARIABLES[*csp*], *assignment*, *csp*)

for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do

if value is consistent with assignment according to CONSTRAINTS[csp] then

add {var=value} to assignment

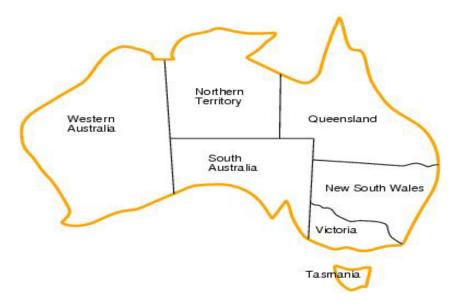
result \leftarrow RRECURSIVE-BACTRACKING(*assignment*, *csp*)

if *result* ≠ *failure* **then return** *result*

remove {var=value} from assignment

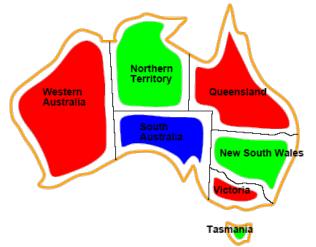
return failure

CSP example: Map coloring problem



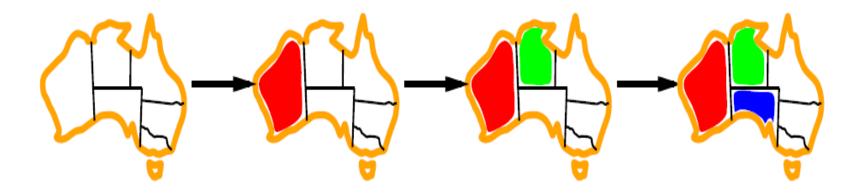
- Variables: WA, NT, Q, NSW, V, SA, T
- Domains: D_i = {red,green,blue}
- Constraints: adjacent regions must have different colors.
 - E.g. $WA \neq NT$

CSP example: Map coloring solution



- A solution is:
 - <u>A complete and consistent assignment.</u>
 - All variables assigned, all constraints satisfied.
- E.g., {*WA*=*red*, *NT*=*green*, *Q*=*red*, *NSW*=*green*, *V*=*red*, *SA*=*blue*, *T*=*green*}

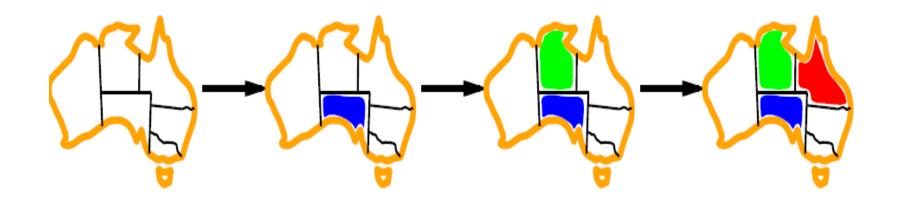
Minimum remaining values (MRV) for next variable



var ← SELECT-UNASSIGNED-VARIABLE(VARIABLES[*csp*],*assignment*,*csp*)

- A.k.a. most constrained variable heuristic
- *Heuristic Rule*: choose variable with the fewest legal moves
 - e.g., will immediately detect failure if X has no legal values

Degree heuristic for next variable



- *Heuristic Rule*: select variable that is involved in the largest number of constraints on other unassigned variables.
- Degree heuristic can be useful as a tie breaker after MRV.
- In what order should a variable's values be tried?

Backtracking search (Figure 6.5)

function BACKTRACKING-SEARCH(csp) return a solution or failure
 return RECURSIVE-BACKTRACKING({}, csp)

function RECURSIVE-BACKTRACKING(assignment, csp) return a solution or failure

if assignment is complete then return assignment

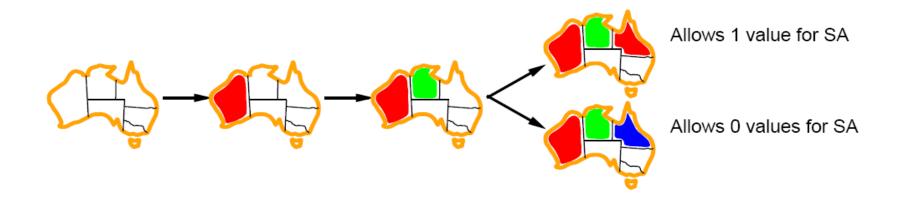
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for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do

if value is consistent with assignment according to CONSTRAINTS[csp] then
 add {var=value} to assignment
 result ← RRECURSIVE-BACTRACKING(assignment, csp)
 if result ≠ failure then return result
 remove {var=value} from assignment

return failure

Least constraining value (LCV) for next value



- Least constraining value heuristic
- Heuristic Rule: given a variable choose the least constraining value
 - leaves the maximum flexibility for subsequent variable assignments

Minimum remaining values (MRV) vs. Least constraining value (LCV)

- Why do we want the MRV (minimum values, most constraining) for variable selection --- but the LCV (maximum values, least constraining) for value selection?
- Isn't there a contradiction here?
- MRV for variable selection to reduces the branching factor.
 - Smaller branching factors lead to faster search.
 - Hopefully, when we get to variables with currently many values, constraint propagation (next lecture) will have removed some of their values and they'll have small branching factors by then too.
- LCV for value selection increases the chance of early success.
 - If we are going to fail at this node, then we have to examine every value anyway, and their order makes no difference at all.
 - If we are going to succeed, then the earlier we succeed the sooner we can stop searching, so we want to succeed early.
 - LCV rules out the fewest possible solutions below this node, so we have the most chances for early success.