# First-Order Logic C: Knowledge Engineering

### CS171, Fall Quarter, 2019 Introduction to Artificial Intelligence Prof. Richard Lathrop



Read Beforehand: R&N 8, 9.1-9.2, 9.5.1-9.5.5



# Outline

- Review --- Syntactic Ambiguity
- Using FOL
  - Tell, Ask
- Example: Wumpus world
- Deducing Hidden Properties
  - Keeping track of change
  - Describing the results of Actions
- Set Theory in First-Order Logic
- Knowledge engineering in FOL
- The electronic circuits domain

### You will be expected to know

Seven steps of Knowledge Engineering (R&N section 8.4.1)

 Given a simple Knowledge Engineering problem, produce a simple FOL Knowledge Base that solves the problem

# Review --- Syntactic Ambiguity

- FOPC provides many ways to represent the same thing.
- E.g., "Ball-5 is red."
  - HasColor(Ball-5, Red)
    - Ball-5 and Red are objects related by HasColor.
  - Red(Ball-5)
    - Red is a unary predicate applied to the Ball-5 object.
  - HasProperty(Ball-5, Color, Red)
    - Ball-5, Color, and Red are objects related by HasProperty.
  - ColorOf(Ball-5) = Red
    - Ball-5 and Red are objects, and ColorOf() is a function.
  - HasColor(Ball-5(), Red())
    - Ball-5() and Red() are functions of zero arguments that both return an object, which objects are related by HasColor.
  - ...
- This can GREATLY confuse a pattern-matching reasoner.
  - Especially if multiple people collaborate to build the KB, and they all have different representational conventions.

### Review --- Syntactic Ambiguity ---Partial Solution

- FOL can be TOO expressive, can offer TOO MANY choices
- Likely confusion, especially for **teams** of Knowledge Engineers
- Different team members can make different representation choices
  - E.g., represent "Ball43 is Red." as:
    - a predicate (= verb)? E.g., "Red(Ball43)"?
    - an object (= noun)? E.g., "Red = Color(Ball43))"?
    - a property (= adjective)? E.g., "HasProperty(Ball43, Red)"?
- PARTIAL SOLUTION:
  - An upon-agreed **ontology** that settles these questions
  - Ontology = what exists in the world & how it is represented
  - The Knowledge Engineering teams agrees upon an ontology BEFORE they begin encoding knowledge

## Using FOL

 We want to TELL things to the KB, e.g. TELL(KB, ∀ x King(x) ⇒ PersonX))
 TELL(KB, King(John))

These sentences are assertions

 We also want to ASK things to the KB, ASK(KB, ∃ x Person(x))

these are queries or goals

The KB should return the list of x's for which Person(x) is true: {x/John, x/Richard,...}

# Knowledge engineering in FOL

- 1. Identify the task
- 2. Assemble the relevant knowledge
- 3. Decide on a vocabulary of predicates, functions, and constants
- 4. Encode general knowledge about the domain
- 5. Encode a description of the specific problem instance
- 6. Pose queries to the inference procedure and get answers
- 7. Debug the knowledge base

## FOL Version of Wumpus World

• Typical percept sentence:

Percept([Stench,Breeze,Glitter,None,None],5)

• Actions:

Turn(Right), Turn(Left), Forward, Shoot, Grab, Release, Climb

• To determine best action, construct query:

∃ a BestAction(a,5)

- ASK solves this and returns {a/Grab}
  - And TELL about the action.

### **Knowledge Base for Wumpus World**

### • Perception

- $\forall$ s,g,x,y,t Percept([s,Breeze,g,x,y],t) ⇒ Breeze(t)
- −  $\forall$ s,b,x,y,t Percept([s,b,Glitter,x,y],t)  $\Rightarrow$  Glitter(t)

### Reflex action

- $\forall$ t Glitter(t) ⇒ BestAction(Grab,t)
- Reflex action with internal state
  - $\forall$ t Glitter(t) ∧¬Holding(Gold,t) ⇒ BestAction(Grab,t)

Holding(Gold,t) can not be observed: keep track of change.

Deducing hidden properties Environment definition:

 $\forall x,y,a,b \ Adjacent([x,y],[a,b]) \Leftrightarrow \\ [a,b] \in \{[x+1,y], [x-1,y], [x,y+1], [x,y-1]\}$ 

Properties of locations:

 $\forall$ s,t At(Agent,s,t)  $\land$  Breeze(t)  $\Rightarrow$  Breezy(s)

### Squares are breezy near a pit:

- Diagnostic rule---infer cause from effect
   ∀s Breezy(s) ⇔ ∃ r Adjacent(r,s) ∧ Pit(r)
- Causal rule---infer effect from cause (model based reasoning)
    $\forall$ r Pit(r) ⇒ [ $\forall$ s Adjacent(r,s) ⇒ Breezy(s)]

#### Keeping track of change

Facts hold in situations, rather than eternally E.g., Holding(Gold, Now) rather than just Holding(Gold)

Situation calculus is one way to represent change in FOL: Adds a situation argument to each non-eternal predicate E.g., Now in Holding(Gold, Now) denotes a situation

Situations are connected by the Result function Result(a, s) is the situation that results from doing a in s



# Yale shooting problem

- The Yale shooting problem illustrates the frame problem. (Its inventors were working at Yale University when they proposed it.)
- Fred (a turkey) is initially alive and a gun is initially unloaded. Loading the gun, waiting for a moment, and then shooting the gun at Fred is expected to kill Fred.
- However, in one solution, Fred indeed dies; in another (also logically correct) solution, the gun becomes mysteriously unloaded and Fred survives.
- By <u>Hanks</u> and <u>McDermott</u>, adapted from Wikipedia

#### Describing actions I

"Effect" axiom—describe changes due to action  $\forall s \ AtGold(s) \Rightarrow Holding(Gold, Result(Grab, s))$ 

"Frame" axiom—describe non-changes due to action  $\forall s \; HaveArrow(s) \Rightarrow HaveArrow(Result(Grab, s))$ 

Frame problem: find an elegant way to handle non-change

- (a) representation—avoid frame axioms
- (b) inference—avoid repeated "copy-overs" to keep track of state

Qualification problem: true descriptions of real actions require endless caveats what if gold is slippery or nailed down or . . .

Ramification problem: real actions have many secondary consequences what about the dust on the gold, wear and tear on gloves, ...

#### Describing actions II

Successor-state axioms solve the representational frame problem

Each axiom is "about" a predicate (not an action per se):

- $\mathsf{P} \ \mathsf{true} \ \mathsf{afterwards} \ \Leftrightarrow \ [\mathsf{an} \ \mathsf{action} \ \mathsf{made} \ \mathsf{P} \ \mathsf{true}$ 
  - $\vee$  P true already and no action made P false]

For holding the gold:

$$\begin{array}{l} \forall a,s \ Holding(Gold,Result(a,s)) \Leftrightarrow \\ [(a = Grab \wedge AtGold(s)) \\ \lor (Holding(Gold,s) \wedge a \neq Release)] \end{array}$$

### Set Theory in First-Order Logic

Can we define set theory using FOL?

- individual sets, union, intersection, etc

Answer is yes.

Basics:

- empty set = constant = { }
- unary predicate Set( ), true for sets

- binary predicates:

 $x \in S$  (true if x is a member of the set s)

 $S_1 \subseteq S_2$  (true if s1 is a subset of s2)

- binary functions:

intersection 
$${\bf S}_1 \cap {\bf S}_2$$
 , union  ${\bf S}_1 \cup {\bf S}_2$  , adjoining {x|s}

### A Possible Set of FOL Axioms for Set Theory

The only sets are the empty set and sets made by adjoining an element to a set

 $\forall s \text{ Set}(s) \Leftrightarrow (s = \{\}) \lor (\exists x, s_2 \text{ Set}(s_2) \land s = \{x \mid s_2\})$ 

The empty set has no elements adjoined to it  $\neg \exists x, s \{x \mid s\} = \{\}$ 

Adjoining an element already in the set has no effect  $\forall x, s \ x \in s \Leftrightarrow s = \{x \mid s\}$ 

The only elements of a set are those that were adjoined into it. Expressed recursively:

 $\forall x,s \quad x \in s \Leftrightarrow [\exists y,s_2 \ (s = \{y \mid s_2\} \land (x = y \lor x \in s_2))]$ 

### A Possible Set of FOL Axioms for Set Theory

A set is a subset of another set iff all the first set's members are members of the 2<sup>nd</sup> set

 $\forall s_1, s_2 \ s_1 \subseteq s_2 \Leftrightarrow (\forall x \ x \in s_1 \Longrightarrow x \in s_2)$ 

Two sets are equal iff each is a subset of the other  $\forall s_1, s_2 \ (s_1 = s_2) \Leftrightarrow (s_1 \subseteq s_2 \land s_2 \subseteq s_1)$ 

An object is in the intersection of 2 sets only if a member of both  $\forall x,s_1,s_2 \ x \in (s_1 \cap s_2) \Leftrightarrow (x \in s_1 \land x \in s_2)$ 

An object is in the union of 2 sets only if a member of either  $\forall x, s_1, s_2 \ x \in (s_1 \cup s_2) \Leftrightarrow (x \in s_1 \lor x \in s_2)$ 

One-bit full adder



Possible queries:

- does the circuit function properly?
- what gates are connected to the first input terminal?
- what would happen if one of the gates is broken?
   and so on

- 1. Identify the task
  - Does the circuit actually add properly?
- 2. Assemble the relevant knowledge
  - Composed of wires and gates; Types of gates (AND, OR, XOR, NOT)
  - \_\_\_\_
    - Irrelevant: size, shape, color, cost of gates
- 3. Decide on a vocabulary
  - Many alternative ways to say X1 is an OR gate:
  - Type(X<sub>1</sub>) = XOR (function)
     Type(X<sub>1</sub>, XOR) (binary predicate)
     XOR(X<sub>1</sub>) (unary predicate)
     etc.

- 4. Encode general knowledge of the domain
  - $\qquad \forall t_1, t_2 \text{ Connected}(t_1, t_2) \Rightarrow \text{Signal}(t_1) = \text{Signal}(t_2)$
  - $\forall t Signal(t) = 1 \lor Signal(t) = 0$
  - 1≠0
  - −  $\forall t_1, t_2 \text{ Connected}(t_1, t_2) \Rightarrow \text{Connected}(t_2, t_1)$
  - $\forall$ g Type(g) = OR ⇒ Signal(Out(1,g)) = 1 ⇔ ∃n Signal(In(n,g)) = 1
  - $\forall$ g Type(g) = AND ⇒ Signal(Out(1,g)) = 0 ⇔ ∃n Signal(In(n,g)) = 0
  - $\qquad \forall g \text{ Type}(g) = XOR \Rightarrow \text{Signal}(\text{Out}(1,g)) = 1 \Leftrightarrow \text{Signal}(\text{In}(1,g)) \neq \text{Signal}(\text{In}(2,g))$
  - $\qquad \forall g Type(g) = NOT \Rightarrow Signal(Out(1,g)) ≠ Signal(In(1,g))$

### 5. Encode the specific problem instance

Type $(X_1) = XOR$ Type $(A_1) = AND$ Type $(O_1) = OR$  Type $(X_2) = XOR$ Type $(A_2) = AND$ 

Connected(Out( $1,X_1$ ),In( $1,X_2$ )) Connected(Out( $1,X_1$ ),In( $2,A_2$ )) Connected(Out( $1,A_2$ ),In( $1,O_1$ )) Connected(Out( $1,A_1$ ),In( $2,O_1$ )) Connected(Out( $1,X_2$ ),Out( $1,C_1$ )) Connected(Out( $1,O_1$ ),Out( $2,C_1$ )) Connected( $In(1,C_1),In(1,X_1)$ ) Connected( $In(1,C_1),In(1,A_1)$ ) Connected( $In(2,C_1),In(2,X_1)$ ) Connected( $In(2,C_1),In(2,A_1)$ ) Connected( $In(3,C_1),In(2,X_2)$ ) Connected( $In(3,C_1),In(1,A_2)$ )

### 6. Pose queries to the inference procedure:

What are the possible sets of values of all the terminals for the adder circuit?

 $\exists i_1, i_2, i_3, o_1, o_2 \text{ Signal}(\text{In}(1, C_1)) = i_1 \land \text{ Signal}(\text{In}(2, C_1)) = i_2 \land \text{ Signal}(\text{In}(3, C_1)) = i_3 \land \text{ Signal}(\text{Out}(1, C_1)) = o_1 \land \text{ Signal}(\text{Out}(2, C_1)) = o_2$ 

### 7. Debug the knowledge base

May have omitted assertions like  $1 \neq 0$ 

### Review --- Knowledge engineering in FOL

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## Summary

- First-order logic:
  - Much more expressive than propositional logic
  - Allows objects and relations as semantic primitives
  - Universal and existential quantifiers
  - syntax: constants, functions, predicates, equality, quantifiers
- Knowledge engineering using FOL

   Capturing domain knowledge in logical form