Knowledge Representation using First-Order Logic (Part III)

This lecture: R&N Chapters 8, 9

Next lecture: Chapter 13; Chapter 14.1-14.2

(Please read lecture topic material before and after each lecture on that topic)

Outline

- Review: $KB \mid = S$ is equivalent to $\mid = (KB \Rightarrow S)$
 - So what does {} |= S mean?
- Review: Follows, Entails, Derives
 - Follows: "Is it the case?"
 - Entails: "Is it true?"
 - Derives: "Is it provable?"
- Review: FOL syntax
- Finish FOL Semantics, FOL examples
- Inference in FOL

Using FOL

• We want to TELL things to the KB, e.g.

TELL(KB,
$$\forall x, King(x) \Rightarrow Person(x)$$
)
TELL(KB, $King(John)$)

These sentences are assertions

• We also want to ASK things to the KB, $ASK(KB, \exists x, Person(x))$

these are queries or goals

The KB should return the list of x's for which Person(x) is true: {x/John,x/Richard,...}

FOL Version of Wumpus World

- Typical percept sentence: Percept([Stench,Breeze,Glitter,None,None],5)
- Actions: Turn(Right), Turn(Left), Forward, Shoot, Grab, Release, Climb
- To determine best action, construct query:
 ∀ a BestAction(a,5)
- ASK solves this and returns {a/Grab}
 - And TELL about the action.

Knowledge Base for Wumpus World

- Perception
 - ∀s,b,g,x,y,t Percept([s,Breeze,g,x,y],t) ⇒ Breeze(t)
 - $\forall s,b,x,y,t \text{ Percept}([s,b,Glitter,x,y],t) \Rightarrow Glitter(t)$
- Reflex action
 - ∀t Glitter(t) ⇒ BestAction(Grab,t)
- Reflex action with internal state
 - – ∀t Glitter(t) ∧¬Holding(Gold,t) ⇒ BestAction(Grab,t)

Holding(Gold,t) can not be observed: keep track of change.

Deducing hidden properties

Environment definition:

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\forall x,y,a,b \; Adjacent([x,y],[a,b]) \Leftrightarrow
[a,b] \in \{[x+1,y], [x-1,y],[x,y+1],[x,y-1]\}
```

Properties of locations:

$$\forall$$
s,t At (Agent,s,t) \land Breeze(t) \Rightarrow Breezy(s)

Squares are breezy near a pit:

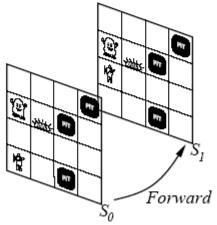
- Diagnostic rule---infer cause from effect
 ∀s Breezy(s) ⇔ ∃ r Adjacent(r,s) ∧ Pit(r)
- Causal rule---infer effect from cause (model based reasoning)
 ∀r Pit(r) ⇒ [∀s Adjacent(r,s) ⇒ Breezy(s)]

Keeping track of change

Facts hold in situations, rather than eternally E.g., Holding(Gold, Now) rather than just Holding(Gold)

Situation calculus is one way to represent change in FOL: Adds a situation argument to each non-eternal predicate E.g., Now in Holding(Gold, Now) denotes a situation

Situations are connected by the Result function Result(a,s) is the situation that results from doing a in s



Describing actions I

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"Effect" axiom—describe changes due to action \forall s \; AtGold(s) \Rightarrow Holding(Gold, Result(Grab, s))
```

"Frame" axiom—describe non-changes due to action $\forall s \; HaveArrow(s) \Rightarrow HaveArrow(Result(Grab, s))$

Frame problem: find an elegant way to handle non-change

- (a) representation—avoid frame axioms
- (b) inference—avoid repeated "copy-overs" to keep track of state

Qualification problem: true descriptions of real actions require endless caveats—what if gold is slippery or nailed down or . . .

Ramification problem: real actions have many secondary consequences—what about the dust on the gold, wear and tear on gloves, . . .

Describing actions II

Successor-state axioms solve the representational frame problem

Each axiom is "about" a predicate (not an action per se):

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P true afterwards \Leftrightarrow [an action made P true \lor P true already and no action made P false]
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For holding the gold:

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 \forall \, a,s \; \, Holding(Gold,Result(a,s)) \; \Leftrightarrow \\ [(a = Grab \land AtGold(s)) \\ \lor (Holding(Gold,s) \land a \neq Release)]
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Set Theory in First-Order Logic

Can we define set theory using FOL?

- individual sets, union, intersection, etc

Answer is yes.

Basics:

- empty set = constant = { }
- unary predicate Set(), true for sets
- binary predicates:

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x \in S (true if x is a member of the set s)

S_1 \subseteq S_2 (true if s1 is a subset of s2)
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- binary functions:

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intersection s_1 \cap s_2, union s_1 \cup s_2, adjoining \{x|s\}
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A Possible Set of FOL Axioms for Set Theory

The only sets are the empty set and sets made by adjoining an element to a set

$$\forall s \ \text{Set}(s) \Leftrightarrow (s = \{\}) \lor (\exists x, s_2 \ \text{Set}(s_2) \land s = \{x | s_2\})$$

The empty set has no elements adjoined to it

$$\neg \exists x, s \{x | s\} = \{\}$$

Adjoining an element already in the set has no effect

$$\forall x, s \ x \in s \Leftrightarrow s = \{x | s\}$$

The only elements of a set are those that were adjoined into it. Expressed recursively:

$$\forall x, s \quad x \in s \Leftrightarrow [\exists y, s_2 (s = \{y | s_2\} \land (x = y \lor x \in s_2))]$$

A Possible Set of FOL Axioms for Set Theory

A set is a subset of another set iff all the first set's members are members of the 2nd set

$$\forall s_1, s_2 \ s_1 \subseteq s_2 \Leftrightarrow (\forall x \ x \in s_1 \Rightarrow x \in s_2)$$

Two sets are equal iff each is a subset of the other

$$\forall s_1, s_2 \ (s_1 = s_2) \Leftrightarrow (s_1 \subseteq s_2 \land s_2 \subseteq s_1)$$

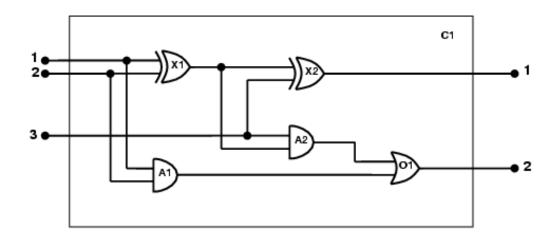
An object is in the intersection of 2 sets only if a member of both $\forall x, s_1, s_2 \ x \in (s_1 \cap s_2) \Leftrightarrow (x \in s_1 \land x \in s_2)$

An object is in the union of 2 sets only if a member of either $\forall x, s_1, s_2 \ x \in (s_1 \cup s_2) \Leftrightarrow (x \in s_1 \lor x \in s_2)$

Knowledge engineering in FOL

- 1. Identify the task
- 2. Assemble the relevant knowledge
- 3. Decide on a vocabulary of predicates, functions, and constants
- 4. Encode general knowledge about the domain
- 5. Encode a description of the specific problem instance
- 6. Pose queries to the inference procedure and get answers
- 7. Debug the knowledge base

One-bit full adder



Possible queries:

- does the circuit function properly?
- what gates are connected to the first input terminal?
- what would happen if one of the gates is broken? and so on

- 1. Identify the task
 - Does the circuit actually add properly?
- 2. Assemble the relevant knowledge
 - Composed of wires and gates; Types of gates (AND, OR, XOR, NOT)

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- Irrelevant: size, shape, color, cost of gates

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- 3. Decide on a vocabulary
 - Alternatives:

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Type(X_1) = XOR (function)
Type(X_1, XOR) (binary predicate)
XOR(X_1)
(unary predicate)
```

- 4. Encode general knowledge of the domain
 - $\forall t_1, t_2 \text{ Connected}(t_1, t_2) \Rightarrow \text{Signal}(t_1) = \text{Signal}(t_2)$
 - $\forall t \ Signal(t) = 1 \lor Signal(t) = 0$
 - 1 \neq 0
 - $\forall t_1, t_2 \text{ Connected}(t_1, t_2) \Rightarrow \text{Connected}(t_2, t_1)$
 - $\forall g \text{ Type}(g) = OR \Rightarrow \text{Signal}(Out(1,g)) = 1 \Leftrightarrow \exists n \text{ Signal}(In(n,g)) = 1$
 - $\forall g \text{ Type}(g) = \text{AND} \Rightarrow \text{Signal}(\text{Out}(1,g)) = 0 \Leftrightarrow \exists n \text{ Signal}(\text{In}(n,g)) = 0$
 - ∀g Type(g) = XOR ⇒ Signal(Out(1,g)) = 1 ⇔ Signal(In(1,g)) ≠ Signal(In(2,g))
 - $\forall g \text{ Type}(g) = \text{NOT} \Rightarrow \text{Signal}(\text{Out}(1,g)) \neq \text{Signal}(\text{In}(1,g))$

5. Encode the specific problem instance

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Type(X_1) = XOR Type(X_2) = XOR

Type(A_1) = AND Type(A_2) = AND

Type(O_1) = OR
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Connected(Out(1, X_1),In(1, X_2))	Connected($In(1,C_1),In(1,X_1)$)
Connected(Out(1, X_1),In(2, A_2))	Connected($In(1,C_1),In(1,A_1)$)
Connected(Out(1, A_2),In(1, O_1))	Connected($In(2,C_1),In(2,X_1)$)
Connected(Out(1, A_1),In(2, O_1))	Connected($In(2,C_1),In(2,A_1)$)
Connected(Out(1, X_2),Out(1, C_1))	Connected($In(3,C_1),In(2,X_2)$)
Connected(Out(1,O ₁),Out(2,C ₁))	Connected($In(3,C_1),In(1,A_2)$)

6. Pose queries to the inference procedure
What are the possible sets of values of all the terminals for the adder circuit?

$$\exists i_1, i_2, i_3, o_1, o_2 \ Signal(In(1, C_1)) = i_1 \land Signal(In(2, C_1)) = i_2 \land Signal(In(3, C_1)) = i_3 \land Signal(Out(1, C_1)) = o_1 \land Signal(Out(2, C_1)) = o_2$$

Debug the knowledge base
 May have omitted assertions like 1 ≠ 0

Syntactic Ambiguity

- FOPC provides many ways to represent the same thing.
- E.g., "Ball-5 is red."
 - HasColor(Ball-5, Red)
 - Ball-5 and Red are objects related by HasColor.
 - Red(Ball-5)
 - Red is a unary predicate applied to the Ball-5 object.
 - HasProperty(Ball-5, Color, Red)
 - Ball-5, Color, and Red are objects related by HasProperty.
 - ColorOf(Ball-5) = Red
 - Ball-5 and Red are objects, and ColorOf() is a function.
 - HasColor(Ball-5(), Red())
 - Ball-5() and Red() are functions of zero arguments that both return an object, which objects are related by HasColor.
 - **—** ...
- This can GREATLY confuse a pattern-matching reasoner.
 - Especially if multiple people collaborate to build the KB, and they all have different representational conventions.

Summary

- First-order logic:
 - Much more expressive than propositional logic
 - Allows objects and relations as semantic primitives
 - Universal and existential quantifiers
 - syntax: constants, functions, predicates, equality, quantifiers

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- Knowledge engineering using FOL
 - Capturing domain knowledge in logical form
- Inference and reasoning in FOL
 - Next lecture
- Required Reading:
 - All of Chapter 8
 - Next lecture: Chapter 9