

Presentation for use with the textbook, *Algorithm Design and Applications*, by M. T. Goodrich and R. Tamassia, Wiley, 2015

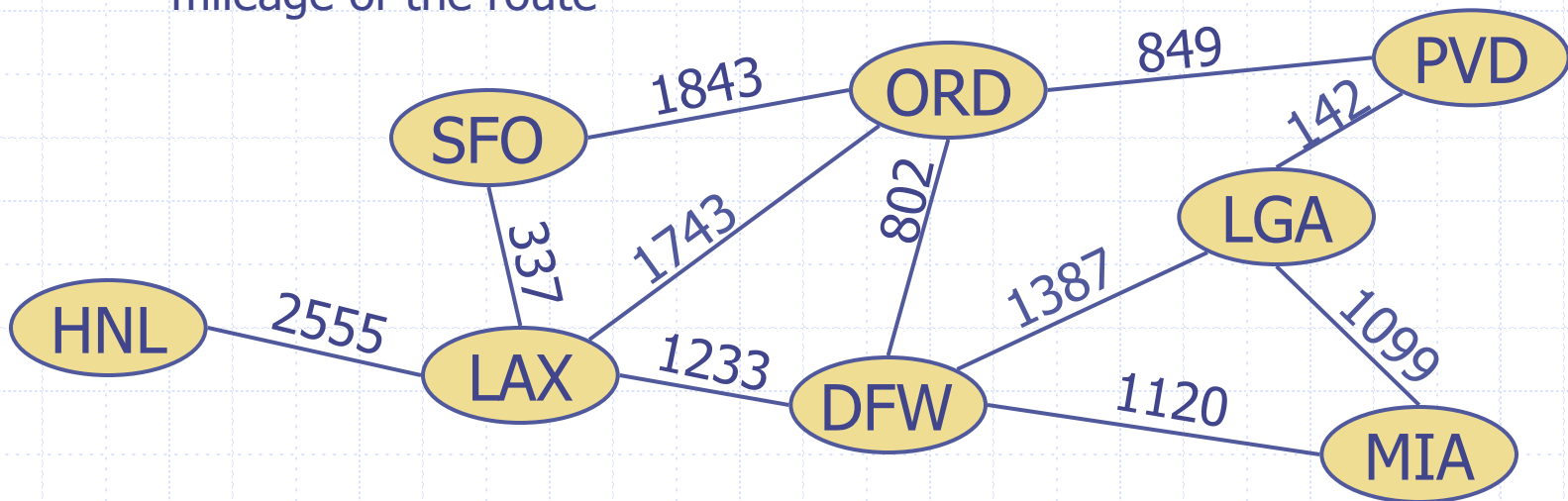
Graph Terminology and Representations



The metropolitan area of Milan, Italy at night. Astronaut photograph ISS026-E-28829, 2011. U.S. government image. NASA-JSC.

Graphs

- A graph is a pair (V, E) , where
 - V is a set of nodes, called **vertices**
 - E is a collection of pairs of vertices, called **edges**
 - Vertices and edges are positions and store elements
- Example:
 - A vertex represents an airport and stores the three-letter airport code
 - An edge represents a flight route between two airports and stores the mileage of the route



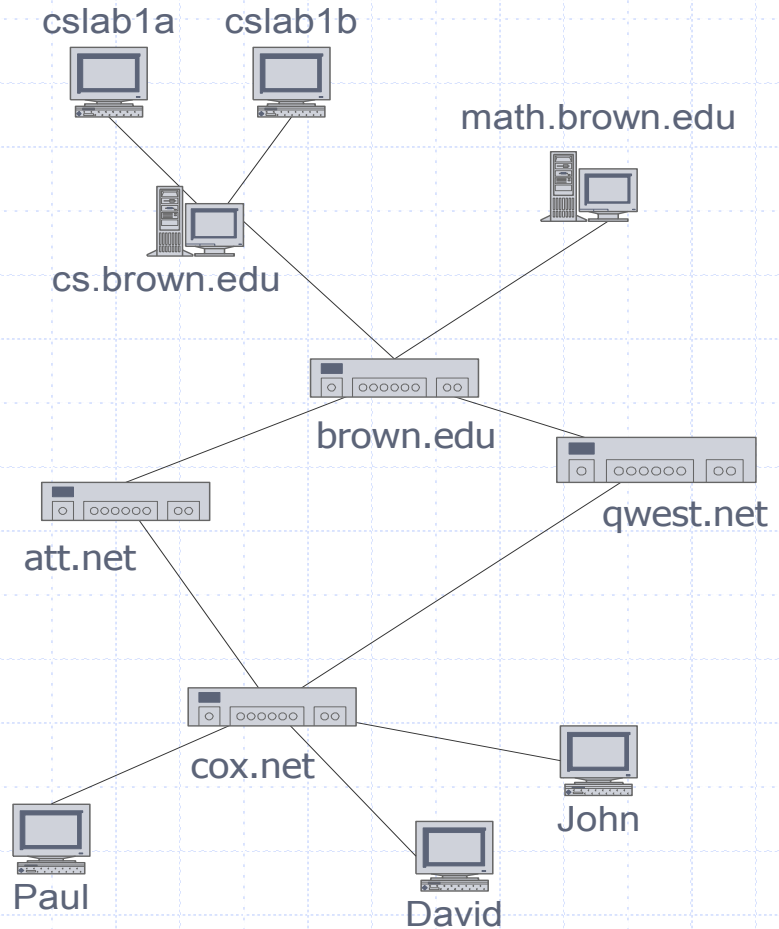
Edge Types

- Directed edge
 - ordered pair of vertices (u,v)
 - first vertex u is the origin
 - second vertex v is the destination
 - e.g., a flight
- Undirected edge
 - unordered pair of vertices (u,v)
 - e.g., a flight route
- Directed graph
 - all the edges are directed
 - e.g., route network
- Undirected graph
 - all the edges are undirected
 - e.g., flight network



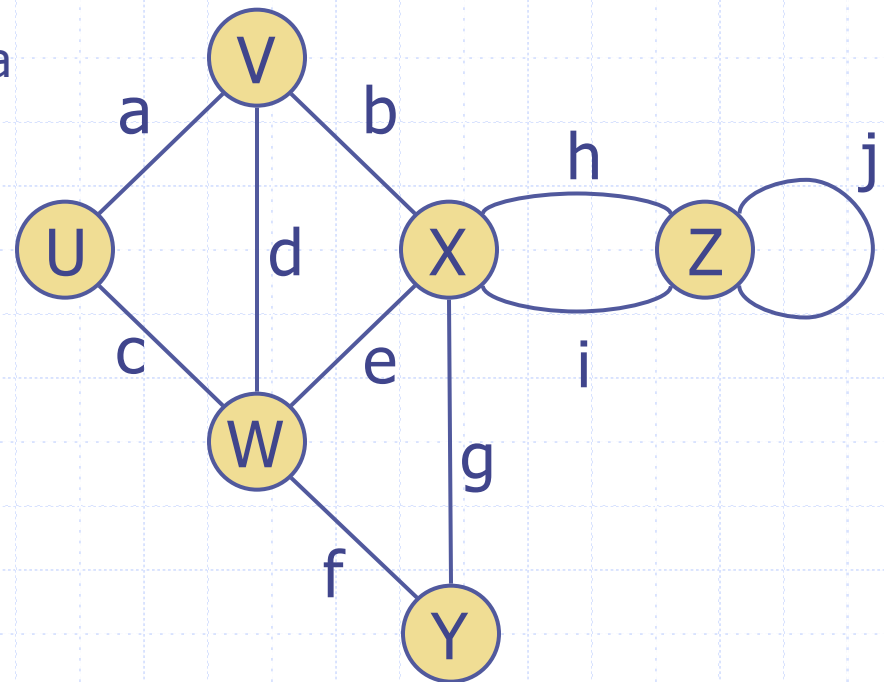
Applications

- ❑ Electronic circuits
 - Printed circuit board
 - Integrated circuit
- ❑ Transportation networks
 - Highway network
 - Flight network
- ❑ Computer networks
 - Local area network
 - Internet
 - Web
- ❑ Databases
 - Entity-relationship diagram



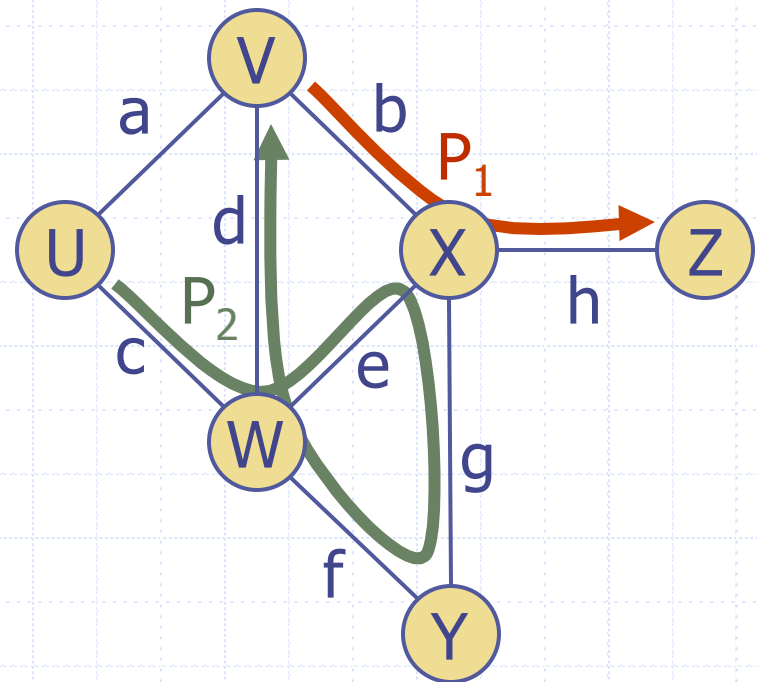
Terminology

- End vertices (or endpoints) of an edge
 - U and V are the endpoints of a
- Edges incident on a vertex
 - a, d, and b are incident on V
- Adjacent vertices
 - U and V are adjacent
- Degree of a vertex
 - X has degree 5
- Parallel edges
 - h and i are parallel edges
- Self-loop
 - j is a self-loop



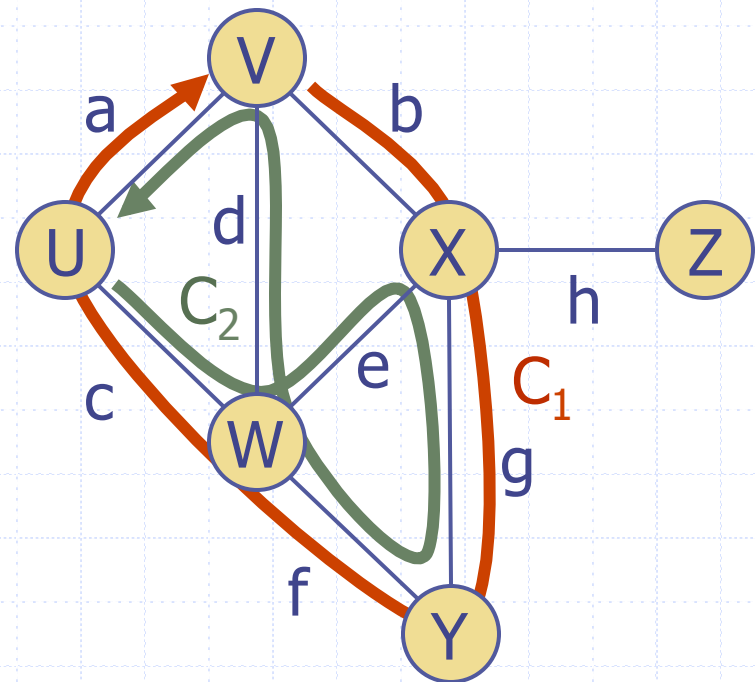
Terminology (cont.)

- Path
 - sequence of alternating vertices and edges
 - begins with a vertex
 - ends with a vertex
 - each edge is preceded and followed by its endpoints
- Simple path
 - path such that all its vertices and edges are distinct
- Examples
 - $P_1 = (V, b, X, h, Z)$ is a simple path
 - $P_2 = (U, c, W, e, X, g, Y, f, W, d, V)$ is a path that is not simple



Terminology (cont.)

- Cycle
 - circular sequence of alternating vertices and edges
 - each edge is preceded and followed by its endpoints
- Simple cycle
 - cycle such that all its vertices and edges are distinct
- Examples
 - $C_1 = (V, b, X, g, Y, f, W, c, U, a, \hookleftarrow)$ is a simple cycle
 - $C_2 = (U, c, W, e, X, g, Y, f, W, d, V, a, \hookleftarrow)$ is a cycle that is not simple



Properties

Property 1

$$\sum_v \deg(v) = 2m$$

Proof: each edge is counted twice

Property 2

In an undirected graph with no self-loops and no multiple edges

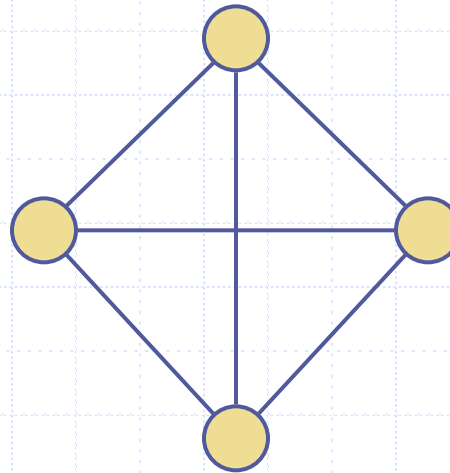
$$m \leq n(n-1)/2$$

Proof: each vertex has degree at most $(n-1)$

What is the bound for a directed graph?

Notation

n	number of vertices
m	number of edges
$\deg(v)$	degree of vertex v



Example

- $n = 4$
- $m = 6$
- $\deg(v) = 3$

Vertices and Edges

- A **graph** is a collection of **vertices** and **edges**.
- A **Vertex** is can be an abstract unlabeled object or it can be labeled (e.g., with an integer number or an airport code) or it can store other objects
- An **Edge** can likewise be an abstract unlabeled object or it can be labeled (e.g., a flight number, travel distance, cost), or it can also store other objects.

Graph Operations

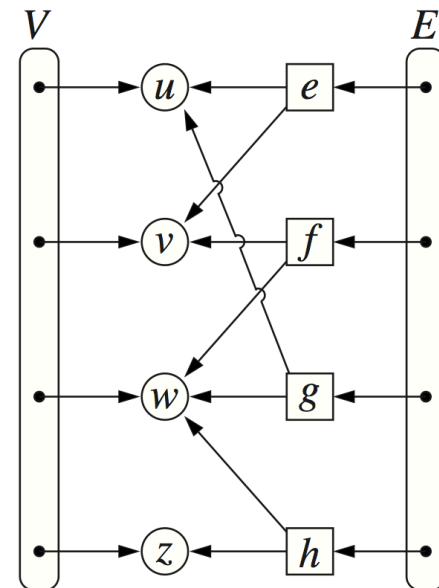
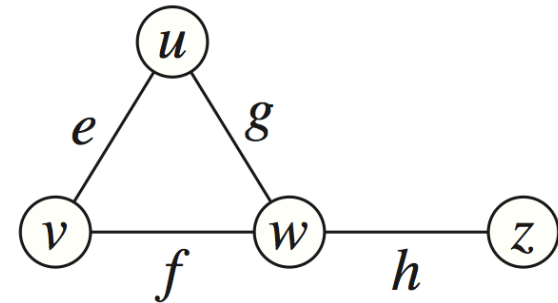
- Return the number, n , of vertices in G .
- Return the number, m , of edges in G .
- Return a set or list containing all n vertices in G .
- Return a set or list containing all m edges in G .
- Return some vertex, v , in G .
- Return the degree, $\deg(v)$, of a given vertex, v , in G .
- Return a set or list containing all the edges incident upon a given vertex, v , in G .
- Return a set or list containing all the vertices adjacent to a given vertex, v , in G .
- Return the two end vertices of an edge, e , in G ; if e is directed, indicate which vertex is the origin of e and which is the destination of e .
- Return whether two given vertices, v and w , are adjacent in G .

Graph Operations, Continued

- Indicate whether a given edge, e , is directed in G .
 - Return the in-degree of v , $\text{inDegree}(v)$.
 - Return a set or list containing all the incoming (or outgoing) edges incident upon a given vertex, v , in G .
 - Return a set or list containing all the vertices adjacent to a given vertex, v , along incoming (or outgoing) edges in G .
-
- Insert a new directed (or undirected) edge, e , between two given vertices, v and w , in G .
 - Insert a new (isolated) vertex, v , in G .
 - Remove a given edge, e , from G .
 - Remove a given vertex, v , and all its incident edges from G .

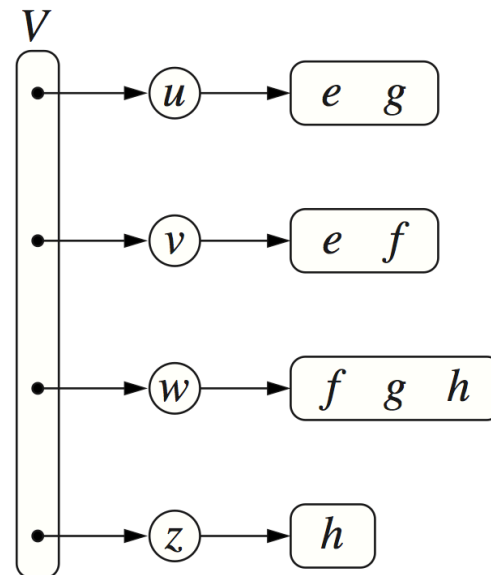
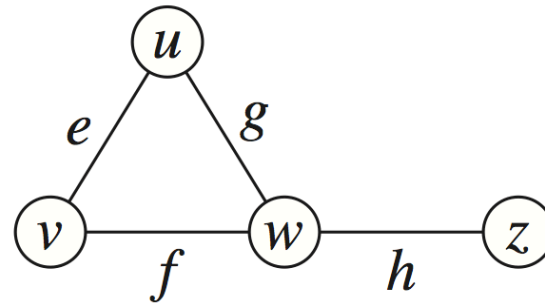
Edge List Structure

- Vertex object
 - element
 - reference to position in vertex sequence
- Edge object
 - element
 - origin vertex object
 - destination vertex object
 - reference to position in edge sequence
- Vertex sequence
 - sequence of vertex objects
- Edge sequence
 - sequence of edge objects



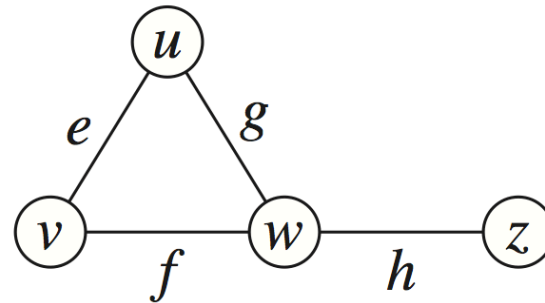
Adjacency List Structure

- Incidence sequence for each vertex
 - sequence of references to edge objects of incident edges
- Augmented edge objects
 - references to associated positions in incidence sequences of end vertices



Adjacency Matrix Structure

- Edge list structure
- Augmented vertex objects
 - Integer key (index) associated with vertex
- 2D-array adjacency array
 - Reference to edge object for adjacent vertices
 - Null for non adjacent vertices
- The “old fashioned” version just has 0 for no edge and 1 for edge



		0	1	2	3	
u	\longrightarrow	0		e	g	
v	\longrightarrow	1	e		f	
w	\longrightarrow	2	g	f		h
z	\longrightarrow	3			h	

Performance

(All bounds are big-oh running times, except for “Space”)

<ul style="list-style-type: none">▪ n vertices, m edges▪ no parallel edges▪ no self-loops	Edge List	Adjacency List	Adjacency Matrix
Space	$n + m$	$n + m$	n^2
incidentEdges(v)	m	$\deg(v)$	n
areAdjacent (v, w)	m	$\min(\deg(v), \deg(w))$	1
insertVertex(o)	1	1	n^2
insertEdge(v, w, o)	1	1	1
removeVertex(v)	m	$\deg(v)$	n^2
removeEdge(e)	1	1	1