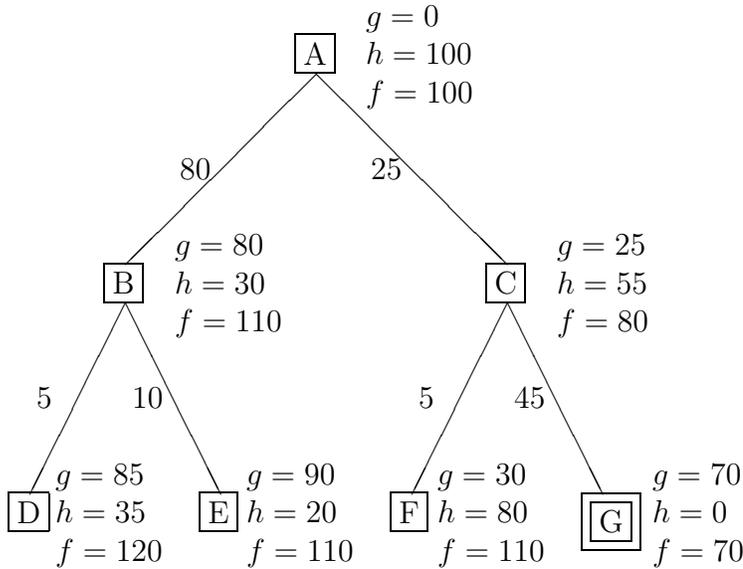


ICS 171 — Quiz #2 — FIFTEEN (15) minutes

1. (5 pts) NAME AND EMAIL ADDRESS: _____

YOUR ID: _____ ID TO RIGHT: _____ ROW: _____ NO. FROM RIGHT: _____

2. (5 pts each, 30 pts total) Use the following tree to indicate the order that nodes are expanded, for different types of search. Assume that G (double box) is the only goal node. Here, path costs are shown to the right of each path, g = cost of path so far, h = estimate of remaining cost to goal, f = estimate of total path cost.



For each search strategy, write down the order in which nodes are expanded. Stop at G.

(2a) DEPTH-FIRST SEARCH:

A B D E C F G _____

(2b) BREADTH-FIRST SEARCH:

A B C D E F G _____

(2c) ITERATIVE DEEPENING DEPTH-FIRST SEARCH:

A A B C A B D E C F G

(2d) UNIFORM-COST SEARCH:

A C F G _____

(2e) GREEDY BEST-FIRST SEARCH:

A B E D C G _____

(2f) A* SEARCH:

A C G _____

3. (5 pts , 10 pts total) Use the tree above.

(3a) Y (Y=yes or N=no) Is the heuristic in problem 2 above **admissible**?

(3b) Y (Y=yes or N=no) Is the heuristic in problem 2 above **consistent (monotonic)**?

***** GO ON TO THE NEXT PAGE *****

4. (10 pts) In general, which is the preferred search method when (a) there is a large search space, (b) the depth of the solution is unknown, (c) an optimal solution is desired, and (d) a consistent admissible heuristic is available? (Mark one blank with “X”)

depth-first search breadth-first search uniform-cost search
 depth-limited search A* search bidirectional search
 greedy best-first search iterative-deepening search

5. (5 pts each, 30 pts total) Recall that

- True path cost so far = $g(n)$.
- Estimated cost to goal = $h'(n)$.
- Estimated total cost = $f'(n) = g(n) + h'(n)$.

The following is a proof that A* search (queue sorted by f') is optimal if the heuristic is admissible. The lines have been labelled A through G. Unfortunately, they have been scrambled.

Let n_g be the first goal node popped off the queue. Let n_o be any other node on the queue. We wish to prove that n_o can never be extended to a path to any goal node that costs less than the path to n_g that we just found.

A : true total cost of n_g
 F : = $g(n_g)$ // because n_g represents a complete path
 D : = $f'(n_g)$ // by definition of f with $h'(n_g) = 0$
 B : ≤ $f'(n_o)$ // because queue is sorted by f
 E : = $g(n_o) + h'(n_o)$ // by definition of f
 C : ≤ $g(n_o) + \text{true cost to goal from } n_o$ // because h' is admissible
 G : = true total cost of n_o

Fill in the blanks with the letters B, C, D, E, F, and G to prove that the true total cost of $n_g \leq$ true total cost of n_o . The first and last letters, A and G, have been done for you as an example.

 A F D B E C G

6. (5 pts each, 15 pts total) Label the following as “T” (true) or “F” (false).

6.a. F An admissible heuristic never underestimates the remaining distance (cost) to the goal.

6.b. F A good way to find a heuristic for a problem is to make the problem harder and then solve the new harder problem.

6.c. F Let h_1 and h_2 be two admissible consistent heuristics. Heuristic h_1 is stronger (better) than heuristic h_2 just in case, for every search node n , we have $h_1(n) \geq h_2(n)$.