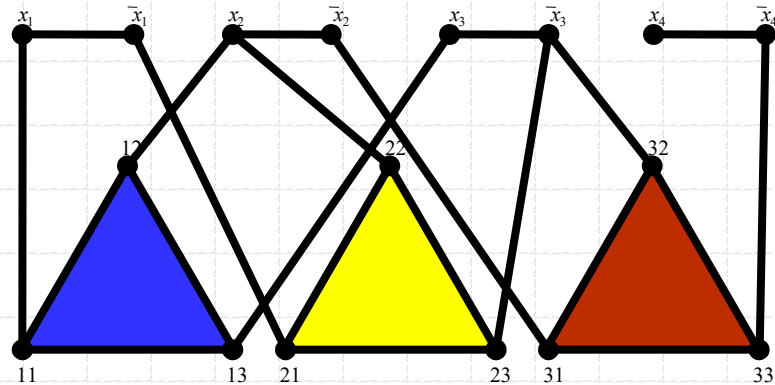


A Gentle Introduction to NP-Completeness



Michael T. Goodrich

Dealing with Hard Problems

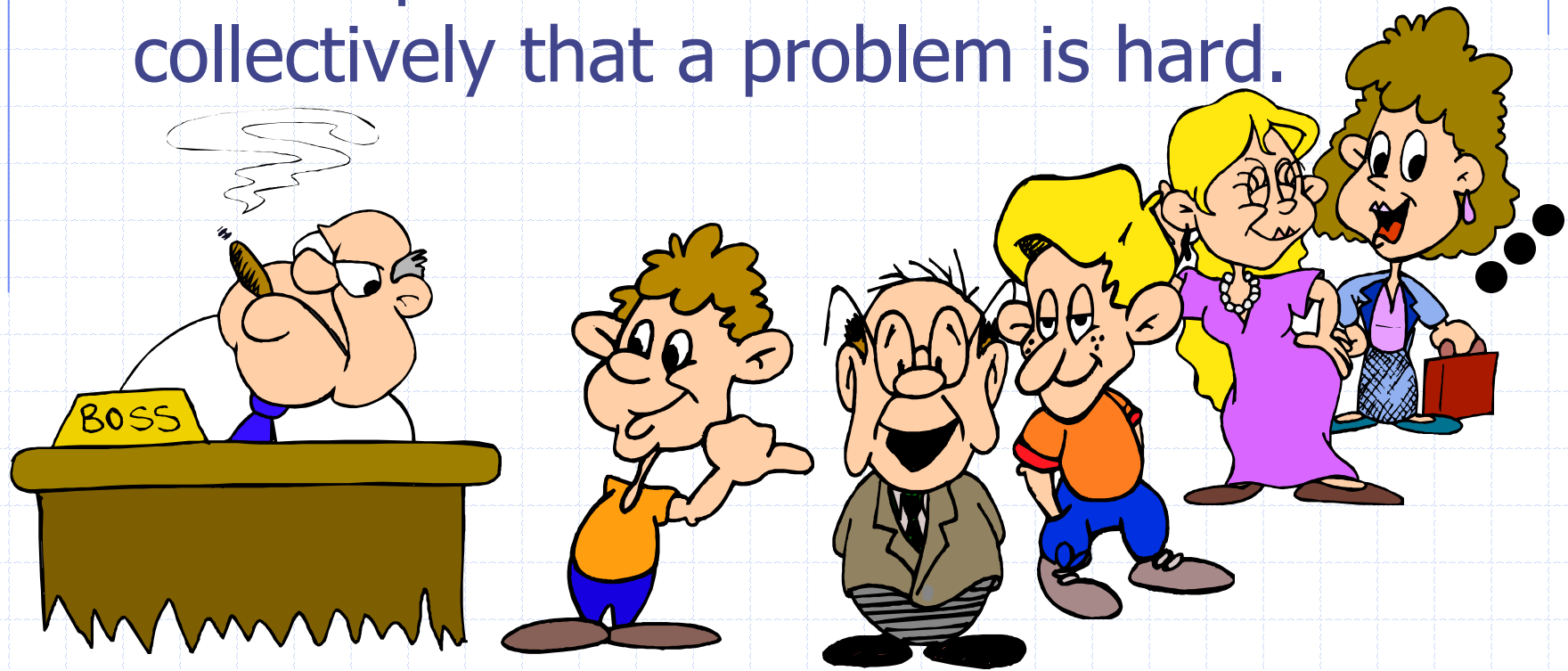
- ◆ What to do when we find a problem that looks hard...



I couldn't find a polynomial-time algorithm;
I guess I'm too dumb.

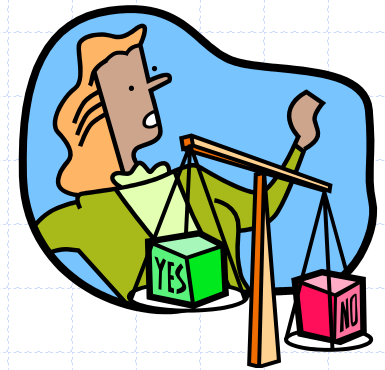
Dealing with Hard Problems

- ◆ NP-completeness let's us show collectively that a problem is hard.



I couldn't find a polynomial-time algorithm,
but neither could all these other smart people.

Polynomial-Time Decision Problems



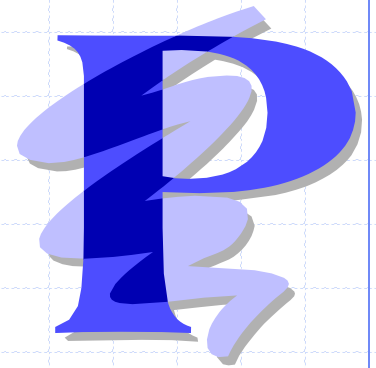
- ◆ To simplify the notion of “hardness,” we will focus on the following:
 - Polynomial-time as the cut-off for efficiency
 - Decision problems: output is 1 or 0 (“yes” or “no”)
 - ◆ Examples:
 - ◆ Does a text T contain a pattern P ?
 - ◆ Is the sequence, S , in sorted order?
 - ◆ Is it possible to graduate with a Computer Science major from UCI in 3 years without any AP credits?

Problems and Languages



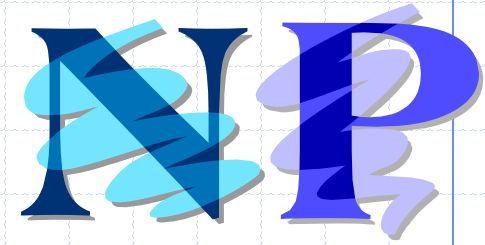
- ◆ A **language** L is a set of strings defined over some alphabet Σ
- ◆ Every decision algorithm A defines a language L
 - L is the set consisting of every string x such that A outputs “yes” on input x .
 - We say “ A **accepts** x ” in this case
 - ◆ Example:
 - ◆ If A determines whether or not a given graph G has an Euler tour, then the language L for A is all graphs with Euler tours.

The Complexity Class P



- ◆ A **complexity class** is a collection of languages
- ◆ P is the complexity class consisting of all languages that are accepted by **polynomial-time** algorithms
- ◆ For each language L in P there is a polynomial-time decision algorithm A for L.
 - If $n=|x|$, for x in L, then A runs in $p(n)$ time on input x .
 - The function $p(n)$ is some polynomial

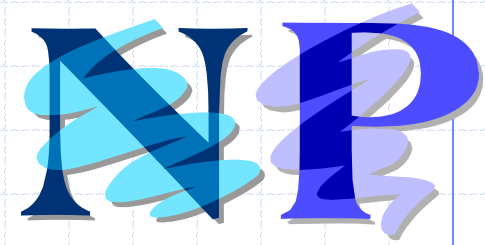
The Complexity Class NP



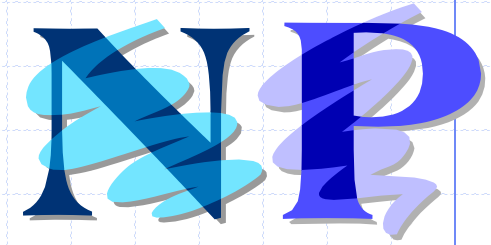
- ◆ We say that an algorithm is non-deterministic if it uses the following operation:
 - Choose(b): chooses a bit b
 - Can be used to choose an entire string y (with $|y|$ choices)
- ◆ We say that a non-deterministic algorithm A **accepts** a string x if there exists some sequence of choose operations that causes A to output “yes” on input x .
- ◆ NP is the complexity class consisting of all languages accepted by **polynomial-time non-deterministic** algorithms.

The Complexity Class NP

Alternate Definition



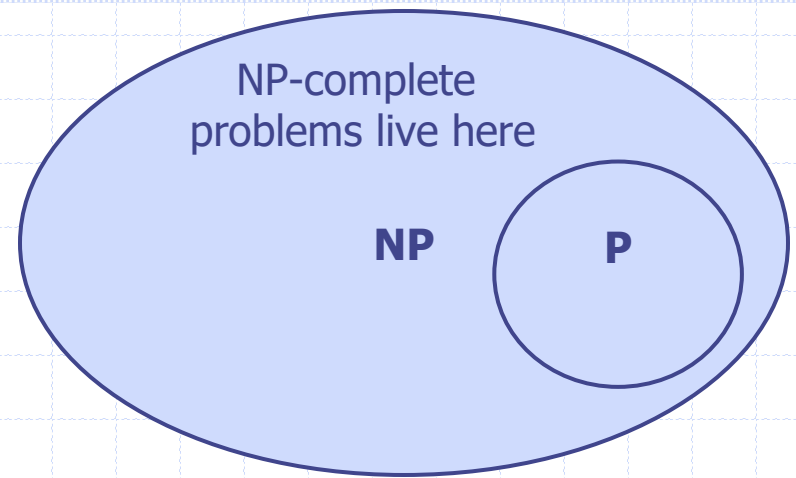
- ◆ We say that an algorithm B **verifies** the acceptance of a language L if and only if, for any x in L, there exists a certificate y such that B outputs “yes” on input (x,y) .
- ◆ NP is the complexity class consisting of all languages verified by **polynomial-time** algorithms.
- ◆ We know: P is a subset of NP.
- ◆ Major open question: $P=NP$?
- ◆ Most researchers believe that P and NP are different.



NP-Completeness

- ◆ A language M is polynomial-time **reducible** to a language L if an instance x for M can be transformed in polynomial time to an instance x' for L such that x is in M if and only if x' is in L .
 - Denote this by $M \rightarrow L$.
- ◆ A problem (language) L is **NP-hard** if every problem in NP is polynomial-time reducible to L .
- ◆ A problem (language) is **NP-complete** if it is in NP and it is NP-hard.

Some Thoughts about P and NP



- ◆ Belief: P is a proper subset of NP.
- ◆ Implication: the NP-complete problems are the hardest in NP.
 - Why: Because if we could solve an NP-complete problem in polynomial time, we could solve every problem in NP in polynomial time.
- ◆ That is, if an NP-complete problem is solvable in polynomial time, then $P=NP$.
- ◆ Since so many people have attempted without success to find polynomial-time solutions to NP-complete problems, showing your problem is NP-complete is equivalent to showing that a lot of smart people have worked on your problem and found no polynomial-time algorithm.
- ◆ If you prove or disprove the $P=NP$, you will win \$1 million.
 - See https://en.wikipedia.org/wiki/Millennium_Prize_Problems

Richard Karp



- ◆ Known for publishing a landmark paper proving 21 problems to be NP-complete.
 - One of these problems is related to the bin-packing problem we will study.
- ◆ 1985: Received the Turing Award.
- ◆ He was also the PhD advisor for UCI Professor Sandy Irani.