## Causal and Probabilistic Reasoning

# Slides Set 5: <br> Exact Inference Algorithms Bucket-elimination 

## Rina Dechter

(Dechter chapter 4, Darwiche chapter 6)


## Inference for probabilistic networks

- Bucket elimination (Dechter chapter 4)
- Belief-updating, P(e), partition function
- Marginals, probability of evidence
- The impact of evidence
- for MPE ( $\rightarrow$ MAP)
- for MAP ( $\rightarrow$ Marginal Map)
- Influence diagrams ?
- Induced-Width (Dechter, Chapter 3.4)


## Inference for probabilistic networks

- Bucket elimination
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## Bayesian Networks: Example <br> (Pearl, 1988)


$P(S, C, B, X, D)=P(S) P(C \mid S) P(B \mid S) P(X \mid C, S) P(D \mid C, B)$
Belief Updating:
P (lung cancer=yes | smoking=no, dyspnoea=yes ) = ?

## A Bayesian Network

| $A$ | $C$ | $\Theta_{C \mid A}$ |
| :--- | :--- | :--- |
| true | true | .8 |
| true | false | .2 |
| false | true | .1 |
| false | false | .9 |


| $B$ | $C$ | $D$ | $\Theta_{D \mid B C}$ |
| :--- | :--- | :--- | :--- |
| true | true | true | .95 |
| true | true | false | .05 |
| true | false | true | .9 |
| true | false | false | .1 |
| false | true | true | .8 |
| false | true | false | .2 |
| false | false | true | 0 |
| false | false | false | 1 |



| $C$ | $E$ | $\Theta_{E \mid C}$ |
| :--- | :--- | :--- |
| true | true | .7 |
| true | false | .3 |
| false | true | 0 |
| false | false | 1 |

## Types of queries



- NP-hard: exponentially many terms
- We will focus on exact and then on approximation algorithms
- Anytime: very fast \& very approximate ! Slower \& more accurate


## Belief Updating is NP-hard

- Each SAT formula can be mapped into a belief updating query in a Bayesian network
- Example

$$
(\neg u \vee \neg w \vee y) \wedge(u \vee \neg v \vee w)
$$

## A Simple Network

Given:


- How can we compute $P(D)$ ?, $P(D \mid A=0)$ ? $P(A \mid D=0)$ ?
- Brute force $\mathrm{O}\left(k^{4}\right)$
- Maybe $O\left(4 k^{2}\right)$


## Elimination as a Basis for Inference



| $A$ | $\Theta_{A}$ |
| :--- | :--- |
| true | .6 |
| false | .4 |


| $A$ | $B$ | $\Theta_{B \mid A}$ |
| :--- | :--- | :--- |
| true | true | .9 |
| true | false | .1 |
| false | true | .2 |
| false | false | .8 |


| $B$ | $C$ | $\Theta_{C \mid B}$ |
| :--- | :--- | :--- |
| true | true | .3 |
| true | false | .7 |
| false | true | .5 |
| false | false | .5 |

To compute the prior marginal on variable $C, \operatorname{Pr}(C)$
we first eliminate variable $A$ and then variable $B$

## Elimination as a Basis for Inference

- There are two factors that mention variable $A, \Theta_{A}$ and $\Theta_{B \mid A}$
- We multiply these factors first and then sum out variable $A$ from the resulting factor.
- Multiplying $\Theta_{A}$ and $\Theta_{B \mid A}$ :

| $A$ | $B$ | $\Theta_{A} \Theta_{B \mid A}$ |
| :--- | :--- | :--- |
| true | true | .54 |
| true | false | .06 |
| false | true | .08 |
| false | false | .32 |

- Summing out variable $A$ :

| $B$ | $\sum_{A} \Theta_{A} \Theta_{B \mid A}$ |
| :--- | :--- |
| true | $.62=.54+.08$ |
| false | $.38=.06+.32$ |

## Elimination as a Basis for Inference

- We now have two factors, $\sum_{A} \Theta_{A} \Theta_{B \mid A}$ and $\Theta_{C \mid B}$, and we want to eliminate variable $B$
- Since $B$ appears in both factors, we must multiply them first and then sum out $B$ from the result.
- Multiplying:

| $B$ | $C$ | $\Theta_{C \mid B} \sum_{A} \Theta_{A} \Theta_{B \mid A}$ |
| :--- | :--- | :--- |
| true | true | .186 |
| true | false | .434 |
| false | true | .190 |
| false | false | .190 |

- Summing out:

| $C$ | $\sum_{B} \Theta_{C \mid B} \sum_{A} \Theta_{A} \Theta_{B \mid A}$ |
| :--- | :--- |
| true | .376 |
| false | .624 |

## Elimination as a Basis for Inference

- We now have two factors, $\sum_{A} \Theta_{A} \Theta_{B \mid A}$ and $\Theta_{C \mid B}$, and we want to eliminate variable $B$
- Since $B$ appears in both factors, we must multiply them first and then sum out $B$ from the result.
- Multiplying:

| $B$ | $C$ | .186 |
| :--- | :--- | :--- |
| true | true | .18 |
| true | false | .434 |
| false | true | .190 |
| false | false | .190 |

- Summing out:



## Belief Updating



P (lung cancer=yes / smoking=no, dyspnoea=yes ) = ?

## Belief updating: $P(X \mid$ evidence $)=$ ?



$$
P(a \mid e=0) \propto P(a, e=0)=
$$

$$
\begin{aligned}
& \sum_{e=0, d, c, b} P(a) \underbrace{P(b \mid a)} P(c \mid a) \underbrace{P(d \mid b, a) P(e \mid b, c)}_{1}= \\
& P(a) \sum_{e=0}^{\sum_{e=0}^{\sum_{d}} \sum_{c} P\left(c \mid a \sum_{c} P(b \mid a) P(d \mid b, a) P(e \mid b, c)\right.} \\
& \text { Variable Elimination }
\end{aligned}
$$

# Bucket elimination Algorithm BE-bel (Dechter 1996) 

$$
P(A \mid E=0)=\alpha \sum_{E=0, D, C, B} P(A) \cdot P(B \mid A) \cdot P(C \mid A) \cdot P(D \mid A, B) \cdot P(E \mid B, C)
$$

$$
\begin{aligned}
& \text { bucket } B \text { : } \\
& \text { bucket } C \text { : } \\
& \text { bucket } D \text { : }
\end{aligned}
$$

$$
\overbrace{P(b \mid a) P(d \mid b, a)} P(e \mid b, c)
$$

bucket E:

$$
e=0 \quad \lambda^{\boldsymbol{D}}(\boldsymbol{a}, \boldsymbol{e})
$$

$$
W^{*}=4
$$

bucket A:

$P(a, e=0)$ slides $P\left(z_{a} \mid e_{4}=0\right)=\frac{P(a, e=0)}{P(e=0)}$


## A Bayesian Network Ordering: A,C,B,E,D,G


(a) Directed acyclic graph

(b) Moral graph

$$
\begin{align*}
P(a, g=1) & =\sum_{c, b, e, d, g=1} P(a, b, c, d, e, g)=\sum_{c, b, f, d, g=1} P(g \mid f) P(f \mid b, c) P(d \mid a, b) P(c \mid a) P(b \mid a) P(a) . \\
P(a, g=1) & =P(a) \sum_{c} P(c \mid a) \sum_{b} P(b \mid a) \sum_{f} P(f \mid b, c) \sum_{d} P(d \mid b, a) \sum_{g=1} P(g \mid f) .  \tag{4.1}\\
P(a, g=1) & =P(a) \sum_{c} P(c \mid a) \sum_{b} P(b \mid a) \sum_{f} P(f \mid b, c) \lambda_{G}(f) \sum_{d} P(d \mid b, a)  \tag{4.2}\\
P(a, g=1) & =P(a) \sum_{c} P(c \mid a) \sum_{b} P(b \mid a) \lambda_{D}(a, b) \sum_{f} P(f \mid b, c) \lambda_{G}(f)  \tag{4.3}\\
P(a, g=1) & =P(a) \sum_{c} P(c \mid a) \sum_{b} P(b \mid a) \lambda_{D}(a, b) \lambda_{F}(b, c)  \tag{4.4}\\
P(a, g=1) & =P(a) \sum_{c} P(c \mid a) \lambda_{B}(a, c) \tag{4.5}
\end{align*}
$$

## A Bayesian Network Ordering: A,C,B,E,D,G


(a) Directed acyclic graph

(b) Moral graph

$$
\begin{align*}
& P(a, g=1)=\sum_{c, b, e, d, g=1} P(a, b, c, d, e, g)=\sum_{c, b, f, d, g=1} P(g \mid f) P(f \mid b, c) P(d \mid a, b) P(c \mid a) P(b \mid a) P(a) . \\
& P(a, g=1)=P(a) \sum_{c} P(c \mid a) \sum_{b} P(b \mid a) \sum_{f} P(f \mid b, c) \sum_{d} P\left(d \mid b, a \sum_{d=1} P(g \mid f) .\right.  \tag{4.1}\\
& P(a, g=1)=P(a) \sum_{c} P(c \mid a) \sum_{b} P(b \mid a) \sum_{f} P(f \mid b, c) \lambda_{G}(f) \sum_{d} P(d \mid b, a) .  \tag{4.2}\\
& P(a, g=1)=P(a) \sum_{c} P(c \mid a) \sum_{b} P(b \mid a) \lambda_{D}(a, b) \sum_{\sum_{f} P(f \mid b, c) \lambda_{G}(f)}  \tag{4.3}\\
& P(a, g=1)=P(a) \sum_{c} P(c \mid a) \underbrace{P(a, g=1)}_{\sum_{b} P(b \mid a) \lambda_{D}(a, b) \lambda_{F}(b, c)}=P(a) \sum_{c} P(c \mid a) \lambda_{B}(a, c) \tag{4.4}
\end{align*}
$$

## A Bayesian Network Ordering: A,C,B,F,D,G



(a) Directed acyclic graph

(b) Moral graph

## A Different Ordering


(a) Directed acyclic graph

(b) Moral graph

## Ordering: $A, F, D, C, B, G$




# A Different Ordering 


(a) Directed acyclic graph

(b) Moral graph

## Ordering: $A, F, D, C, B, G$

$P(a, g=1)=P(a) \sum_{f} \sum_{d} \sum_{c} P(c \mid a) \sum_{b} P(b \mid a) P(d \mid a, b) P(f \mid b, c) \sum_{g=1} P(g \mid f)$
$=P(a) \sum_{f} \lambda_{G}(f) \sum_{d} \sum_{c} P(c \mid a) \sum_{b} P(b \mid a) P(d \mid a, b) P(f \mid b, c)$
$=P(a) \sum_{f} \lambda_{G}(f) \sum_{d} \sum_{c} P(c \mid a) \lambda_{B}(a, d, c, f)$
$=P(a) \sum_{f} \lambda_{g}(f) \sum_{d} \lambda_{C}(a, d, f)$
$=P(a) \sum_{f} \lambda_{G}(f) \lambda_{D}(a, f)$
$=P(a) \lambda_{F}(a)$

(a)

(b)

## A Bayesian Network Processed Along 2 Orderings

$\Sigma \Pi$

(a) Directed acyclic graph

(a)

(b) Moral graph

(b)
$d 1=A, C, B, F, D, G$
Figure 4.4: The bucket's output when processing along $d_{2}=A, F, D, C, B, G$.

## The Operation In a Bucket

- Multiplying functions
- Marginalizing (summing-out) functions


## Combination of Cost Functions

| $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{f}(\mathbf{A}, \mathbf{B})$ |
| :---: | :---: | :---: |
| b | b | 0.4 |
| b | g | 0.1 |
| g | b | 0 |
| g | g | 0.5 |



## Factors: Sum-Out Operation

## The result of summing out variable $X$ from factor $f(\mathbf{X})$

is another factor over variables $\mathbf{Y}=\mathbf{X} \backslash\{X\}$ :

$$
\left(\sum_{x} f\right)(\mathbf{y}) \stackrel{\text { def }}{=} \sum_{x} f(x, \mathbf{y})
$$

| $B$ | $C$ | $D$ | $f_{1}$ |
| :--- | :--- | :--- | :--- |
| true | true | true | .95 |
| true | true | false | .05 |
| true | false | true | .9 |
| true | false | false | .1 |
| false | true | true | .8 |
| false | true | false | .2 |
| false | false | true | 0 |
| false | false | false | 1 |


| $B$ | $C$ | $\sum_{D} f_{1}$ |
| :--- | :--- | :--- |
| true | true | 1 |
| true | false | 1 |
| false | true | 1 |
| false | false | 1 |

$$
\top \quad \sum_{B} \sum_{C} \sum_{D} f_{1}
$$

## Bucket Elimination and Induced Width



## Bucket Elinnination and induced Midth



## Algorithm BE-bel

Input: A belief network $\mathcal{B}=\left\langle\mathbf{X}, \mathbf{D}, \mathbf{P}_{G}, \Pi\right\rangle$, an ordering $d=\left(X_{1}, \ldots, X_{n}\right)$; evidence $e$ output: The belief $P\left(X_{1} \mid \mathrm{e}\right)$ and probability of evidence $P(\mathrm{e})$

1. Partition the input functions (CPTs) into bucket $_{1}, \ldots$, bucket $_{n}$ as follows: for $i \leftarrow n$ downto 1 , put in bucket $_{i}$ all unplaced functions mentioning $X_{i}$. Put each observed variable in its bucket. Denote by $\psi_{i}$ the product of input functions in bucket ${ }_{i}$.
2. backward: for $p \leftarrow n$ downto 1 do
3. for all the functions $\psi_{S_{0}}, \lambda_{S_{1}}, \ldots, \lambda_{S_{j}}$ in bucket $_{p}$ do

If (observed variable) $X_{p}=x_{p}$ appears in bucket $_{p}$, assign $X_{p}=x_{p}$ to each function in bucket $_{p}$ and then put each resulting function in the bucket of the closest variable in its scope. else,
4. $\quad \lambda_{p} \leftarrow \sum_{X_{p}} \psi_{p} \cdot \prod_{i=1}^{j} \lambda_{S_{i}}$
5. place $\lambda_{p}$ in bucket of the latest variable in scope $\left(\lambda_{p}\right)$,
6. return (as a result of processing bucket $_{1}$ ):

$$
\begin{aligned}
& P(\mathrm{e})=\alpha=\sum_{X_{1}} \psi_{1} \cdot \prod_{\lambda \in \text { bucket }_{1}} \lambda \\
& P\left(X_{1} \mid \mathrm{e}\right)=\frac{1}{\alpha} \psi_{1} \cdot \prod_{\lambda \in \text { bucket }_{1}} \lambda
\end{aligned}
$$

Figure 4.5: BE-bel: a sum-product bucket-elimination algorithm.

## Student Network Example



## Induced Width (continued)

$w^{*}(d)$ - the induced width of the primal graph along ordering $d$ The effect of the ordering:



$$
w^{*}\left(d_{1}\right)=4
$$



$$
w^{*}\left(d_{2}\right)=2
$$

## Inference for Probabilistic Networks

- Bucket elimination
- Belief-updating, P(e), partition function
- Marginals, probability of evidence
- The impact of evidence
- for MPE ( $\rightarrow$ MAP)
- for MAP ( $\rightarrow$ Marginal Map)
- Induced-Width

The Impact of Evidence? Algorithm BE-bel


## The Impact of Evidence?

 Algorithm BE-bel$$
P(A \mid E=0)=\alpha \sum_{E=0, D, C, B} P(A) \cdot P(B \mid A) \cdot P(C \mid A) \cdot P(D \mid A, B) \cdot P(E \mid B, C)
$$

$$
P(A / E=0, B=1) ?
$$

bucket B:

$$
\overbrace{P(b \mid a) \quad P(d \mid b, a) \quad P(e \mid b, c)}^{b} \quad B=1
$$

bucket $C$ :

$$
P(c \mid a) \quad P(e \mid b=1, c)
$$

bucket E:

$$
e=0
$$

bucket A:

${ }_{\text {slides 5 }}^{276}\left(\boldsymbol{a} \mid{ }_{20}=\mathbf{e}=\mathbf{0}\right)=\frac{\boldsymbol{P}(\boldsymbol{a}, \boldsymbol{e}=\boldsymbol{0})}{\boldsymbol{P}(\boldsymbol{e}=0)}$

## The Impact of Observations



(a)

(b)

(c)

Figure 4.9: Adjusted induced graph relative to observing $B$.

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## $\mathrm{MPE}=\max \mathrm{P}(\overline{\mathrm{x}})$ <br> $\overline{\mathrm{X}}$



$$
\begin{aligned}
& \sum_{M P E} \text { is replaced by } \max : \\
& \max _{a, e, d, c, b} P(a) P(c \mid a) P(b \mid a) P(d \mid a, b) P(e \mid b, c)
\end{aligned}
$$

## $\mathrm{MPE}=\max \mathrm{P}(\overline{\mathrm{x}})$ <br> $\overline{\mathrm{x}}$



$$
\sum \text { is replaced by } \max :
$$

$M P E=\max _{a, e, d, c, b} P(a) P(c \mid a) P(b \mid a) P(d \mid a, b) P(e \mid b, c)$
bucket D:
bucket E:
bucket A:


## Generating the MPE-tuple

5. $b^{\prime}=\arg \max P\left(b / a^{\prime}\right) \times$

$$
\times P\left(d^{\prime} \mid b, a^{b^{\prime}}\right) \times P\left(e^{\prime} \mid b, c^{\prime}\right)
$$

4. $c^{\prime}=\arg \max P\left(c / a^{\prime}\right) \times$

$$
\times h^{B}\left(a^{\prime}, d^{c}, c, e^{\prime}\right)
$$

3. $d^{\prime}=\arg \max _{d} h^{c}\left(a^{\prime}, d, e^{\prime}\right)$
4. $e^{\prime}=0$
5. $a^{\prime}=\arg \max _{a} P(a) \cdot h^{E}(a)$
$B: P(b \mid a) \quad P(d \mid b, a) \quad P(e \mid b, c)$
$C: \quad P(c \mid a) \quad h^{B}(\boldsymbol{a}, \boldsymbol{d}, \boldsymbol{c}, \boldsymbol{e})$

D:
$h^{c}(a, d, e)$
$E: \quad e=0 \quad h^{D}(\mathbf{a}, \boldsymbol{e})$

A: $\quad P(a) \quad h^{E}(a)$ Return ( $\left.a^{\prime}, b^{\prime}, c^{\prime}, d^{\prime}, e^{\prime}\right)$

## Induced Width

- Width is the max number of parents in the ordered graph
- Induced-width is the width of the induced ordered graph: recursively connecting parents going from last node to first.
- Induced-width $\mathrm{w}^{*}(\mathrm{~d})$ is the max induced-width over all nodes in ordering d
- Induced-width of a graph, $\mathrm{w}^{*}$ is the $\min \mathrm{w}^{*}(\mathrm{~d})$ over all orderings d

primal graph

$w^{*}\left(d_{1}\right)=4$

$w^{*}\left(d_{2}\right)=2$


## Complexity of Bucket Elimination

 Bucket-Elimination is time and space $O\left(r \exp \left(w_{d}^{*}\right)\right)$$w_{d}^{*}$ : the induced width of the primal graph along ordering d
$r=$ number of functions $\quad$ The effect of the ordering:

primal
graph Finding smallest induced-width is hard!

## A Bayesian Network

## Example with mpe?



| $A$ | $C$ | $\Theta_{C \mid A}$ |
| :--- | :--- | :--- |
| true | true | .8 |
| true | false | .2 |
| false | true | .1 |
| false | false | .9 |


| $B$ | $C$ | $D$ | $\Theta_{D \mid B C}$ |
| :--- | :--- | :--- | :--- |
| true | true | true | .95 |
| true | true | false | .05 |
| true | false | true | .9 |
| true | false | false | .1 |
| false | true | true | .8 |
| false | true | false | .2 |
| false | false | true | 0 |
| false | false | false | 1 |


| $C$ | $E$ | $\Theta_{E \mid C}$ |
| :--- | :--- | :--- |
| true | true | .7 |
| true | false | .3 |
| false | true | 0 |
| false | false | 1 |

## Try to compute MPE when E=0



| $A$ | $C$ | $\Theta_{C \mid A}$ |
| :--- | :--- | :--- |
| true | true | .8 |
| true | false | .2 |
| false | true | .1 |
| false | false | .9 |


| $B$ | $C$ | $D$ | $\Theta_{D \mid B C}$ |
| :--- | :--- | :--- | :--- |
| true | true | true | .95 |
| true | true | false | .05 |
| true | false | true | .9 |
| true | false | false | .1 |
| false | true | true | .8 |
| false | true | false | .2 |
| false | false | true | 0 |
| false | false | false | 1 |


| $C$ | $E$ | $\Theta_{E \mid C}$ |
| :--- | :--- | :--- |
| true | true | .7 |
| true | false | .3 |
| false | true | 0 |
| false | false | 1 |

## Complexity of Bucket-Elimination

- Theorem:
$B E$ is $O\left(n \exp \left(w^{*}+1\right)\right)$ time and $O\left(n \exp \left(w^{*}\right)\right)$ space, when $\mathrm{w}^{*}$ is the induced-width of the moral graph along d when evidence nodes are processed (edges from evidence nodes to earlier variables are removed.)

More accurately: $O\left(r \exp \left(w^{*}(d)\right)\right.$ where $r$ is the number of CPTs. For Bayesian networks r=n. For Markov networks?

## Inference for probabilistic networks

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- Marginals, probability of evidence
- The impact of evidence
- for MPE ( $\rightarrow$ MAP)
- for MAP ( $\rightarrow$ Marginal Map)
. Induced-Width (Dechter 3.4,3.5)


## Finding a Small Induced-Width

- NP-complete
- A tree has induced-width of ?
- Greedy algorithms:
- Min width
- Min induced-width
- Max-cardinality and chordal graphs
- Fill-in (thought as the best)
- Anytime algorithms
- Search-based [Gogate \& Dechter 2003]
- Stochastic (CVO) [Kask, Gelfand \& Dechter 2010]


## Finding a Small Induced-Width

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## Finding a Small Induced-Width

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## Min-width Ordering

MIN-WIDTH (MW)
input: a graph $G=(V, E), V=\left\{v_{1}, \ldots, v_{n}\right\}$
output: A min-width ordering of the nodes $d=\left(v_{1}, \ldots, v_{n}\right)$.

1. for $j=n$ to 1 by -1 do
2. $\quad r \leftarrow$ a node in $G$ with smallest degree.
3. put $r$ in position $j$ and $G \leftarrow G-r$.
(Delete from $V$ node $r$ and from $E$ all its adjacent edges)
4. endfor

Proposition: algorithm min-width finds a min-width ordering of a graph What is the Complexity of MW?

## Greedy Orderings Heuristics

- Min-induced-width
- From last to first, pick a node with smallest width, then connect parent and remove
. Min-Fill
- From last to first, pick a node with smallest fill-edges

Complexity? $O\left(n^{3}\right)$

## Min-Fill Heuristic

- Select the variable that creates the fewest "fill-in" edges



## Example


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## Different Induced-Graphs


(a)

(b)

(c)

(d)

A Miw ordering

## Which Greedy Algorithm is Best?

- Min-Fill, prefers a node who add the least number of fill-in arcs.
- Empirically, fill-in is the best among the greedy algorithms (MW,MIW,MF,MC)
- Complexity of greedy orderings?
- MW is $\mathrm{O}(\mathrm{e})$, MIW: $\mathrm{O}\left(n^{3}\right)$ MF $\mathrm{O}\left(n^{3}\right)$ MC is $\mathrm{O}(\mathrm{e}+\mathrm{n})$


## Propagation in Both Directions

- Messages can propagate both ways and we get beliefs for each variable

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## Inference for probabilistic networks

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- for MAP ( $\rightarrow$ Marginal Map)
- Influence diagrams ?
- Induced-Width (Dechter, Chapter 3.4)


## Marginal Map

| Max-Inference | $f\left(\mathbf{x}^{*}\right)=\max _{\mathbf{x}} \prod_{\alpha} f_{\alpha}\left(\mathbf{x}_{\alpha}\right)$ |
| :---: | :---: |
| Sum-Inference | $Z=\sum_{\mathbf{x}} \prod_{\alpha} f_{\alpha}\left(\mathbf{x}_{\alpha}\right)$ |
| Mixed-Inference | $f\left(\mathbf{x}_{M}^{*}\right)=\max _{\mathbf{x}_{\mathrm{M}}} \sum_{\mathbf{x}_{S}} \prod_{\alpha} f_{\alpha}\left(\mathbf{x}_{\alpha}\right)$ |

- NP-hard: exponentially many terms


## Example for MMAP Applications <br> 6 people, 3 markers

- Haplotype in Family pedigrees
- Coding networks

- Probabilistic planning
- Diagnosis



## Marginal MAP is Not Easy on Trees

- Pure MAP or summation tasks
- Dynamic programming
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## Bucket Elimination for MMAP

## Bucket Elimination


$\mathbf{X}_{M}=\{A, D, E\}$
$\mathbf{X}_{S}=\{B, C\}$
$\max _{\mathbf{X}_{M}} \sum_{\mathbf{X}_{S}} P(\mathbf{X})$

$M A P^{*}$ is the marginal MAP value


## Why is MMAP harder?



## Inference for probabilistic networks

- Bucket elimination (Dechter chapter 4)
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- for MPE ( $\rightarrow$ MAP)
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## Finding a Small Induced-Width

- NP-complete
- A tree has induced-width of ?
- Greedy algorithms:
- Min width
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- Max-cardinality and chordal graphs
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## Min-width Ordering

MIN-WIDTH (MW)
input: a graph $G=(V, E), V=\left\{v_{1}, \ldots, v_{n}\right\}$
output: A min-width ordering of the nodes $d=\left(v_{1}, \ldots, v_{n}\right)$.

1. for $j=n$ to 1 by -1 do
2. $\quad r \leftarrow$ a node in $G$ with smallest degree.
3. put $r$ in position $j$ and $G \leftarrow G-r$.
(Delete from $V$ node $r$ and from $E$ all its adjacent edges)
4. endfor

Proposition: algorithm min-width finds a min-width ordering of a graph What is the Complexity of MW?

## Greedy Orderings Heuristics

- Min-induced-width
- From last to first, pick a node with smallest width, then connect parent and remove
. Min-Fill
- From last to first, pick a node with smallest fill-edges

Complexity? $O\left(n^{3}\right)$

## Min-Fill Heuristic

- Select the variable that creates the fewest "fill-in" edges



## Example


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## Different Induced-Graphs


(a)

(b)

(c)

(d)

A Miw ordering

## Which Greedy Algorithm is Best?

- Min-Fill, prefers a node who add the least number of fill-in arcs.
- Empirically, fill-in is the best among the greedy algorithms (MW,MIW,MF,MC)
- Complexity of greedy orderings?
- MW is $\mathrm{O}(\mathrm{e})$, MIW: $\mathrm{O}\left(n^{3}\right)$ MF $\mathrm{O}\left(n^{3}\right)$ MC is $\mathrm{O}(\mathrm{e}+\mathrm{n})$


## Propagation in Both Directions

- Messages can propagate both ways and we get beliefs for each variable

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## Marginal Map

| Max-Inference | $f\left(\mathbf{x}^{*}\right)=\max _{\mathbf{x}} \prod_{\alpha} f_{\alpha}\left(\mathbf{x}_{\alpha}\right)$ |
| :---: | :---: |
| Sum-Inference | $Z=\sum_{\mathbf{x}} \prod_{\alpha} f_{\alpha}\left(\mathbf{x}_{\alpha}\right)$ |
| Mixed-Inference | $f\left(\mathbf{x}_{M}^{*}\right)=\max _{\mathbf{x}_{\mathrm{M}}} \sum_{\mathbf{x}_{S}} \prod_{\alpha} f_{\alpha}\left(\mathbf{x}_{\alpha}\right)$ |

- NP-hard: exponentially many terms


## Example for MMAP Applications <br> 6 people, 3 markers

- Haplotype in Family pedigrees
- Coding networks

- Probabilistic planning
- Diagnosis



## Marginal MAP is Not Easy on Trees

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