

ICS 6B Boolean Algebra & Logic

Lecture Notes for Summer Quarter, 2008

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Set 3 – Ch. 1.4, 2.1, 2.2

Announcements

- Quiz schedule online*
 - Will allow you to drop 1 quiz
 - Next Quiz is on Thursday
 - * Subject to change
- Homework is online

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2

Today's Lecture

- Chapter 1 (Section 1.4)
 - Nested Quantifiers(1.4)
- Chapter 2 (Sections 2.1 & 2.2)
 - Sets(2.1)
 - Set Operations (2.2)

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3

Chapter 1: Section 1.4

Nested Quantifiers

What are Nested quantifiers?

If one quantifier is within the scope of the other.

- Eg.

U:R

$$\forall x \exists y(x + y = 0)$$

This is the same as $\forall x Q(x)$,
where $Q(x)$ is $\exists y P(x, y)$,
where $P(x, y)$ is $(x + y = 0)$

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5

Translating to English

Translate:

U: R

$$\forall x \forall y((x > 0) \wedge (y < 0) \rightarrow (xy < 0))$$

"For every real number x and every real number y, if x > 0 and y < 0, then xy < 0"

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6

Switching order

- If the **quantifiers** are the **same** switching order **doesn't matter**
 - (ie. All \forall 's or all \exists 's)
 - $\forall x \forall y P(x,y) \equiv \forall y \forall x P(x,y)$
 - $\exists x \exists y P(x,y) \equiv \exists y \exists x P(x,y)$
- If the **quantifiers** are **different** then **order matters**
 - $\forall x \exists y P(x,y) \not\equiv \exists y \forall x P(x,y)$

NOT Equivalent

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Thinking of Quantification as Loops

To **prove** or **disprove** nested quantifications
 → think in terms of **nested loops**

$\forall x \forall y P(x,y)$
 Loop through **x** values
 For each **x** value loop through the **y** values
 If we find that $P(x,y)$ is true for **all values of y**
 for **every x**,
 then $\forall x \forall y P(x,y)$ is **True**
 If we find **one y** for any **x** such that $P(x,y)$ is False
 then $\forall x \forall y P(x,y)$ is **False**

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Thinking of Quant. as Loops (2)

$\forall x \exists y P(x,y)$
 Loop through **x** values
 For each **x** value loop through the **y** values
 If we find **one y** for **each x** such that $P(x,y)$ is true
 then $\forall x \exists y P(x,y)$ is **True**
 If for any **one x** we can't find a **y** such that $P(x,y)$ is true
 then $\forall x \exists y P(x,y)$ is **False**

$\exists x \forall y P(x,y)$
 Loop through **x** values
 For each **x** value loop through the **y** values
 If we find an **x** such $P(x,y)$ is true for **all y**'s
 then $\exists x \forall y P(x,y)$ is **True**
 If we can't find such an **x**
 then $\exists x \forall y P(x,y)$ is **False**

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Thinking of Quant. as Loops (3)

$\exists x \exists y P(x,y)$
 Loop through **x** values
 For each **x** value loop through the **y** values
 If we find **one y** for **one x** such that $P(x,y)$ is true
 then $\exists x \exists y P(x,y)$ is **True**
 If we can't find **one x** and **one y** such that $P(x,y)$ is true
 then $\exists x \exists y P(x,y)$ is **False**

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Quantification of Two Variables

Statement	When True?	When False?
$\forall x \forall y$	$P(x,y)$ is true for every x, y pair	There is an x, y pair for which $P(x,y)$ is false
$\forall x \exists y$	For every x , there is at least one y for which $P(x,y)$ is true	There is an x such that $P(x,y)$ is false for every y
$\exists x \forall y$	There is an x for which $P(x,y)$ is true for every y	For every x there is at least one y for which $P(x,y)$ is false
$\exists x \exists y$	There is at least one x, y pair for which $P(x,y)$ is true	$P(x,y)$ is false for every x, y pair

Note: These are not equivalent.

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Translating from English to Nested Quantifiers

"The product of two positive numbers is positive."

$U: \mathbb{R}$
 $\forall x \forall y [(x > 0 \wedge y > 0) \rightarrow (xy > 0)]$

Is there an easier way?
 $U: \mathbb{R} > 0$
 $\forall x \forall y (xy > 0)$

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Translating from English to Nested Quantifiers (2)

"Given a number, there is a number greater than it."

In other words:

"For every number x we choose, there is a number y such that $y > x$."

$U: \mathbb{R}$

$\forall x \exists y (y > x)$

Homework for Section 1.4

- 3(b,f)
- 5(b,f)
- 9(b,d,h,j)
- 11(b,f,h)
- 15(b,d,f)
- Feel free to do more if you need the practice

Chapter 2: Section 2.1

Sets

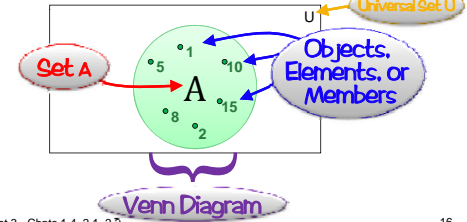
What is a Set?

○ **Set:** An unordered *collection of objects*

- The point \rightarrow to group objects together
- Often objects have some similar properties

○ **Objects:** *elements* or *members* of the set

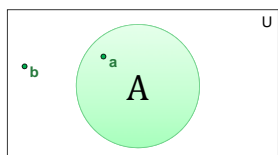
- A set is said to *contain* its element



Some Notations

$a \in A \rightarrow a$ is an **element** of the set A

$b \notin A \rightarrow a$ is **not an element** of the set A



Describing sets

○ $\{ \}$ denote all the elements in the set

Eg.

$V = \{a, e, i, o, u\}$



Sets can also have unrelated objects

$O = \{26, \text{Paul}, \text{Pot}, a\}$

... - ellipses denote a pattern

$I = \{2, 4, 6, \dots, 98\}$

Set Builders (describing sets)

- You can also use **set builders** so that you don't have to name every element
 - Just state the properties

$I = \{x \mid x \text{ is a positive even integer less than } 100\}$

Or

$I = \{x \in \mathbb{Z}^+ \mid x \text{ is even and } x < 100\}$

All integers *All positive*

- You can also use Predicates

$I = \{x \mid P(x)\}$

I contains all elements from U which make P true

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Common Universal Sets

$\mathbb{N} = \{0, 1, 2, 3, \dots\}$, the set of **natural numbers**

$\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$, the set of **integers**

$\mathbb{Z}^+ = \{1, 2, 3, \dots\}$, the set of **positive integers**

$\mathbb{Q} = \{p/q \mid p \in \mathbb{Z}, q \in \mathbb{Z}, \text{ and } q \neq 0\}$, the set of **rational numbers**

\mathbb{R} , the set of **real numbers**

Note: Sometimes 0 is not considered a part of the set of natural numbers.

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Special Sets

- The **empty set, void set, or null set**
 - A set with no elements
 - Notation: \emptyset or $\{\}$
 - The assertion $x \in \emptyset$ is always false
- The **singleton set**
 - A set with one element

Is this the empty set? $\{\emptyset\}$

No! It is the singleton set with the empty set as its element

Think of the empty set as an empty folder

Think of this $\{\emptyset\}$ as a folder with only an empty folder in it

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Subsets

A is a **subset** of B iff every element of A is also an element of B .

- In other words A is a **subset** of B iff

$\forall x (x \in A \rightarrow x \in B)$.

- Notation

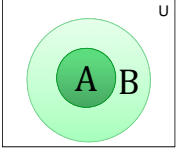
$A \subseteq B$

Example:

$A = \{2, 4, 6\}$

$B = \{1, 2, 3, 4, 5, 6, 7, 8\}$

$A \subseteq B$



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Proper Subsets

- A **proper subset** is a subset in which $A \neq B$
- In other words A is a proper subset of B iff
 - $\forall x (x \in A \rightarrow x \in B) \wedge \exists x (x \in B \wedge x \notin A)$
 - Notation: $A \subset B$

Eg.

$A = \{1, 2, 3\}$

$B = \{0, 1, 2, 3, 4, 5, 6\}$

$A \subset B$ because B has more elements than A

If $A = \{0, 1, 2, 3, 4, 5, 6\}$

then $A \subseteq B$, but $A \not\subset B$

What if $A = \{1, 2, 9\}$?

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Theorem 1

For every set S ,

(1) $\emptyset \subseteq S$ and (2) $S \subseteq S$

Proof

Let S be a set

To show that $\emptyset \subseteq S$ we must show that

$\forall x (x \in \emptyset \rightarrow x \in S)$ is **true**

Because \emptyset contains no elements, it follows that $x \in \emptyset$ is **always false**.

It follows that $x \in \emptyset \rightarrow x \in S$ is **always true**, because its hypothesis is **always false** and a **conditional statement with a false hypothesis is always true**.

Therefore $\forall x (x \in \emptyset \rightarrow x \in S)$ is **true**

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Equal Sets

Two sets are **equal** iff they have **the same elements**

- In other words A and B are **equal** iff

$$\forall x (x \in A \leftrightarrow x \in B).$$

- Notation

$$A=B$$

Eg.

$$A=\{x, y, z\}$$

$$B=\{z, x, y\}$$

$$C=\{z,z,z,z,z,y,y,y,y,y,y,y,x\}$$

$$A=B=C$$

25

Equal Sets (2)

We can prove 2 sets (A & B) are equal if we can show:

$$A \subseteq B \text{ and } B \subseteq A$$

Remember :

$A \subseteq B$ is the same as $\forall x (x \in A \rightarrow x \in B)$ and

$B \subseteq A$ is the same as $\forall x (x \in B \rightarrow x \in A)$

Which is the same as saying $\forall x (x \in A \leftrightarrow x \in B)$

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26

Sets as Elements/Members

- Sets may have other sets as elements

Like with the empty set ...

$$A=\{\emptyset, \{a\}, \{b\}, \{a,b\}\}$$

Is $a \in A$?

No!

Is $\{a\} \in A$?

Yes!

Is $B=\{x \mid x \text{ is a subset of the set } \{a,b\}\}$ equivalent to A?

Yes!

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27

Cardinality, Finite & Infinite Sets

Let S be a set.

If there is **exactly n** distinct elements in S where n is a nonnegative integer, we say that S is a **finite set** and that N is the **cardinality** of S.

- Notation for cardinality: $|S|$

Eg.

Let V be the set of vowels in the alphabet

$$|V| = 5$$

A set is **infinite** if it is **not finite**.

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28

The Power Set

Given a set S, the **power set** is the **set of all subsets** of S.

- Notation

$$P(S)$$

Eg.

$$\text{Let } S=\{a,b,c\}$$

What is P(S)?

$$P(S)=\{\emptyset, \{a\}, \{b\}, \{c\}, \{a,b\}, \{a,c\}, \{b,c\}, \{a,b,c\}\}$$

How many elements does P(S) have if |S| is 6?

2⁶ or 64 elements

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29

Ordered n-tuples

- Sometimes order matters

- Sets are unordered

- We use **ordered n-tuples**

The **ordered n-tuple** (a_1, a_2, \dots, a_n) is the **ordered collection** that has a_1 as its first element, a_2 as its 2nd element, ..., and a_n as its nth element.

- They are equal iff each corresponding pair of elements is equal.

- $\{a_1, a_2, \dots, a_n\} = \{b_1, b_2, \dots, b_n\}$ iff $a_i = b_i$

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30

Cartesian Product

Let A & B be sets.
The **Cartesian product** of A and B is the set of all ordered pairs (a,b) where $a \in A$ and $b \in B$.

- **Notation**

$$A \times B$$

Eg.

Let $A = \{1,2\}$ and $B = \{a,b,c\}$

What is $A \times B$

- $A \times B = \{(1,a), (1,b), (1,c), (2,a), (2,b), (2,c)\}$

Cartesian Product (2)

- **Some things to note:**

$A \times B \neq B \times A$ **unless**

- $A = \emptyset$ or $B = \emptyset$ (thus $A \times B = \emptyset$) or
- $A = B$

Relations

A subset R of the Cartesian product $A \times B$ is called a **relation** from set A to set B.

The elements of R are ordered pairs where the 1st element belongs to A and the 2nd to B.

Eg.

Let $A = \{1,2\}$ and $B = \{a,b,c\}$

$R = \{(1,a), (1,b), (1,c), (2,a), (2,b), (2,c)\}$

R is a relations from Set A to Set B

Cartesian Product: More than 2 sets

The **Cartesian product** of the sets A_1, A_2, \dots, A_n is the set of ordered n-tuples (a_1, a_2, \dots, a_n) , where a_i belongs to A_i for $i=1,2,\dots,n$

- **Notated**

$$A_1 \times A_2 \times \dots \times A_n$$

Eg.

Let $A = \{1,2\}$, $B = \{a,b,c\}$, $C = \{y,z\}$

What is $A \times B \times C$

- $A \times B \times C = \{(1,a,y), (1,a,z), (1,b,y), (1,b,z), (1,c,y), (1,c,z), (2,a,y), (2,a,z), (2,b,y), (2,b,z), (2,c,y), (2,c,z)\}$

Thus $A \times B \times C$ is all possible ordered tuples (a,b,c) where $a \in A$, $b \in B$, and $c \in C$.

Homework for Section 2.1

- 1, 3, 5, 7, 13, 17, 23, 25, 27, 31

Chapter 2: Section 2.2

Set Operations

Union

- Let A & B be sets.

The **union** of the sets A & B is the set that contains those elements **in either A or B or both**.

- Notation: $A \cup B$

- In other words: $\{x|x \in A \vee x \in B\}$

Eg

$$A = \{1, 2, 3\}$$

$$B = \{3, 4, 5\}$$

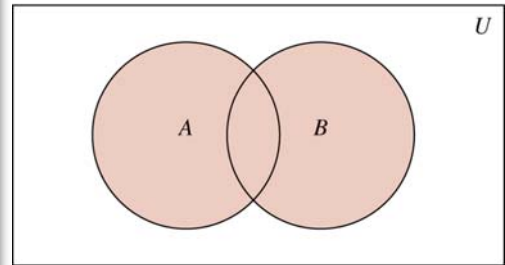
$$A \cup B = \{1, 2, 3, 4, 5\}$$

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37

Union – Venn Diagram

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$A \cup B$ is shaded.

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38

Intersection

- Let A & B be sets.

The **intersection** of the sets A & B is the set that contains those elements **in both A and B**.

- Notation: $A \cap B$

- In other words: $\{x|x \in A \wedge x \in B\}$

Eg

$$A = \{1, 2, 3\}$$

$$B = \{3, 4, 5\}$$

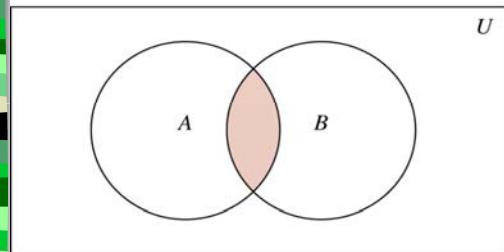
$$A \cap B = \{3\}$$

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39

Intersection – Venn Diagram

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$A \cap B$ is shaded.

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40

Disjoint

- Let A & B be sets.

A & B are **disjoint** if the **intersection** of the sets A & B is **the empty set**.

Eg

$$A = \{1, 2, 3\}$$

$$B = \{4, 5, 6\}$$

$$A \cap B = \emptyset$$

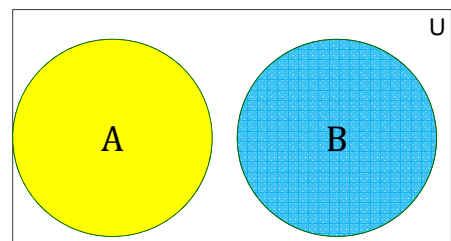
A & B are Disjoint

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41

Disjoint – Venn Diagram

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42

Difference / Complement

- Let A & B be sets.
- The **difference** in A & B is the set containing those elements that **are in A , but not in B** .
- AKA the **complement** of B with respect to A .
- Notation: $A - B$
 - In other words: $\{x | x \in A \wedge x \notin B\}$

Eg

$A = \{1, 2, 3\}$
 $B = \{3, 4, 5, 6\}$
 $A - B = \{1, 2\}$

43

Difference / Complement – Venn Diagram

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$A - B$ is shaded.

44

(Absolute) Complement

- Once the universal set U can be specified, the complement can be defined
- Let U be the universal set.
- The **(absolute) complement** of set A , is the **complement of A with respect to U** .
- In other words, the **(absolute) complement** of set A is $U - A$
- Notation: \bar{A} or A^c
 - In other words: $\{x | x \notin A\}$ or $\{x | \neg(x \in A)\}$

Eg

$U = \{2, 4, 6, 8, 10\}$
 $A = \{2, 4, 6\}$
 $\bar{A} = \{8, 10\}$

45

Complement

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\bar{A} is shaded.

46

Symmetric Difference

- Let A & B be sets.
- The **symmetric difference** in A & B is the set containing those elements that **are in A , but not in B and the elements in B , that are not in A** .
- Notation: $A \oplus B$
 - In other words: $(A - B) \cup (B - A)$

Eg

$A = \{1, 2, 3\}$
 $B = \{3, 4, 5, 6\}$
 $A - B = \{1, 2\}$
 $B - A = \{4, 5, 6\}$
 $A \oplus B = \{1, 2, 4, 5, 6\}$

47

Examples

$U = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$
 $A = \{1, 2, 3, 4, 5\}$
 $B = \{4, 5, 6, 7, 8\}$
 $A \cup B = \{1, 2, 3, 4, 5, 6, 7, 8\}$
 $A \cap B = \{4, 5\}$
 $A^c = \{0, 6, 7, 8, 9, 10\}$
 $B^c = \{0, 1, 2, 3, 9, 10\}$
 $A - B = \{1, 2, 3\}$
 $B - A = \{6, 7, 8\}$
 $A \oplus B = \{1, 2, 3, 6, 7, 8\}$

48

TABLE 1 Set Identities.	
Identity	Name
$A \cup \emptyset = A$ $A \cap U = A$	Identity laws
$A \cup U = U$ $A \cap \emptyset = \emptyset$	Domination laws
$A \cup A = A$ $A \cap A = A$	Idempotent laws
$\overline{(\overline{A})} = A$	Complementation law
$A \cup B = B \cup A$ $A \cap B = B \cap A$	Commutative laws
$A \cup (B \cup C) = (A \cup B) \cup C$ $A \cap (B \cap C) = (A \cap B) \cap C$	Associative laws

More Set Identities

$A \cup (B \cup C) = (A \cup B) \cup C$ $A \cap (B \cap C) = (A \cap B) \cap C$	Associative laws
$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$ $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$	Distributive laws
$\overline{A \cup B} = \overline{A} \cap \overline{B}$ $\overline{A \cap B} = \overline{A} \cup \overline{B}$	De Morgan's laws
$A \cup (A \cap B) = A$ $A \cap (A \cup B) = A$	Absorption laws
$A \cup \overline{A} = U$ $A \cap \overline{A} = \emptyset$	Complement laws

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