CS-171, Intro to A.I., Winter Quarter, 2014 — Quiz # 1 — 20 minutes

1. NAME:

YOUR ID: ______ ID TO RIGHT: ______ ROW: ____ NO. FROM RIGHT: _____

2. (25 pts total, -5 pts each error, but not negative) Search Properties. Fill in the

values of the four evaluation criteria for each search search where b is the finite branching factor; d is the maximum depth of the search tree; l is the d equal to some positive ε ; in bidirectional search bo more than ±1, e.g., O(b^d) vs. O(b^(d+1)).

Note: These assumptions are the same as in Figure 3.21 of your textbook.

	Complete?	Time complexity	Space complexity	Optimal?
Depth-First	No	O(b^m)	O(bm)	No
Breadth-First	Yes	O(b^d)	O(b^d)	Yes
Uniform-Cost	Yes	O(b ^{(1+floor(C*/ε)))} O(b ^(d+1)) also OK	O(b ^{(1+floor(C*/ε)))} O(b ^(d+1)) also OK	Yes
Depth-Limited	No	O(b시)	O(bl)	No
Iterative Deepening	Yes	O(b^d)	O(bd)	Yes
Bidirectional (if applicable)	Yes	O(b^(d/2))	O(b^(d/2))	Yes

3. (30 pts total, -5 pts each error, but not negative) Reasoning about Search.

Note: Assumptions are DIFFERENT from problem 2 above. REASON about them.

Assume that you are doing Tree Search, the state space is infinitely deep, the branching factor is finite, there are cycles and loops, multiple goal nodes exist with different costs, step costs may differ from each other and are always greater than some given positive constant, in bidirectional search both directions use breadth-first search, and the heuristic function is consistent.

These assumptions represent a typical ill-conditioned search space.

3a. Is depth-first search complete? N optimal? N ("Y" = yes, "N" = no) 3b. Is breadth-first search complete? Y optimal? N ("Y" = yes, "N" = no) 3c. Is uniform-cost search complete? Y optimal? Y ("Y" = yes, "N" = no) 3d. Is depth-limited search complete? N optimal? N ("Y" = yes, "N" = no) 3e. Is iterated-deepening search complete? Y optimal? N ("Y" = yes, "N" = no) 3f. Is bidirectional search complete? Y optimal? N ("Y" = yes, "N" = no) 3g. Is greedy best-first search complete? N optimal? N ("Y" = yes, "N" = no) 3h. Is A* search complete? Y optimal? N ("Y" = yes, "N" = no)

**** TURN PAGE OVER AND CONTINUE ON THE OTHER SIDE ****

4. (**45 pts total**, **9 pts each**) Execute Tree Search through this graph (i.e., do not remember visited nodes). Step costs are given next to each arc. Heuristic values are given next to each node (as h=x). The successors of each node are indicated by the arrows out of that node. Successors are returned in left-to-right order. Specifically, the children of S are (A, B) and the children of B are (C, B), in that order.

Specifically, the children of S are (A, B) and the children of B are (C, D), ill that order.				
For each search strategy below, show the order in Minor errors will receive partial credit. The Reader will determine how much credit to allow in any particular case.	ee the lecture slides for Uninformed Search, topic to do Goal-Test? When generated? When popped?" fication about exactly what to do in practical cases.				
A h=15 $h=8 C$	$\begin{array}{c} 1 \\ 4 \\ 9 \\ \end{array} \\ \hline \mathbf{G} \\ \mathbf{h}=0 \\ \end{array} \\ \begin{array}{c} 10 \\ 10 \\ \mathbf{B} $				
4.a. DEPTH FIRST SEARCH.					
Order of node expansion: <u>S A C G</u> See section 3.4.3.					
Path found: <u>S A C G</u>	Cost of path found: 25				
4 b (0 pts) RDFADTH FIDST SFADCH	BFS does the Goal-test before the				
4.0. (9 pts) BREADTH FIRST SEARCH. See section 3.4.1.					
Order of node expansion: <u>SABCG</u> goal is found when C is expanded.					
Path found: SACG	Cost of path found: 25				
4.c. (9 pts) UNIFORM COST SEARCH.	s goaltast when node is nonned off guoue				
Order of node expansion: <u>SBCABG</u>					
Path found: SBCG	Cost of path found: 14				
4.d. (9 pts) GREEDY (BEST-FIRST) SEARCH.					
Order of node expansion: SBBBBBBB	has lower h(=4) than any other node on queue.				
Path found: none See section 3.5.1.	Cost of path found: none				
4.e. (9 pts) ITERATED DEEPENING SEARCH.	IDS does the Goal-test before the child				
	is pushed onto the queue. The goal is				
Order of node expansion: <u>SSABSACG</u>	found when D is expanded.				
Path found: <u>SACG</u>	Cost of path found: 25				
4.f. (9 pts) A* SEARCH. See section 3.5.2. A* does goa	altest when node is popped off queue.				
Order of node expansion: <u>SBCG</u>					
Path found: <u>SBCG</u>	Cost of path found: 14				
Is the heuristic admissible? (Yes or No) Y					