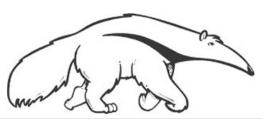
First-Order Logic C: Knowledge Engineering

CS171, Winter Quarter, 2019
Introduction to Artificial Intelligence
Prof. Richard Lathrop



Read Beforehand: R&N 8, 9.1-9.2, 9.5.1-9.5.5





Outline

- Review --- Syntactic Ambiguity
- Using FOL
 - Tell, Ask
- Example: Wumpus world
- Deducing Hidden Properties
 - Keeping track of change
 - Describing the results of Actions
- Set Theory in First-Order Logic
- Knowledge engineering in FOL
- The electronic circuits domain

You will be expected to know

Seven steps of Knowledge Engineering (R&N section 8.4.1)

 Given a simple Knowledge Engineering problem, produce a simple FOL Knowledge Base that solves the problem

Review --- Syntactic Ambiguity

- FOPC provides many ways to represent the same thing.
- E.g., "Ball-5 is red."
 - HasColor(Ball-5, Red)
 - Ball-5 and Red are objects related by HasColor.
 - Red(Ball-5)
 - Red is a unary predicate applied to the Ball-5 object.
 - HasProperty(Ball-5, Color, Red)
 - Ball-5, Color, and Red are objects related by HasProperty.
 - ColorOf(Ball-5) = Red
 - Ball-5 and Red are objects, and ColorOf() is a function.
 - HasColor(Ball-5(), Red())
 - Ball-5() and Red() are functions of zero arguments that both return an object, which objects are related by HasColor.
 - **–** ...
- This can GREATLY confuse a pattern-matching reasoner.
 - Especially if multiple people collaborate to build the KB, and they all have different representational conventions.

Review --- Syntactic Ambiguity --Partial Solution

- FOL can be TOO expressive, can offer TOO MANY choices
- Likely confusion, especially for teams of Knowledge Engineers
- Different team members can make different representation choices
 - E.g., represent "Ball43 is Red." as:
 - a predicate (= verb)? E.g., "Red(Ball43)"?
 - an object (= noun)? E.g., "Red = Color(Ball43))"?
 - a property (= adjective)? E.g., "HasProperty(Ball43, Red)"?

PARTIAL SOLUTION:

- An upon-agreed ontology that settles these questions
- Ontology = what exists in the world & how it is represented
- The Knowledge Engineering teams agrees upon an ontology BEFORE they begin encoding knowledge

Using FOL

We want to TELL things to the KB, e.g.

```
TELL(KB, \forall x King(x) \Rightarrow PersonX) )
TELL(KB, King(John))
```

These sentences are assertions

 We also want to ASK things to the KB, ASK(KB, ∃ x Person(x))

these are queries or goals

The KB should return the list of x's for which Person(x) is true: {x/John, x/Richard,...}

Knowledge engineering in FOL

- 1. Identify the task
- 2. Assemble the relevant knowledge
- 3. Decide on a vocabulary of predicates, functions, and constants
- 4. Encode general knowledge about the domain
- 5. Encode a description of the specific problem instance
- 6. Pose queries to the inference procedure and get answers
- 7. Debug the knowledge base

FOL Version of Wumpus World

Typical percept sentence:

Percept([Stench,Breeze,Glitter,None,None],5)

Actions:

Turn(Right), Turn(Left), Forward, Shoot, Grab, Release, Climb

To determine best action, construct query:

∃ a BestAction(a,5)

- ASK solves this and returns {a/Grab}
 - And TELL about the action.

Knowledge Base for Wumpus World

Perception

- ∀s,g,x,y,t Percept([s,Breeze,g,x,y],t) ⇒ Breeze(t)
- \forall s,b,x,y,t Percept([s,b,Glitter,x,y],t) \Rightarrow Glitter(t)

Reflex action

- \forall t Glitter(t) \Rightarrow BestAction(Grab,t)

Reflex action with internal state

- \forall t Glitter(t) \land ¬Holding(Gold,t) \Rightarrow BestAction(Grab,t)

Holding(Gold,t) can not be observed: keep track of change.

Deducing hidden properties

Environment definition:

```
\forall x,y,a,b Adjacent([x,y],[a,b]) \Leftrightarrow [a,b] \in {[x+1,y], [x-1,y],[x,y+1],[x,y-1]}
```

Properties of locations:

```
\foralls,t At(Agent,s,t) \land Breeze(t) \Rightarrow Breezy(s)
```

Squares are breezy near a pit:

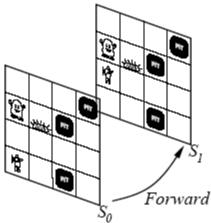
- Diagnostic rule---infer cause from effect
 ∀s Breezy(s) ⇔ ∃ r Adjacent(r,s) ∧ Pit(r)
- Causal rule---infer effect from cause (model based reasoning)
 \forall r Pit(r) ⇒ [\forall s Adjacent(r,s) ⇒ Breezy(s)]

Keeping track of change

Facts hold in situations, rather than eternally E.g., Holding(Gold, Now) rather than just Holding(Gold)

Situation calculus is one way to represent change in FOL: Adds a situation argument to each non-eternal predicate E.g., Now in Holding(Gold, Now) denotes a situation

Situations are connected by the Result function Result(a,s) is the situation that results from doing a in s



Yale shooting problem

- The Yale shooting problem illustrates the frame problem. (Its inventors were working at Yale University when they proposed it.)
- Fred (a turkey) is initially alive and a gun is initially unloaded. Loading the gun, waiting for a moment, and then shooting the gun at Fred is expected to kill Fred.
- However, in one solution, Fred indeed dies; in another (also logically correct) solution, the gun becomes mysteriously unloaded and Fred survives.
- By Hanks and McDermott, adapted from Wikipedia

Describing actions I

```
"Effect" axiom—describe changes due to action
\forall s \ AtGold(s) \Rightarrow Holding(Gold, Result(Grab, s))
```

"Frame" axiom—describe non-changes due to action $\forall s \; HaveArrow(s) \Rightarrow HaveArrow(Result(Grab, s))$

Frame problem: find an elegant way to handle non-change

- (a) representation—avoid frame axioms
- (b) inference—avoid repeated "copy-overs" to keep track of state

Qualification problem: true descriptions of real actions require endless caveats what if gold is slippery or nailed down or . . .

Ramification problem: real actions have many secondary consequences what about the dust on the gold, wear and tear on gloves, ...

Describing actions II

Successor-state axioms solve the representational frame problem

Each axiom is "about" a predicate (not an action per se):

```
P true afterwards \Leftrightarrow [an action made P true \lor P true already and no action made P false]
```

For holding the gold:

```
\forall a, s \; Holding(Gold, Result(a, s)) \Leftrightarrow
[(a = Grab \land AtGold(s))
\lor (Holding(Gold, s) \land a \neq Release)]
```

Set Theory in First-Order Logic

Can we define set theory using FOL?

- individual sets, union, intersection, etc

Answer is yes.

Basics:

- empty set = constant = { }
- unary predicate Set(), true for sets
- binary predicates:

$$x \in S$$
 (true if x is a member of the set s)

$$S_1 \subseteq S_2$$
 (true if s1 is a subset of s2)

- binary functions:

intersection
$$\mathsf{S}_1 \cap \mathsf{S}_2$$
, union $\mathsf{S}_1 \cup \mathsf{S}_2$, adjoining $\{\mathsf{x} \mid \mathsf{s}\}$

A Possible Set of FOL Axioms for Set Theory

The only sets are the empty set and sets made by adjoining an element to a set

$$\forall s \text{ Set}(s) \Leftrightarrow (s = \{\}) \lor (\exists x, s_2 \text{ Set}(s_2) \land s = \{x \mid s_2\})$$

The empty set has no elements adjoined to it

$$\neg \exists x,s \{x \mid s\} = \{\}$$

Adjoining an element already in the set has no effect

$$\forall x, s \ x \in s \Leftrightarrow s = \{x \mid s\}$$

The only elements of a set are those that were adjoined into it. Expressed recursively:

$$\forall x,s \quad x \in s \Leftrightarrow [\exists y,s_2 \ (s = \{y \mid s_2\} \land (x = y \lor x \in s_2))]$$

A Possible Set of FOL Axioms for Set Theory

A set is a subset of another set iff all the first set's members are members of the 2nd set

$$\forall s_1, s_2 \ s_1 \subseteq s_2 \Leftrightarrow (\forall x \ x \in s_1 \Rightarrow x \in s_2)$$

Two sets are equal iff each is a subset of the other

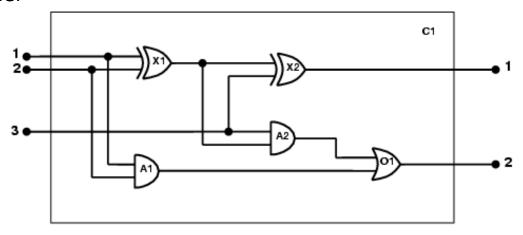
$$\forall s_1, s_2 \ (s_1 = s_2) \Leftrightarrow (s_1 \subseteq s_2 \land s_2 \subseteq s_1)$$

An object is in the intersection of 2 sets only if a member of both

$$\forall x, s_1, s_2 \ x \in (s_1 \cap s_2) \Leftrightarrow (x \in s_1 \land x \in s_2)$$

An object is in the union of 2 sets only if a member of either $\forall x,s_1,s_2 \ x \in (s_1 \cup s_2) \Leftrightarrow (x \in s_1 \vee x \in s_2)$

One-bit full adder



Possible queries:

- does the circuit function properly?
- what gates are connected to the first input terminal?
- what would happen if one of the gates is broken? and so on

1. Identify the task

Does the circuit actually add properly?

2. Assemble the relevant knowledge

- Composed of wires and gates; Types of gates (AND, OR, XOR, NOT)
- Irrelevant: size, shape, color, cost of gates

3. Decide on a vocabulary

Many alternative ways to say X1 is an OR gate:

Type(X₁) = XOR (function)
 Type(X₁, XOR) (binary predicate)
 XOR(X₁) (unary predicate)
 etc.

4. Encode general knowledge of the domain

```
\forall t_1, t_2 \text{ Connected}(t_1, t_2) \Rightarrow \text{Signal}(t_1) = \text{Signal}(t_2)
\forallt Signal(t) = 1 \vee Signal(t) = 0
1 ≠ 0
\forall t_1, t_2 \text{ Connected}(t_1, t_2) \Rightarrow \text{Connected}(t_2, t_1)
\forall g \text{ Type}(g) = OR \Rightarrow Signal(Out(1,g)) = 1 \Leftrightarrow \exists n \text{ Signal}(In(n,g)) = 1
\forall g \text{ Type}(g) = AND \Rightarrow Signal(Out(1,g)) = 0 \Leftrightarrow \exists n \text{ Signal}(In(n,g)) = 0
\forall g \text{ Type}(g) = XOR \Rightarrow Signal(Out(1,g)) = 1 \Leftrightarrow Signal(In(1,g)) \neq Signal(In(2,g))
\forall g \text{ Type}(g) = \text{NOT} \Rightarrow \text{Signal}(\text{Out}(1,g)) \neq \text{Signal}(\text{In}(1,g))
```

5. Encode the specific problem instance

```
Type(X_1) = XOR Type(X_2) = XOR

Type(A_1) = AND Type(A_2) = AND

Type(O_1) = OR
```

 $\begin{array}{lll} \text{Connected}(\text{Out}(1, X_1), \text{In}(1, X_2)) & \text{Connected}(\text{In}(1, C_1), \text{In}(1, X_1)) \\ \text{Connected}(\text{Out}(1, X_1), \text{In}(2, A_2)) & \text{Connected}(\text{In}(1, C_1), \text{In}(1, A_1)) \\ \text{Connected}(\text{Out}(1, A_2), \text{In}(1, O_1)) & \text{Connected}(\text{In}(2, C_1), \text{In}(2, X_1)) \\ \text{Connected}(\text{Out}(1, A_1), \text{In}(2, O_1)) & \text{Connected}(\text{In}(2, C_1), \text{In}(2, A_1)) \\ \text{Connected}(\text{Out}(1, X_2), \text{Out}(1, C_1)) & \text{Connected}(\text{In}(3, C_1), \text{In}(2, X_2)) \\ \text{Connected}(\text{Out}(1, O_1), \text{Out}(2, C_1)) & \text{Connected}(\text{In}(3, C_1), \text{In}(1, A_2)) \\ \end{array}$

6. Pose queries to the inference procedure:

What are the possible sets of values of all the terminals for the adder circuit?

$$\exists i_1, i_2, i_3, o_1, o_2 \text{ Signal}(In(1,C_1)) = i_1 \land \text{Signal}(In(2,C_1)) = i_2 \land \text{Signal}(In(3,C_1)) = i_3 \land \text{Signal}(Out(1,C_1)) = o_1 \land \text{Signal}(Out(2,C_1)) = o_2$$

7. Debug the knowledge base

May have omitted assertions like $1 \neq 0$

Review --- Knowledge engineering in FOL

- 1. Identify the task
- 2. Assemble the relevant knowledge
- 3. Decide on a vocabulary of predicates, functions, and constants
- 4. Encode general knowledge about the domain
- 5. Encode a description of the specific problem instance
- 6. Pose queries to the inference procedure and get answers
- 7. Debug the knowledge base

Summary

- First-order logic:
 - Much more expressive than propositional logic
 - Allows objects and relations as semantic primitives
 - Universal and existential quantifiers
 - syntax: constants, functions, predicates, equality, quantifiers
- Knowledge engineering using FOL
 - Capturing domain knowledge in logical form